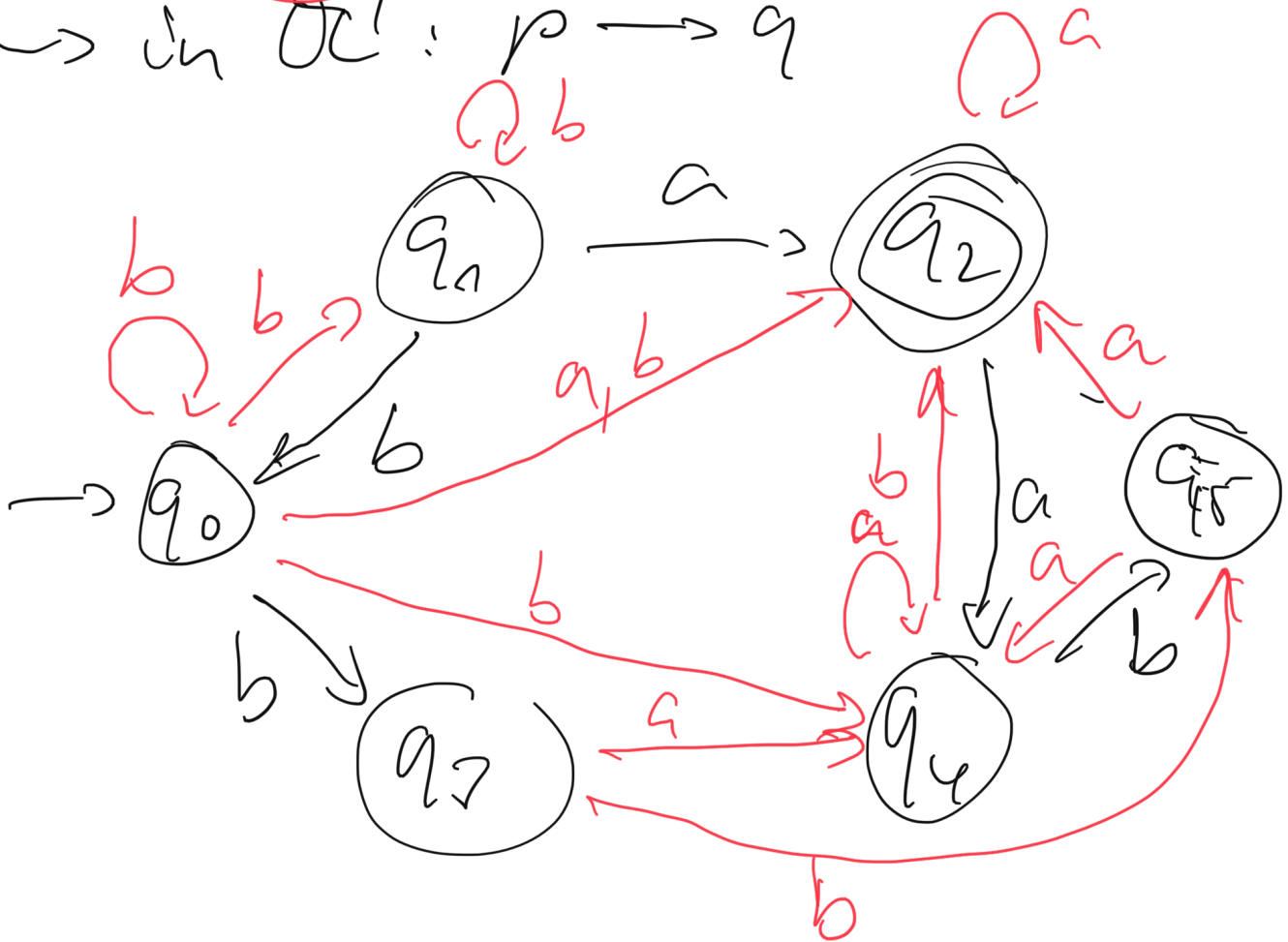
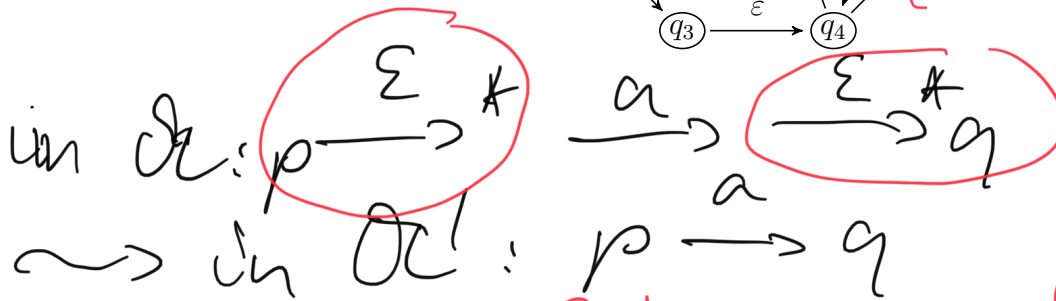
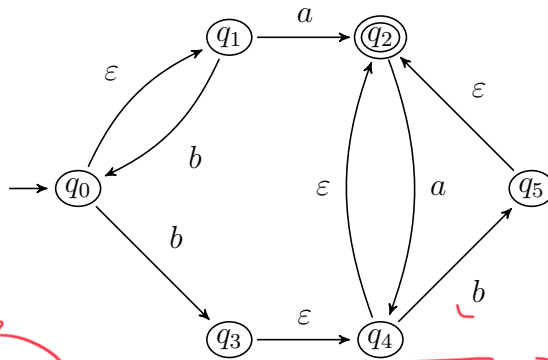


A6: Elimination of ε -Transitions

Task: Eliminate all ε -transitions of the following ε -NFA \mathfrak{A} over $\Sigma := \{a, b\}$ to obtain an equivalent NFA.



A8: Construction of Regular Expressions

Task: Give regular expressions that describe the following languages.

(a) $L := \{w \in \{a, b\}^* \mid |w| \text{ divisible by } 3\}$

(b) $L := \{w \in \{a, b, c\}^* \mid w \text{ does not contain } a, b, \text{ or } c\}$

(c) $L := \{w \in \{a, b\}^* \mid \text{substring } ab \text{ occurs exactly twice in } w, \text{ but not at the end}\}$

Reg. expr.: $\emptyset, \Sigma, a \in \Sigma$

$\alpha | \beta, \alpha \cdot \beta, \alpha^*$

(a) $((a|b)(a|b)(a|b))^*$ Σ, aaa, aba, ...
aaabbb, ...

(b) $(b|c)^* | (a|c)^* | (a|b)^*$

(c) $b^* \underbrace{a^+}_{1st} b^+ \underbrace{a^+}_{2nd} b^+ (b^+ a^* | b^* a^+)$

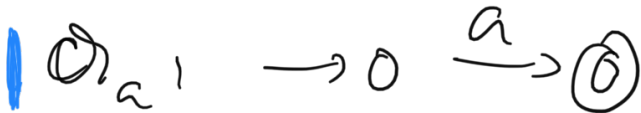
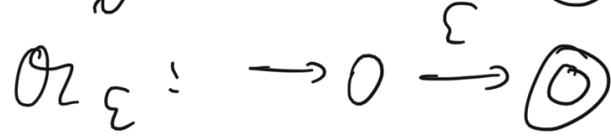
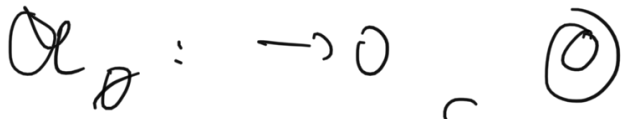
where $a^+ := a a^*$

$b^* \underbrace{a^+ b^+}_{ab} \underbrace{a^+ b^+}_{ab} \underbrace{b^* a^+}_{\Sigma}$ y

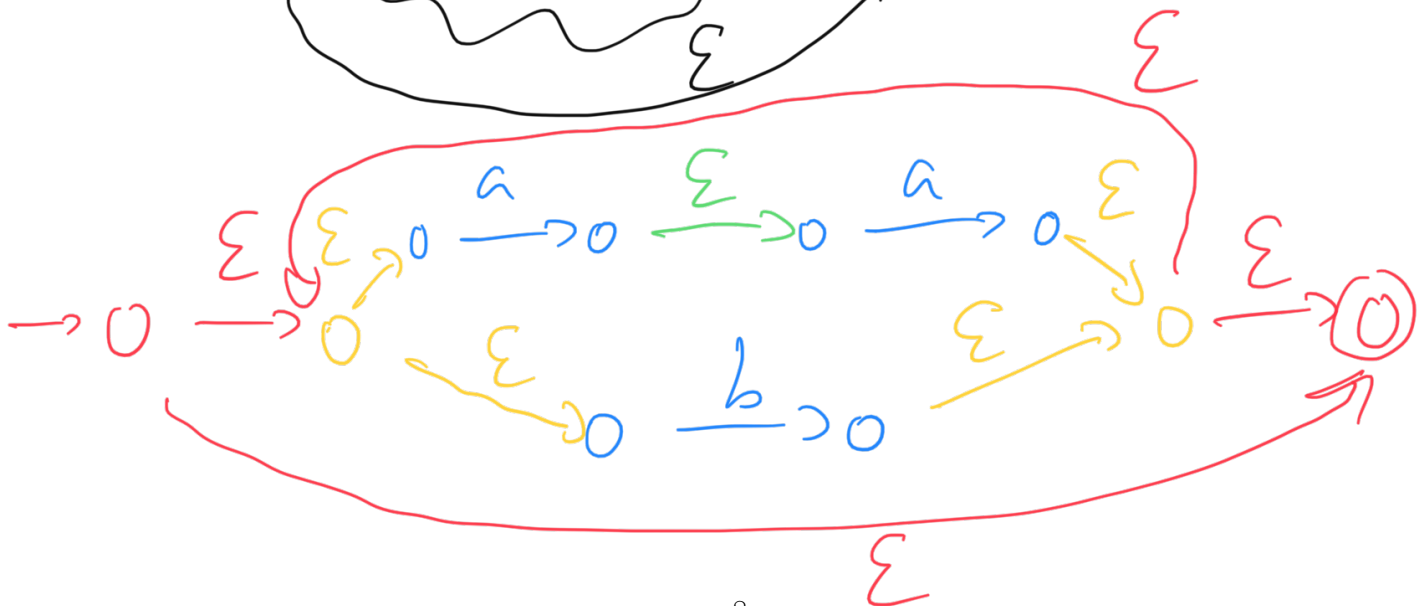
A9: From Regular Expressions to Finite Automata

Task: Using Kleene's construction, give the ϵ -NFA for the regular expression $(aa | b)^*$.

General: reg. expr. $\alpha \rightsquigarrow \epsilon\text{-NFA } \mathcal{A}_\alpha$



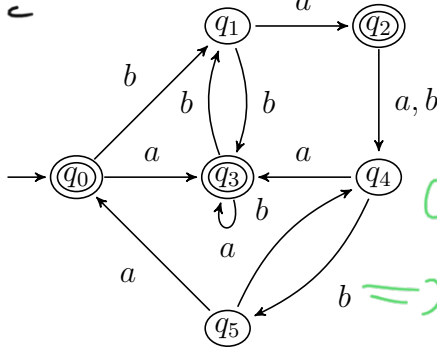
$$L^* = \bigcup_{n \geq 0} L^n$$



A10: Minimisation of Deterministic Finite Automata

Task: Minimise the following DFA.

(i) $p \in F, q \notin F$
 $\rightarrow p, q$ disting.
 by Σ

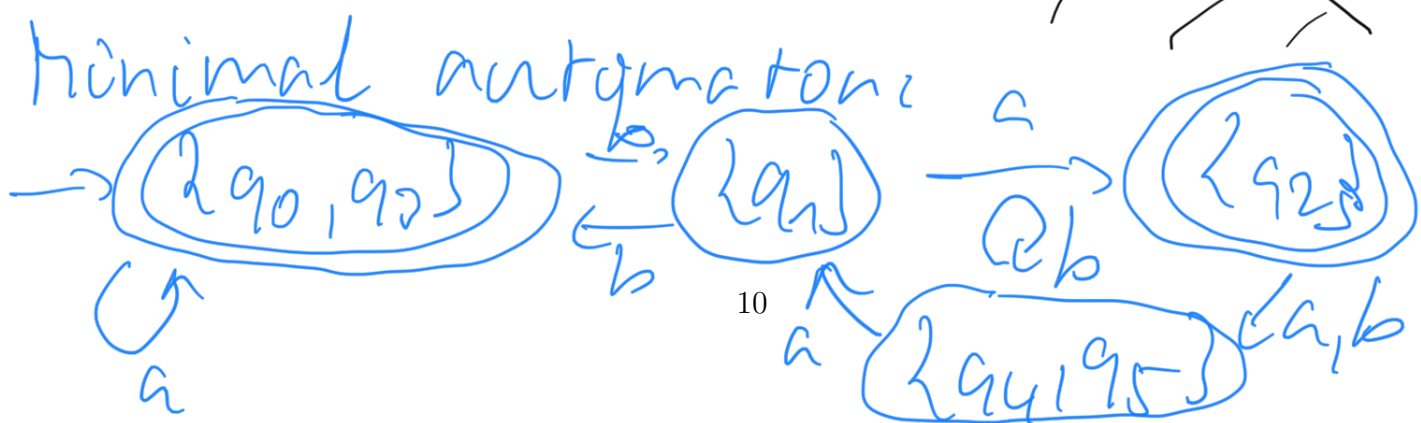


(ii) $\delta(q, a) = q'$
 $\delta(p, a) = p'$
 q', p' disting. by w
 $\Rightarrow q, p$ disting. by aw

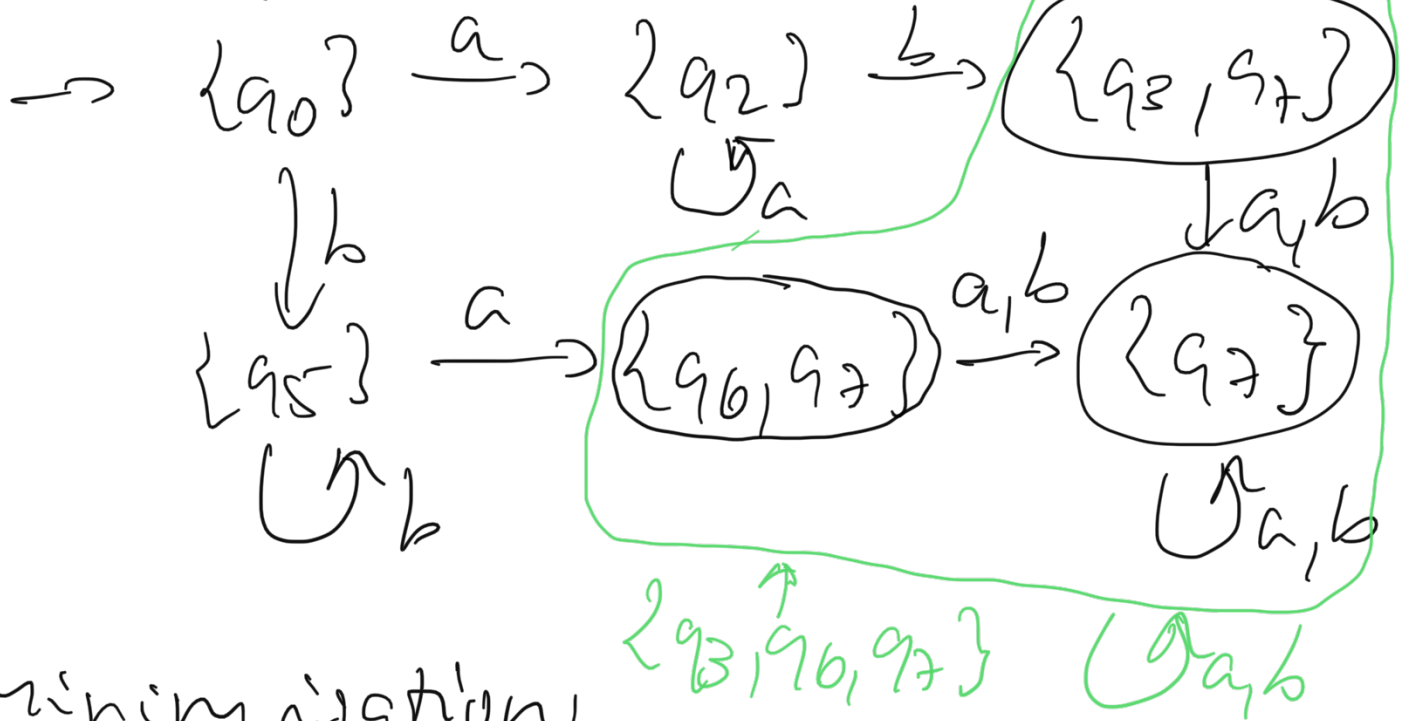
$p \sim q \hat{=} \forall w \in \Sigma^*: \delta^*(p, w) \in F \Leftrightarrow \delta^*(q, w) \in F$

	q_0	q_1	q_2	q_3	q_4	q_5
q_0	Σ	a	\emptyset	Σ	Σ	Σ
q_1		Σ	Σ	b	b	b
q_2			a	Σ	Σ	Σ
q_3				Σ	Σ	Σ
q_4					\emptyset	\emptyset
q_5						\emptyset

$\rightarrow q_0 \sim q_3$
 $q_4 \sim q_5$



Power set construction:



Minimization:

	q_0	q_2	q_5	q_3, q_7	q_6, q_7	q_7
q_0		b	a	Σ	Σ	Σ
q_2			a	Σ	Σ	Σ
q_5				Σ	Σ	Σ
q_3, q_7					\times	\times
q_6, q_7						\times
q_7						

$\Rightarrow q_3, q_7 \sim q_6, q_7 \sim q_7$