

Model Checking

(Winter Term 2019/2020)

— Exercise Sheet 7 (due 20th December) —

General Remarks

- The exercises are to be solved in groups of *three* students.
- You may hand in your solutions for the exercises just before the exercise class starts at 10:30 or by dropping them into the “Model Checking” box at our chair *before 10:25*. Do *not* hand in your solutions via Moodle or via e-mail.

Exercise 1★

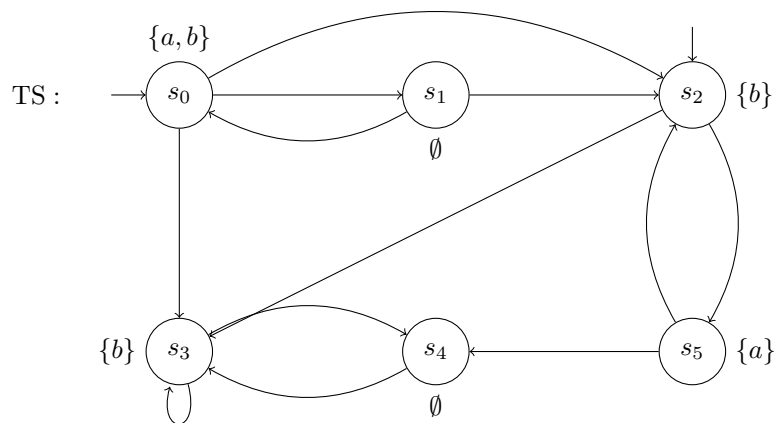
(3+2+3 Points)

Consider the following CTL formula Φ and the fairness assumption $sfair$:

$$\Phi = \forall \square \forall \diamond a$$

$$sfair = \square \diamond \underbrace{(b \wedge \neg a)}_{\Phi_1} \rightarrow \square \diamond \underbrace{\exists (b \cup (a \wedge \neg b))}_{\Psi_1}$$

and transition system TS over $AP = \{a, b\}$ which is given below.



Here, we abstract from the actions in TS as they are not relevant to the task.

- Determine $Sat(\Phi_1)$ and $Sat(\Psi_1)$ (without fairness). Justify your answer.
- Determine $Sat_{sfair}(\exists \square true)$. Justify your answer.
- Determine $Sat_{sfair}(\Phi)$. Justify your answer.

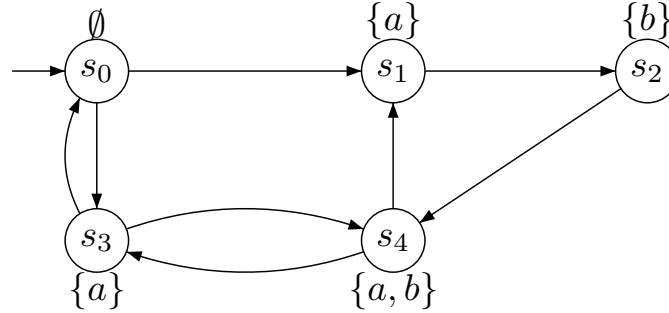
Exercise 2

(6 Points)

Consider the CTL-formula $\Phi = \forall \square (a \rightarrow \forall \diamond (b \wedge \neg a))$ together with the following CTL fairness assumption

$$\begin{aligned} \text{fair} = & \square \diamond \forall \bigcirc (a \wedge \neg b) \rightarrow \square \diamond \forall \bigcirc (b \wedge \neg a) \\ & \wedge \diamond \square \exists \diamond b \rightarrow \square \diamond b. \end{aligned}$$

Check whether $TS \models_{\text{fair}} \Phi$ for the transition system TS below.



Exercise 3★

(3 Points)

For each of the following, give pairs of transition systems such that:

- T_1, T_2 are simulation-equivalent, but not bisimilar.
- T_3, T_4 are trace-equivalent, but not bisimilar.
- T_5, T_6 are finite-trace equivalent, but not simulation-equivalent.
- T_7, T_8 are finite-trace equivalent, but not trace-equivalent.
- T_9, T_{10} are trace-equivalent and simulation-equivalent.

Exercise 4

(3 Points)

Let $TS = (S, Act, \rightarrow, I, AP, L)$ be a transition system. The relations $\sim_n \subseteq S \times S$ are inductively defined by:

- $s_1 \sim_0 s_2$ iff $L(s_1) = L(s_2)$.
- $s_1 \sim_{n+1} s_2$ iff:
 - $L(s_1) = L(s_2)$,
 - for all $s'_1 \in Post(s_1)$ there exists $s'_2 \in Post(s_2)$ with $s'_1 \sim_n s'_2$,
 - for all $s'_2 \in Post(s_2)$ there exists $s'_1 \in Post(s_1)$ with $s'_1 \sim_n s'_2$.

Questions:

- (a) Show that for *finite* TS it holds that $\sim_{TS} = \bigcap_{n \geq 0} \sim_n$, i.e.,

$$s_1 \sim_{TS} s_2 \text{ iff } s_1 \sim_n s_2 \text{ for all } n \geq 0$$

- (b) Does this also hold for infinite transition systems (provide a proof or a counterexample)?