

Model Checking (Winter Term 2019/2020)

— Exercise Sheet 6 (due 13th December) —

General Remarks

- The exercises are to be solved in groups of *three* students.
- You may hand in your solutions for the exercises just before the exercise class starts at 10:30 or by dropping them into the “Model Checking” box at our chair *before 10:25*. Do *not* hand in your solutions via Moodle or via e-mail.

Exercise 1★

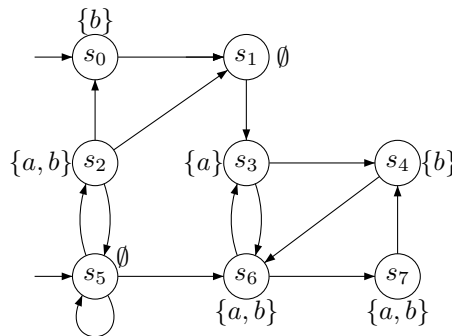
- (a) Give a transition system TS without terminal states that contains two states s_1 and s_2 such that $s_1 \not\equiv_{LTL} s_2$ and there is *no* LTL formula φ with $s_2 \models \varphi$ and $s_1 \not\models \varphi$.
- (b) Let TS_1 and TS_2 be transition systems over AP without terminal states such that $TS_1 \not\equiv_{CTL} TS_2$. Prove or disprove: there exists a CTL formula Φ over AP such that $TS_1 \models \Phi$ and $TS_2 \not\models \Phi$.

Exercise 2

5 Consider the CTL*-formula (with derived operators) over $AP = \{a, b\}$

$$\Phi = \forall \diamond \square \exists \bigcirc (a \cup \exists \square b)$$

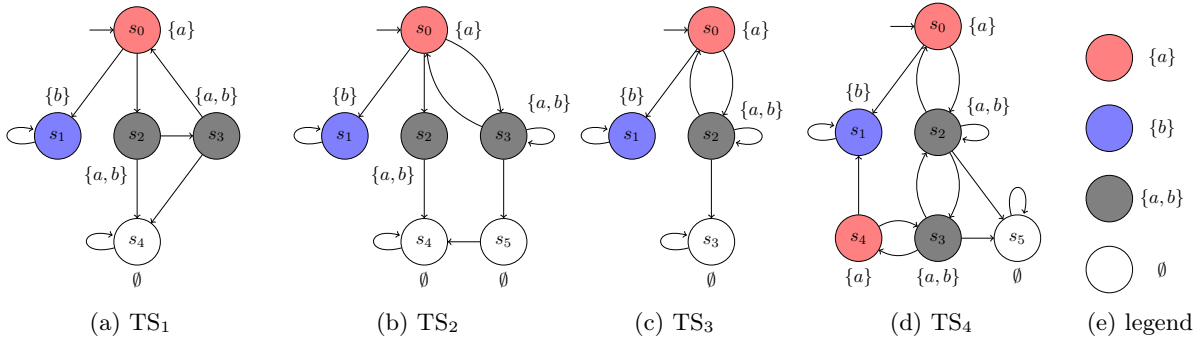
and the transition system TS outlined below:



Apply the CTL* Model Checking algorithm to compute $Sat(\Phi)$ and decide whether $TS \models \Phi$.
Hint: You may infer the satisfaction sets for LTL formulas directly.

Exercise 3

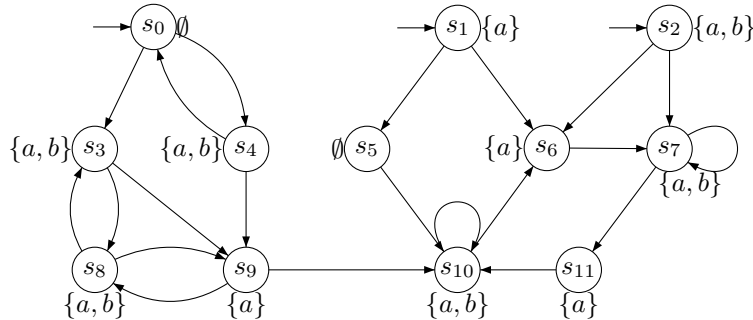
Consider the following transition systems TS_1, \dots, TS_4 .



- Which transition systems are trace equivalent? Justify your answers by either providing the set of traces or a counterexample trace.
- Which transition systems are bisimulation equivalent? Justify your answers by either providing a bisimulation relation or a CTL formula that distinguishes the considered transition systems.

Exercise 4

Consider the transition system TS over $AP = \{a, b\}$ outlined below:



- Determine the bisimulation equivalence \sim_{TS} and depict the bisimulation quotient system TS/\sim .
- For each bisimulation equivalence class C , provide a CTL formula Φ_C that holds only in the states in C .