

# Model Checking (Winter Term 2019/2020)

## — Exercise Sheet 3 (due 15th November) —

### General Remarks

- The exercises are to be solved in groups of *three* students.
- You may hand in your solutions before the exercise class starts at 10:30 or by dropping them into the “Model Checking” box at our chair *before 10:20*. Do *not* hand in your solutions electronically.

### Exercise 1

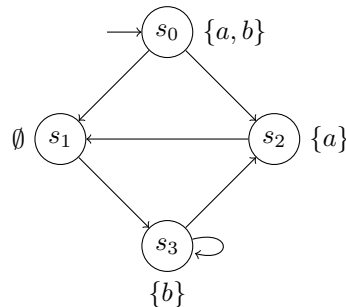
**(5 Points)**

Let the  $\omega$ -regular LT properties  $P_1$  and  $P_2$  over the set of atomic propositions  $AP = \{a, b\}$  be given by

$P_1 :=$  “if  $a$  holds infinitely often, then  $b$  holds finitely often”

$P_2 :=$  “ $a$  holds infinitely often and  $b$  holds infinitely often”

The model is given by the transition system TS as follows:



Algorithmically check whether  $TS \models P_1$  and  $TS \models P_2$ . For this, proceed as follows.

- Derive *suitable* NBA  $\mathcal{A}_{P_1}$ ,  $\mathcal{A}_{P_2}$ , where suitable means “appropriate for part b)-d)”.  
*Hint:* For  $P_1$  you can find an automaton with 3 states and for  $P_2$  4 states suffice. Derive the automata directly.
- Outline the reachable fragments of the product transition systems  $TS \otimes \mathcal{A}_{P_1}$  and  $TS \otimes \mathcal{A}_{P_2}$ .
- Decide whether  $TS \models P_1$  by checking an appropriate persistence property via nested depth-first search on  $TS \otimes \mathcal{A}_{P_1}$ . Document *all* changes to the contents of  $U$ ,  $V$ ,  $\pi$  and  $\xi$  (the state sets and stacks of the nested depth-first search, see lecture). If the property is violated, provide a counterexample *based on the execution of the algorithm*.
- Decide whether  $TS \models P_2$  by checking an appropriate persistence property via SCC analysis on  $TS \otimes \mathcal{A}_{P_2}$ . If the property is violated, provide a counterexample *based on your analysis*.

### Exercise 2

**(3 Points)**

Provide an example for a liveness property  $P_{live}$  that is *not*  $\omega$ -regular. Show that  $P_{live}$  is indeed a liveness property and prove that  $P_{live}$  is not  $\omega$ -regular.

*Hint:* Think about words of the form  $\{a\} \{b\} \{a\} \{a\} \{b\} \{a\} \{a\} \{a\} \{a\} \{b\} \dots$

### Exercise 3

(2+2+2+1 Points)

- (a) Provide NBA  $\mathcal{A}_1$  and  $\mathcal{A}_2$  for the languages given by the  $\omega$ -regular expressions  $\alpha_1 = (AC + B)^*B^\omega$  and  $\alpha_2 = (B^*AC)^\omega$ .
- (b) Apply the product construction to obtain a GNBA  $\mathcal{G}$  with  $\mathcal{L}_\omega(\mathcal{G}) = \mathcal{L}_\omega(\mathcal{A}_1) \cap \mathcal{L}_\omega(\mathcal{A}_2)$ .
- (c) Transform the GNBA  $\mathcal{G}$  into an NBA  $\mathcal{A}$  with  $\mathcal{L}_\omega(\mathcal{A}) = \mathcal{L}_\omega(\mathcal{G})$ .
- (d) Justify, why  $\mathcal{L}_\omega(\mathcal{G}) = \emptyset$  on the level of the GNBA  $\mathcal{G}$ .

*Hint:* For a GNBA  $\mathcal{G} = (Q, \Sigma, \delta, Q_0, \mathcal{F})$  with at least one element in  $\mathcal{F} = \{F_1, \dots, F_k\}$ . Let  $\mathcal{A} = (Q', \Sigma, \delta', Q'_0, F')$  be an NBA with

- $Q' = Q \times \{1, \dots, k\}$ ,
- $Q'_0 = Q_0 \times \{1\}$ ,
- $F' = F_1 \times \{1\}$ , and

for all  $A \in \Sigma$  it is

$$\delta'(\langle q, i \rangle, A) = \begin{cases} \{\langle q', i \rangle \mid q' \in \delta(q, A)\} & \text{if } q \notin F_i \\ \{\langle q', (i \bmod k) + 1 \rangle \mid q' \in \delta(q, A)\} & \text{if } q \in F_i. \end{cases}$$

Then  $\mathcal{L}_\omega(\mathcal{G}) = \mathcal{L}_\omega(\mathcal{A})$ .

### Exercise 4★

(4 Points)

Recall the following LT properties from exercise sheet 1.

- (i) **“Winter is coming.”**  
 $P_1 = \emptyset^* \{winter\} (2^{\text{AP}})^\omega$
- (ii) **“Everything is awesome.”**  
 $P_2 = \{awesome\}^\omega$
- (iii) **“I’ll be back.”**  
 $P_3 = \{here\}^+ \emptyset^+ \{here\}^+ (2^{\text{AP}})^\omega$
- (iv) **“You either die a hero, or you live long enough to see yourself become the villain.”**  
 $P_4 = \{live, hero\}^+ \{hero\} (2^{\text{AP}})^\omega + \{live, hero\}^+ \{live\} (2^{\text{AP}})^\omega$
- (v) **“By night one way, by day another  
 Thus shall be the norm  
 Till you receive true love’s kiss  
 then, take love’s true form.”**  
 $P_5 = ((\{form_1\} \{day, form_2\})^+ + \{form_1\} (\{day, form_2\} \{form_1\})^*) \{kiss, true\_form\} (\{true\_form\} \{true\_form, day\})^\omega$
- (vi) **“A Lannister always pays his debts.”**  
 $P_6 = \emptyset^* (\{in\_debt\}^+ \emptyset^+)^* \emptyset^\omega$
- (vii) **“Anything is possible [if you just believe]”**  
 $P_7 = (2^{\text{AP}})^\omega$
- (viii) **“It’s gonna be legen... wait for it... dary!”**  
 $P_8 = \{legen\} \{wait\_for\_it\}^+ \{dary\} (2^{\text{AP}})^\omega$

Express each property  $P_i$  as an LTL formula  $\varphi_i$ .