

Compiler Construction 2018/19

— Exercise Sheet 0 —

General Remarks

- This exercise sheet is meant to recap topics of automata theory and will not be graded.
- The solution for this exercise sheet will be presented in the first exercise class on October 15th.
- You are expected to hand in your solutions in groups of four. If you are looking for a group or your group has less than four members, please post in the L2P forum or come to the first exercise class.

Exercise 1

(0 Points)

Which of the following statements hold?

- Deterministic finite automata (DFA) are strictly less expressive than regular expressions.
- Non-deterministic finite automata (NFA) are strictly more expressive than DFA.
- The languages of regular expressions are closed under:
 - union,
 - intersection,
 - complement,
 - concatenation,
 - Kleene closure.
- Context Free Languages (CFL) are closed under:
 - union,
 - intersection,
 - complement,
 - concatenation,
 - Kleene closure.
- DCFL is the set of context free languages that are accepted by deterministic push down automata. Is $DCFL = CFL$?

Exercise 2

(0 Points)

- Describe the language of the following regular expression in words:

$$r = (0 + 1)^*0(0 + 1)^*0(0 + 1)^*.$$

- Construct the regular expression for...
 - the set of all strings with at most one pair of consecutive 0's and at most one pair of consecutive 1's,
 - the set of all strings with equal number of 0's and 1's such that no prefix has two more 0's than 1's nor two more 1's than 0's.

- (c) Construct a context free grammar (CFG) for a set of strings of $\{(,)\}^*$ such that every string of the set has equal number of left and right parenthesis, and every prefix has at least as many left parenthesis as right parenthesis.

Exercise 3

(0 Points)

- (a) Let r and s be regular expressions. Consider the set X such that $X = r.X + s$. Under the assumption that the language of r does not contain ε (i.e., $\varepsilon \notin L(r)$), find X .
- (b) (i) Show that the language $L = \{0^{i^2} \mid i \in \mathbb{N}\}$ is not regular.
- (ii) Show that the language $L = \{a^i b^i c^i \mid i \in \mathbb{N}\}$ is not a CFL.