Semantics and Verification of Software

Winter Semester 2017/18

Lecture 17: Nondeterminism and Parallelism II (Channel Communication)

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http://moves.rwth-aachen.de/teaching/ws-1718/sv-sw/

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Outline of Lecture 17

Recap: Shared-Variables Communication

Channel Communication

CSP Examples

Fairness in CSP

Summary: Nondeterminism and Parallelism

The ParWHILE Language

Definition (Syntax of ParWHILE)

$$
\begin{aligned}
a &::= z \mid x \mid a_1\texttt{+}a_2 \mid a_1\texttt{-}a_2 \mid a_1\texttt{*}a_2 \in AExp \\
b &::= t \mid a_1\texttt{=}a_2 \mid a_1\texttt{>}a_2 \mid \neg b \mid b_1 \wedge b_2 \mid b_1 \vee b_2 \in BExp \\
c &::= \texttt{skip} \mid x \texttt{ := } a \mid c_1\texttt{;}c_2 \mid \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2 \texttt{ end} \mid \texttt{while } b \texttt{ do } c \texttt{ end} \mid \\
&\quad\; c_1 \parallel c_2 \in Cmd
\end{aligned}
$$

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Semantics of ParWHILE

Definition (Small-step execution relation for ParWHILE)

The small-step execution relation, $\to_1 \subseteq (Cmd \times \Sigma) \times (Cmd \times \Sigma)$, is defined by the following rules:

$$\frac{}{\langle \texttt{skip}, \sigma\rangle \to_1 \langle \downarrow, \sigma\rangle} \qquad \frac{\langle a, \sigma\rangle \to z}{\langle x := a, \sigma\rangle \to_1 \langle \downarrow, \sigma[x \mapsto z]\rangle}$$

$$\frac{\langle c_1, \sigma\rangle \to_1 \langle c_1', \sigma'\rangle}{\langle c_1; c_2, \sigma\rangle \to_1 \langle c_1'; c_2, \sigma'\rangle} \qquad \frac{\langle b, \sigma\rangle \to \mathsf{true}}{\langle \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2 \texttt{ end}, \sigma\rangle \to_1 \langle c_1, \sigma\rangle}$$

$$\frac{\langle b, \sigma\rangle \to \mathsf{false}}{\langle \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2 \texttt{ end}, \sigma\rangle \to_1 \langle c_2, \sigma\rangle} \qquad \frac{\langle b, \sigma\rangle \to \mathsf{false}}{\langle \texttt{while } b \texttt{ do } c \texttt{ end}, \sigma\rangle \to_1 \langle \downarrow, \sigma\rangle}$$

$$\frac{\langle b, \sigma\rangle \to \mathsf{true}}{\langle \texttt{while } b \texttt{ do } c \texttt{ end}, \sigma\rangle \to_1 \langle c; \texttt{while } b \texttt{ do } c \texttt{ end}, \sigma\rangle}$$

$$\frac{\langle c_1, \sigma\rangle \to_1 \langle c_1', \sigma'\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \to_1 \langle c_1' \parallel c_2, \sigma'\rangle} \qquad \frac{\langle c_2, \sigma\rangle \to_1 \langle c_2', \sigma'\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \to_1 \langle c_1 \parallel c_2', \sigma'\rangle}$$

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Communicating Sequential Processes

- Approach: Communicating Sequential Processes (CSP) by T. Hoare and R. Milner

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  - processes communicate along channels
  - process can send/receive on a channel if another process simultaneously performs the complementary I/O operation

$\Rightarrow$ no buffering (synchronous communication)

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  - run a sequential program ("process")
- Communication proceeds in the following way:
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  - process can send/receive on a channel if another process simultaneously performs the complementary I/O operation

  $\Rightarrow$ no buffering (synchronous communication)
- New syntactic domains:

Channel names:	$\alpha, \beta, \gamma, \ldots \in Chn$
Input operations:	$\alpha?x$ where $\alpha \in Chn$, $x \in Var$
Output operations:	$\alpha!a$ where $\alpha \in Chn$, $a \in AExp$
Guarded commands:	$gc \in GCmd$

Syntax of CSP

Definition 17.1 (Syntax of CSP)

The syntax of CSP is given by

$$
\begin{aligned}
a &::= z \mid x \mid a_1 + a_2 \mid a_1 - a_2 \mid a_1 * a_2 \in AExp \\
b &::= t \mid a_1 = a_2 \mid a_1 > a_2 \mid \neg b \mid b_1 \wedge b_2 \mid b_1 \vee b_2 \in BExp \\
c &::= \texttt{skip} \mid x := a \mid \alpha?x \mid \alpha!a \mid \\
&\qquad c_1 ; c_2 \mid \texttt{if}\ gc\ \texttt{fi} \mid \texttt{do}\ gc\ \texttt{od} \mid c_1 \parallel c_2 \in Cmd \\
gc &::= b \to c \mid b \wedge \alpha?x \to c \mid b \wedge \alpha!a \to c \mid gc_1 \,\square\, gc_2 \in GCmd
\end{aligned}
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gc &::= b \to c \mid b \wedge \alpha?x \to c \mid b \wedge \alpha!a \to c \mid gc_1 \,\Box\, gc_2 \in GCmd
\end{aligned}
$$

- In $c_1 \parallel c_2$, commands c_1 and c_2 must not use common variables (only local store)
- Guarded command $gc_1 \,\Box\, gc_2$ represents an alternative
- In $b \to c$, b acts as a guard that enables the execution of c only if evaluated to true
- $b \wedge \alpha?x \to c$ and $b \wedge \alpha!a \to c$ additionally require the respective I/O operation to be enabled
- If none of its alternatives is enabled, a guarded command gc fails (configuration fail)
- `if` nondeterministically picks an enabled alternative
- A `do` loop is iterated until its body fails

Semantics of CSP I

- Most important aspect: I/O operations
- E.g., $\langle \alpha?x; c, \sigma\rangle$ can only execute if a parallel command provides corresponding output

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Semantics of CSP I

- Most important aspect: I/O operations
- E.g., $\langle \alpha?x; c, \sigma\rangle$ can only execute if a parallel command provides corresponding output

$\Rightarrow$ Indicate communication potential by labels

$$L := \{\alpha?z \mid \alpha \in Chn, z \in \mathbb{Z}\} \cup \{\alpha!z \mid \alpha \in Chn, z \in \mathbb{Z}\}$$

- Yields following labelled transitions:

$$\langle \alpha?x; c_1, \sigma\rangle \xrightarrow{\alpha?z} \langle c_1, \sigma[x \mapsto z]\rangle \quad \text{(for all } z \in \mathbb{Z})$$
$$\langle \alpha!a; c_2, \sigma\rangle \xrightarrow{\alpha!z} \langle c_2, \sigma\rangle \quad \text{(if } \langle a, \sigma\rangle \to z)$$

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$$\langle \alpha?x; c_1, \sigma\rangle \xrightarrow{\alpha?z} \langle c_1, \sigma[x \mapsto z]\rangle \quad (\text{for all } z \in \mathbb{Z})$$

$$\langle \alpha!a; c_2, \sigma\rangle \xrightarrow{\alpha!z} \langle c_2, \sigma\rangle \quad (\text{if } \langle a, \sigma\rangle \to z)$$

- Now both commands, if running in parallel, can communicate:

$$\langle (\alpha?x; c_1) \parallel (\alpha!a; c_2), \sigma\rangle \to \langle c_1 \parallel c_2, \sigma[x \mapsto z]\rangle.$$

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- Now both commands, if running in parallel, can communicate:

$$\langle(\alpha?x; c_1) \parallel (\alpha!a; c_2), \sigma\rangle \to \langle c_1 \parallel c_2, \sigma[x \mapsto z]\rangle.$$

- To allow communication with other processes, the following transitions should also be enabled (for $\langle a, \sigma\rangle \to z$ and all $z' \in \mathbb{Z}$):

$$\langle(\alpha?x; c_1) \parallel (\alpha!a; c_2), \sigma\rangle \xrightarrow{\alpha?z'} \langle c_1 \parallel (\alpha!a; c_2), \sigma[x \mapsto z']\rangle$$
$$\langle(\alpha?x; c_1) \parallel (\alpha!a; c_2), \sigma\rangle \xrightarrow{\alpha!z} \langle(\alpha?x; c_1) \parallel c_2, \sigma\rangle$$

Semantics of CSP II

Definition of transition relation

$$\xrightarrow{\lambda} \subseteq (\mathit{Cmd} \times \Sigma) \times (\mathit{Cmd} \times \Sigma) \cup (\mathit{GCmd} \times \Sigma) \times (\mathit{Cmd} \times \Sigma \cup \{\mathrm{fail}\})$$

(see following slides)

- Marking λ can be a label or empty: $\lambda \in L \cup \{\varepsilon\}$
- Again: uniform treatment of configurations of the form $\langle c, \sigma \rangle \in \mathit{Cmd} \times \Sigma$ and $\sigma \in \Sigma$:
  - σ interpreted as $\langle \downarrow, \sigma \rangle$ with "terminated" command $\downarrow$
  - $\downarrow$ satisfies $\downarrow; c = \downarrow \parallel c = c \parallel \downarrow = c$

Semantics of CSP III

Definition 17.2 (Semantics of CSP – Commands (*Cmd*))

$$\frac{}{\langle \texttt{skip}, \sigma\rangle \rightarrow \langle \downarrow, \sigma\rangle} \qquad \frac{\langle a, \sigma\rangle \rightarrow z}{\langle x := a, \sigma\rangle \rightarrow \langle \downarrow, \sigma[x \mapsto z]\rangle}$$

$$\frac{}{\langle \alpha?x, \sigma\rangle \xrightarrow{\alpha?z} \langle \downarrow, \sigma[x \mapsto z]\rangle} \qquad \frac{\langle a, \sigma\rangle \rightarrow z}{\langle \alpha!a, \sigma\rangle \xrightarrow{\alpha!z} \langle \downarrow, \sigma\rangle}$$

$$\frac{\langle c_1, \sigma\rangle \xrightarrow{\lambda} \langle c_1', \sigma'\rangle}{\langle c_1; c_2, \sigma\rangle \xrightarrow{\lambda} \langle c_1'; c_2, \sigma'\rangle} \qquad \frac{\langle gc, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}{\langle \texttt{if}\ gc\ \texttt{fi}, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}$$

$$\frac{\langle gc, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}{\langle \texttt{do}\ gc\ \texttt{od}, \sigma\rangle \xrightarrow{\lambda} \langle c; \texttt{do}\ gc\ \texttt{od}, \sigma'\rangle} \qquad \frac{\langle gc, \sigma\rangle \rightarrow \text{fail}}{\langle \texttt{do}\ gc\ \texttt{od}, \sigma\rangle \rightarrow \langle \downarrow, \sigma\rangle}$$

$$\frac{\langle c_1, \sigma\rangle \xrightarrow{\lambda} \langle c_1', \sigma'\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \xrightarrow{\lambda} \langle c_1' \parallel c_2, \sigma'\rangle} \qquad \frac{\langle c_2, \sigma\rangle \xrightarrow{\lambda} \langle c_2', \sigma'\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \xrightarrow{\lambda} \langle c_1 \parallel c_2', \sigma'\rangle}$$

$$\frac{\langle c_1, \sigma\rangle \xrightarrow{\alpha?z} \langle c_1', \sigma'\rangle \quad \langle c_2, \sigma\rangle \xrightarrow{\alpha!z} \langle c_2', \sigma\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \rightarrow \langle c_1' \parallel c_2', \sigma'\rangle} \qquad \frac{\langle c_1, \sigma\rangle \xrightarrow{\alpha!z} \langle c_1', \sigma\rangle \quad \langle c_2, \sigma\rangle \xrightarrow{\alpha?z} \langle c_2', \sigma'\rangle}{\langle c_1 \parallel c_2, \sigma\rangle \rightarrow \langle c_1' \parallel c_2', \sigma'\rangle}$$

Semantics of CSP IV

Definition 17.2 (Semantics of CSP – Guarded commands (*GCmd*))

$$\frac{\langle b, \sigma\rangle \to \mathsf{true}}{\langle b \to c, \sigma\rangle \to \langle c, \sigma\rangle} \qquad \frac{\langle b, \sigma\rangle \to \mathsf{false}}{\langle b \to c, \sigma\rangle \to \mathsf{fail}}$$

$$\frac{\langle b, \sigma\rangle \to \mathsf{true}}{\langle b \wedge \alpha?x \to c, \sigma\rangle \xrightarrow{\alpha?z} \langle c, \sigma[x \mapsto z]\rangle} \qquad \frac{\langle b, \sigma\rangle \to \mathsf{false}}{\langle b \wedge \alpha?x \to c, \sigma\rangle \to \mathsf{fail}}$$

$$\frac{\langle b, \sigma\rangle \to \mathsf{true} \quad \langle a, \sigma\rangle \to z}{\langle b \wedge \alpha!a \to c, \sigma\rangle \xrightarrow{\alpha!z} \langle c, \sigma\rangle} \qquad \frac{\langle b, \sigma\rangle \to \mathsf{false}}{\langle b \wedge \alpha!a \to c, \sigma\rangle \to \mathsf{fail}}$$

$$\frac{\langle gc_1, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}{\langle gc_1 \,\square\, gc_2, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle} \qquad \frac{\langle gc_2, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}{\langle gc_1 \,\square\, gc_2, \sigma\rangle \xrightarrow{\lambda} \langle c, \sigma'\rangle}$$

$$\frac{\langle gc_1, \sigma\rangle \to \mathsf{fail} \quad \langle gc_2, \sigma\rangle \to \mathsf{fail}}{\langle gc_1 \,\square\, gc_2, \sigma\rangle \to \mathsf{fail}}$$

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CSP Examples

Example 17.3

(on the board)

1. $\texttt{do}\ (\text{true} \wedge \alpha?x \rightarrow \beta!x)\ \texttt{od}$
   describes a process that repeatedly receives a value along α and forwards it along β (i.e., a one-place buffer)

CSP Examples

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1. $\texttt{do}\ (\text{true} \wedge \alpha?x \rightarrow \beta!x)\ \texttt{od}$
   describes a process that repeatedly receives a value along α and forwards it along β (i.e., a one-place buffer)
2. $\texttt{do}\ \text{true} \wedge \alpha?x \rightarrow \beta!x\ \texttt{od} \parallel \texttt{do}\ \text{true} \wedge \beta?y \rightarrow \gamma!y\ \texttt{od}$
   specifies a two-place buffer that receives along α and sends along γ (using β for internal communication)

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Example 17.3

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   specifies a two-place buffer that receives along α and sends along γ (using β for internal communication)
3. Nondeterministic choice between input channels:
   i. $\texttt{if}\ (\text{true} \wedge \alpha?x \rightarrow c_1 \,\square\, \text{true} \wedge \beta?y \rightarrow c_2)\ \texttt{fi}$
   ii. $\texttt{if}\ (\text{true} \rightarrow (\alpha?x; c_1) \,\square\, \text{true} \rightarrow (\beta?y; c_2))\ \texttt{fi}$

   Expected: progress whenever environment provides data on α or β
   i. correct
   ii. incorrect (can deadlock)

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Summary: Nondeterminism and Parallelism

Fairness I

- Informally: unfair behaviour excludes processes from being executed
- Here: consider parallel composition of $n \geq 1$ sequential programs with executions of the form $\kappa_0 \to \kappa_1 \to \kappa_2 \to \ldots$ where $\kappa_j = \langle c_1^{(j)} \parallel \ldots \parallel c_n^{(j)}, \sigma_j \rangle$ and, for some $1 \leq i \leq n$ and $k_0 \in \mathbb{N}$, $c_i^{(k)} = c_i^{(k_0)}$ for all $k \geq k_0$
- But: only unfair if c_i not enabled

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- But: only unfair if c_i not enabled

Definition 17.4 (Enabledness)

c_i is enabled in configuration $\kappa = \langle c_1 \parallel \ldots \parallel c_n, \sigma \rangle$ if there exists $\kappa' = \langle c_1' \parallel \ldots \parallel c_n', \sigma' \rangle$ with $\kappa \to \kappa'$ and $c_i' \neq c_i$.

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- Here: consider parallel composition of $n \geq 1$ sequential programs with executions of the form $\kappa_0 \to \kappa_1 \to \kappa_2 \to \ldots$ where $\kappa_j = \langle c_1^{(j)} \parallel \ldots \parallel c_n^{(j)}, \sigma_j \rangle$ and, for some $1 \leq i \leq n$ and $k_0 \in \mathbb{N}$, $c_i^{(k)} = c_i^{(k_0)}$ for all $k \geq k_0$
- But: only unfair if c_i not enabled

Definition 17.4 (Enabledness)

c_i is enabled in configuration $\kappa = \langle c_1 \parallel \ldots \parallel c_n, \sigma \rangle$ if there exists $\kappa' = \langle c_1' \parallel \ldots \parallel c_n', \sigma' \rangle$ with $\kappa \to \kappa'$ and $c_i' \neq c_i$.

Example 17.5

1. $x := 0$ enabled in $\langle x := 0 \parallel y := 1, \sigma \rangle$ (actually always enabled)
2. $\alpha?x$ enabled in $\langle \alpha?x \parallel \alpha!0, \sigma \rangle$
3. $\alpha?x$ not enabled in $\langle \alpha?x \parallel \beta!1, \sigma \rangle$

Fairness II

Definition 17.6 (Fairness)

An execution $\kappa_0 \to \kappa_1 \to \kappa_2 \to \ldots$ where $\kappa_j = \langle c_1^{(j)} \parallel \ldots \parallel c_n^{(j)}, \sigma_j \rangle$ and, for some $1 \leq i \leq n$ and $k_0 \in \mathbb{N}$, $c_i^{(k)} = c_i^{(k_0)}$ for all $k \geq k_0$ is called

1. strongly unfair if $c_i^{(k)}$ is enabled in κ_k for all $k \geq k_0$
2. weakly unfair if $c_i^{(k)}$ is enabled in κ_k for infinitely many $k \geq k_0$

Fairness III

Example 17.7

1. $$\begin{aligned} & \langle \mathtt{do}\ \mathsf{true} \to x := x + 1\ \mathtt{od} \parallel y := y + 1, \ldots \rangle \\ \to\ & \langle x := x + 1; \mathtt{do}\ \mathsf{true} \to x := x + 1\ \mathtt{od} \parallel y := y + 1, \ldots \rangle \\ \to\ & \langle \mathtt{do}\ \mathsf{true} \to x := x + 1\ \mathtt{od} \parallel y := y + 1, \ldots \rangle \to \ldots \end{aligned}$$

   is strongly unfair since $y := y + 1$ is always enabled

Fairness III

Example 17.7

1. $\quad\langle \text{do true} \to x := x+1 \text{ od} \parallel y := y+1, \ldots\rangle$
   $\to \langle x := x+1; \text{do true} \to x := x+1 \text{ od} \parallel y := y+1, \ldots\rangle$
   $\to \langle \text{do true} \to x := x+1 \text{ od} \parallel y := y+1, \ldots\rangle \to \ldots$

   is strongly unfair since $y := y+1$ is always enabled
2. $\quad\langle \text{do true} \to x := x+1 \text{ od} \parallel \alpha!1 \parallel \alpha?y, \ldots\rangle$
   $\to \langle x := x+1; \text{do true} \to x := x+1 \text{ od} \parallel \alpha!1 \parallel \alpha?y, \ldots\rangle$
   $\to \langle \text{do true} \to x := x+1 \text{ od} \parallel \alpha!1 \parallel \alpha?y, \ldots\rangle \to \ldots$

   is strongly unfair since both I/O operations are always enabled

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   is strongly unfair since $y := y + 1$ is always enabled
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   is strongly unfair since both I/O operations are always enabled
3. $$\begin{array}{ll} & \langle \texttt{do } \alpha!1 \rightarrow \texttt{skip od} \parallel \texttt{do } \alpha?x \rightarrow \texttt{skip od} \parallel \alpha?y, \ldots \rangle \\ \rightarrow & \langle \texttt{skip}; \texttt{do } \alpha!1 \rightarrow \texttt{skip od} \parallel \texttt{skip}; \texttt{do } \alpha?x \rightarrow \texttt{skip od} \parallel \alpha?y, \ldots \rangle \\ \rightarrow & \langle \texttt{skip}; \texttt{do } \alpha!1 \rightarrow \texttt{skip od} \parallel \texttt{do } \alpha?x \rightarrow \texttt{skip od} \parallel \alpha?y, \ldots \rangle \\ \rightarrow & \langle \texttt{do } \alpha!1 \rightarrow \texttt{skip od} \parallel \texttt{do } \alpha?x \rightarrow \texttt{skip od} \parallel \alpha?y, \ldots \rangle \rightarrow \ldots \end{array}$$
   is weakly unfair since $\alpha?y$ is (only) enabled in every third configuration

Summary: Nondeterminism and Parallelism

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- Important modelling aspects:
  - parallelism (here: interleaving = nondeterminism + sequential execution)
  - interaction (here: via shared variables/channels)
- Interleaving requires small-step execution relation
- Communication between parallel processes is represented by labels on transitions
- Parallelism raises new issues such as fairness