Concurrency Theory

Lecture 17: Interleaving Semantics of Petri Nets

Joost-Pieter Katoen and Thomas Noll

Lehrstuhl für Informatik 2
Software Modeling and Verification Group

http://moves.rwth-aachen.de/teaching/ws-1718/ct

December 11, 2017

Overview

Overview

Carl Adam Petri (1926-2010)

The original work[1] does not contain a single (graphical) Petri net!

[1] Petri's PhD dissertation, 1962.

Semantics: executions and traces

Semantics: executions and traces

Models in the 60s: lambda calculus, finite automata, Turing machines, . . .

Semantics: executions and traces

Models in the 60s: lambda calculus, finite automata, Turing machines, ...

States: current configurations of the machine

- One or more initial states
- Possibly some distinguished final states

Transitions: moves between configurations

Lambda calculus	$(\lambda x.xx)(\lambda y.y)$	$\longrightarrow$	$(\lambda y.y)(\lambda z.z)$
Turing machine	$0010q_1011$	$\longrightarrow$	$001q_201011$
Finite automaton	q_1	$\xrightarrow{a}$	q_2
Pushdown automaton	$(q_1, XYYZ)$	$\xrightarrow{a}$	$(q_2, XYXYYZ)$

Executions: alternating sequences of states and transitions

Petri's question

Petri's question

C.A. Petri points out a discrepancy between how Theoretical Physics and Theoretical Computer Science described systems in 1962:

Theoretical Physics describes systems as a collection of interacting particles (subsystems), without a notion of global clock or simultaneity

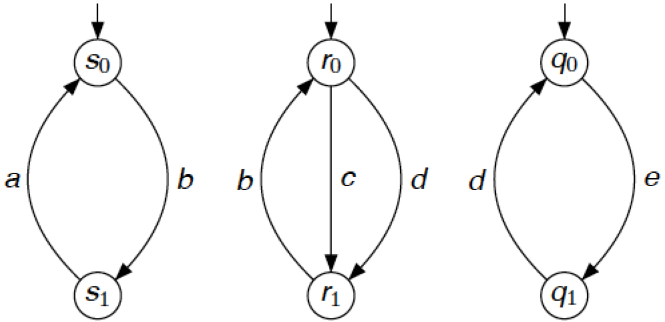
Theoretical Computer Science describes systems as sequential virtual machines going through a temporally ordered sequence of global states

Petri's question:

Which kind of abstract machine should be used to describe the physical implementation of a Turing machine?

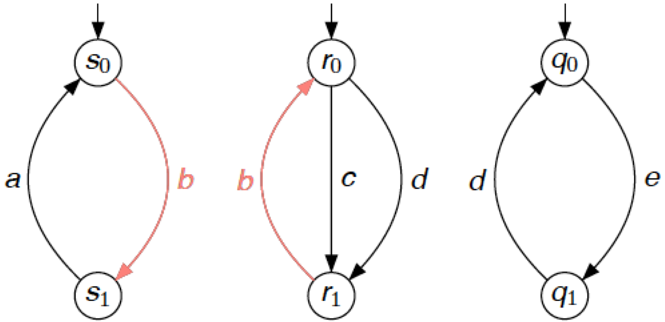
Petri net

A graphical representation of interacting finite automata:



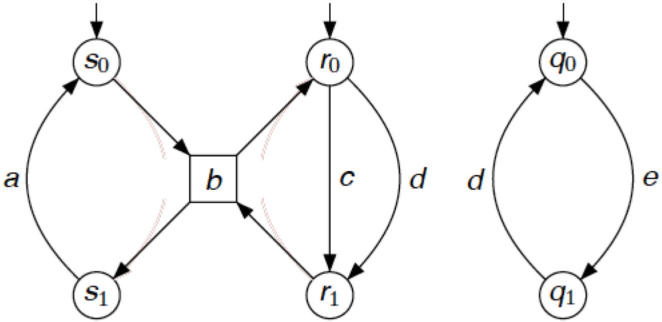
Petri net

A graphical representation of interacting finite automata:



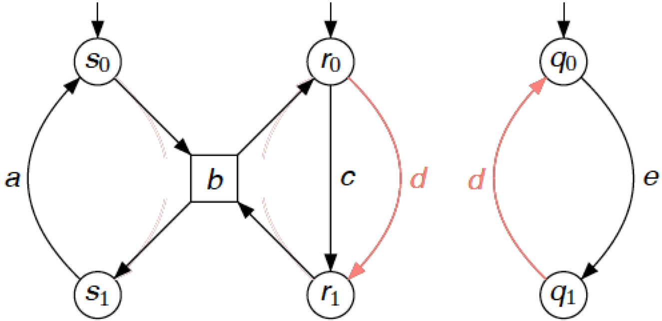
Petri net

A graphical representation of interacting finite automata:



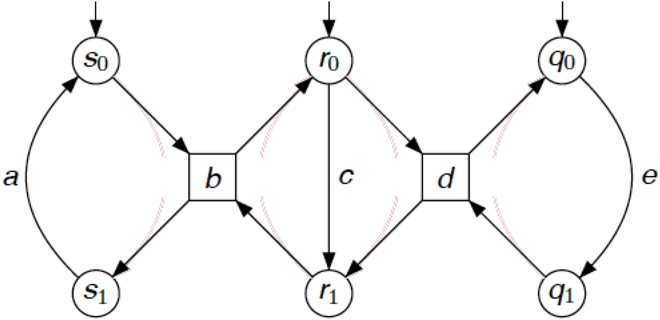
Petri net

A graphical representation of interacting finite automata:



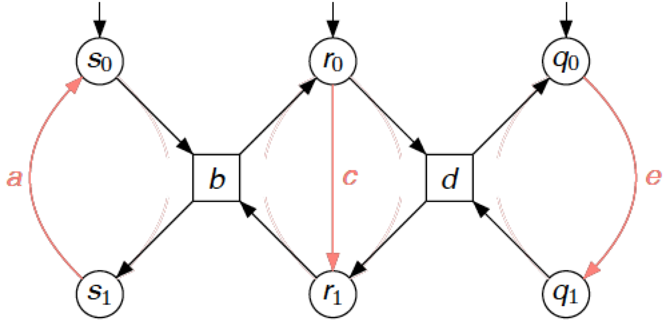
Petri net

A graphical representation of interacting finite automata:



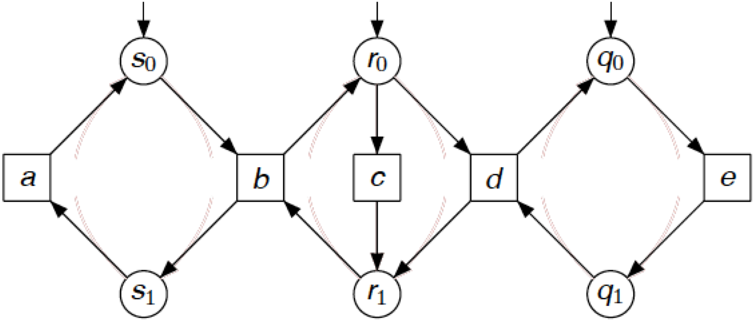
Petri net

A graphical representation of interacting finite automata:



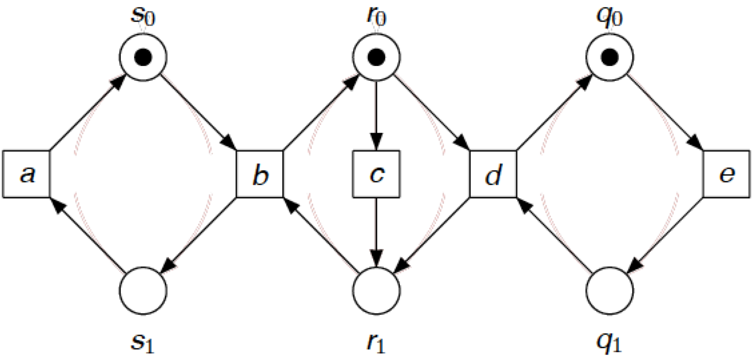
Petri net

A graphical representation of interacting finite automata:



Petri net

A graphical representation of interacting finite automata:



Overview

Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.

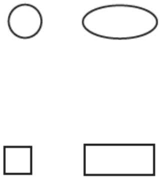
Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.

A place is represented by a circle or ellipse. A place p always models a passive component: p can store, accumulate or show things.

Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.



A place is represented by a circle or ellipse. A place p always models a passive component: p can store, accumulate or show things.

A transition is represented by a square or rectangle. A transition t always models an active component: t can produce things, consume, transport or change them.

Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.

A place is represented by a circle or ellipse. A place p always models a passive component: p can store, accumulate or show things.

A transition is represented by a square or rectangle. A transition t always models an active component: t can produce things, consume, transport or change them.

Places and transitions are connected to each other by directed arcs. Graphically, an arc is represented by an arrow.

Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.

A place is represented by a circle or ellipse. A place p always models a passive component: p can store, accumulate or show things.

A transition is represented by a square or rectangle. A transition t always models an active component: t can produce things, consume, transport or change them.

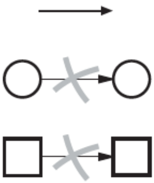
Places and transitions are connected to each other by directed arcs. Graphically, an arc is represented by an arrow. An arc models an abstract, sometimes only notional relation between components.

Components of a net

A Petri net is a structure with two kinds of elements: places and transitions. They are connected by arcs.

A place is represented by a circle or ellipse. A place *p* always models a passive component: *p* can store, accumulate or show things.

A transition is represented by a square or rectangle. A transition *t* always models an active component: *t* can produce things, consume, transport or change them.



Places and transitions are connected to each other by directed arcs. Graphically, an arc is represented by an arrow. An arc models an abstract, sometimes only notional relation between components. Arcs run from places to transitions or vice versa.

Nets

[2] F is also called the flow relation.

Nets

Net

A Petri net N is a triple (P, T, F) where:

- P is the finite set of places
- T is the finite set of transitions with $P \cap T = \varnothing$
- $F \subseteq (P \times T) \cup (T \times P)$ are the arcs[2]

Places and transitions are generically called nodes.

[2] F is also called the flow relation.

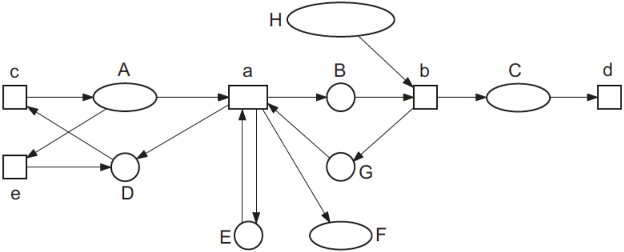
Nets

Net

A Petri net N is a triple (P, T, F) where:

- P is the finite set of places
- T is the finite set of transitions with $P \cap T = \varnothing$
- $F \subseteq (P \times T) \cup (T \times P)$ are the arcs[2]

Places and transitions are generically called nodes.



[2] F is also called the flow relation.

The pre- and post-sets

The pre- and post-sets

Pre- and post-sets

Let node $x \in P \cup T$.
The pre-set of x is defined by: ${}^\bullet x = \{ y \mid (y, x) \in F \}$.
The post-set of x is defined by: $x^\bullet = \{ y \mid (x, y) \in F \}$.

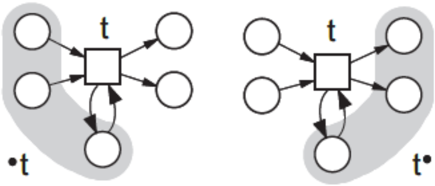
Two nodes $x, y \in N$ form a loop if $x \in {}^\bullet y$ and $y \in {}^\bullet x$.

The pre- and post-sets

Pre- and post-sets

Let node $x \in P \cup T$.
The pre-set of x is defined by: ${}^\bullet x = \{ y \mid (y, x) \in F \}$.
The post-set of x is defined by: $x^\bullet = \{ y \mid (x, y) \in F \}$.

Two nodes $x, y \in N$ form a loop if $x \in {}^\bullet y$ and $y \in {}^\bullet x$.



Markings

Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \to \mathbb{N}$.

Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \to \mathbb{N}$.

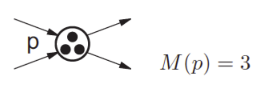
For net $N = (P, T, F)$ and marking M_0, the tuple (P, T, F, M_0) is called an elementary system net. M_0 is the initial marking of N.

Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \to \mathbb{N}$.

For net $N = (P, T, F)$ and marking M_0, the tuple (P, T, F, M_0) is called an elementary system net. M_0 is the initial marking of N.

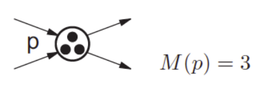


Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \rightarrow \mathbb{N}$.

For net $N = (P, T, F)$ and marking M_0, the tuple (P, T, F, M_0) is called an elementary system net. M_0 is the initial marking of N.



Intuition

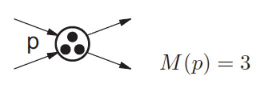
Note: a marking is a multiset.

Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \to \mathbb{N}$.

For net $N = (P, T, F)$ and marking M_0, the tuple (P, T, F, M_0) is called an elementary system net. M_0 is the initial marking of N.



Intuition

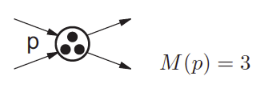
Note: a marking is a multiset. It defines a distribution of tokens across places.

Markings

Marking

A marking M of a net $N = (P, T, F)$ is a mapping $M : P \to \mathbb{N}$.

For net $N = (P, T, F)$ and marking M_0, the tuple (P, T, F, M_0) is called an elementary system net. M_0 is the initial marking of N.



Intuition

Note: a marking is a multiset. It defines a distribution of tokens across places. Tokens are depicted as black dots.

Transition firing

Transition firing

Enabling and occurrence of a transition

Let (P, T, F, M) be an elementary system net. Marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in {}^\bullet t$.

Transition firing

Enabling and occurrence of a transition

Let (P, T, F, M) be an elementary system net. Marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in^\bullet t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \ = \ M(p) - F(p, t) + F(t, p).$$

where we represent F by its characteristic function.

Transition firing

Enabling and occurrence of a transition

Let (P, T, F, M) be an elementary system net. Marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in {}^{\bullet}t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \; = \; M(p) - F(p, t) + F(t, p).$$

where we represent F by its characteristic function.

Intuition

Transition firing

Enabling and occurrence of a transition

Let (P, T, F, M) be an elementary system net. Marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in^{\bullet} t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \ = \ M(p) - F(p, t) + F(t, p).$$

where we represent F by its characteristic function.

Intuition

Transition t is enabled whenever every $p \in^{\bullet} t$ holds at least one token.

Transition firing

Enabling and occurrence of a transition

Let (P, T, F, M) be an elementary system net. Marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in {}^{\bullet}t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \;=\; M(p) - F(p,t) + F(t,p).$$

where we represent F by its characteristic function.

Intuition

Transition t is enabled whenever every $p \in {}^{\bullet}t$ holds at least one token. On t's occurrence, one token is removed from each place in ${}^{\bullet}t$, and one token is put in each place in $t^{\bullet}$:

$$M'(p) \;=\; \begin{cases} M(p) - 1 & \text{if } p \in {}^{\bullet}t \text{ and } p \notin t^{\bullet} \\ M(p) + 1 & \text{if } p \in t^{\bullet} \text{ and } p \notin {}^{\bullet}t \\ M(p) & \text{otherwise} \end{cases}$$

Transition occurrence

Transition occurrence

Enabling and occurrence of a transition

A marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in^{\bullet} t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \;=\; M(p) - F(p, t) + F(t, p).$$

where we represent F by its characteristic function.

Transition occurrence

Enabling and occurrence of a transition

A marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in^{\bullet} t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \;=\; M(p) - F(p,t) + F(t,p).$$

where we represent F by its characteristic function.
$M \xrightarrow{t} M'$ is also called a step of the net N.

Transition occurrence

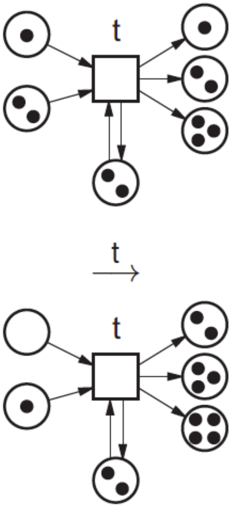
Enabling and occurrence of a transition

A marking M enables a transition t if $M(p) \geqslant 1$ for each place $p \in^{\bullet} t$.

Transition t can occur in marking M if t is enabled at M. Its occurrence leads to marking M', denoted $M \xrightarrow{t} M'$, defined for place $p \in P$ by:

$$M'(p) \;=\; M(p) - F(p,t) + F(t,p).$$

where we represent F by its characteristic function.
$M \xrightarrow{t} M'$ is also called a step of the net N.



Overview

The interleaving semantics of Petri nets

The interleaving semantics of Petri nets

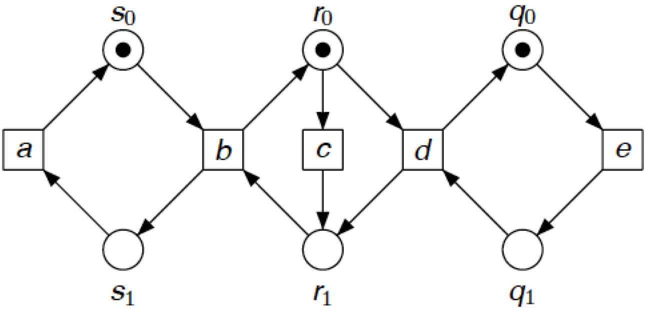
An execution semantics

State: marking (distribution of tokens over the net)

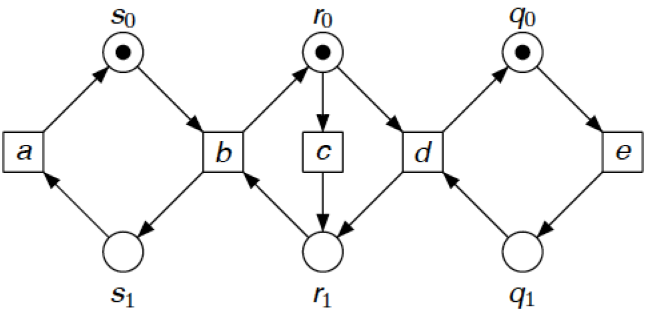
Transitions: $M \xrightarrow{t} M'$

Sequential runs: $M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} M_2 \xrightarrow{t_3} \ldots$

The interleaving semantics of Petri nets



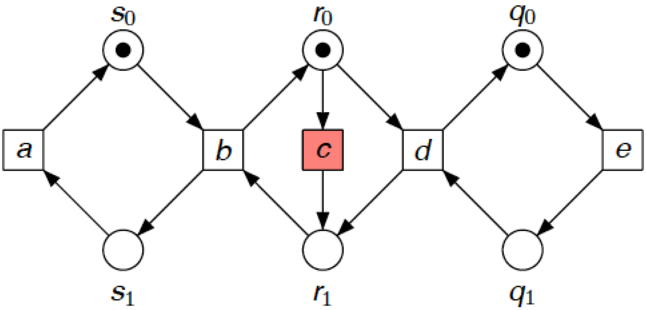
The interleaving semantics of Petri nets



$$\begin{array}{c} s_1 \\ r_1 \\ q_1 \end{array} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

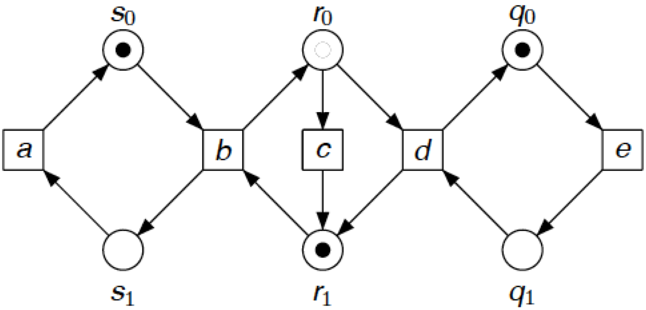
As the marking for s_0 is the complement of s_1, the marking for s_0 is omitted. The same applies to the places r_0 and q_0.

The interleaving semantics of Petri nets



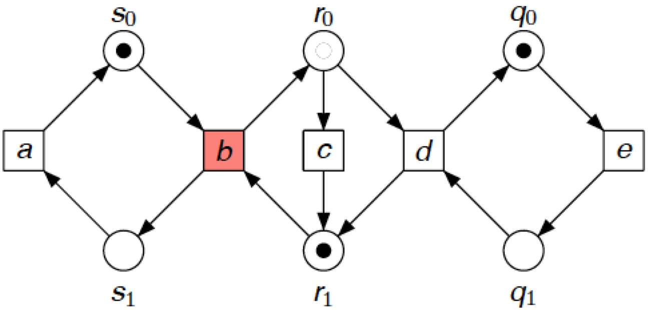
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c}$$

The interleaving semantics of Petri nets



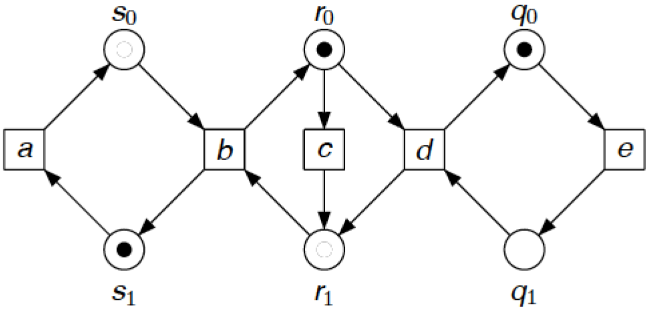
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$$

The interleaving semantics of Petri nets



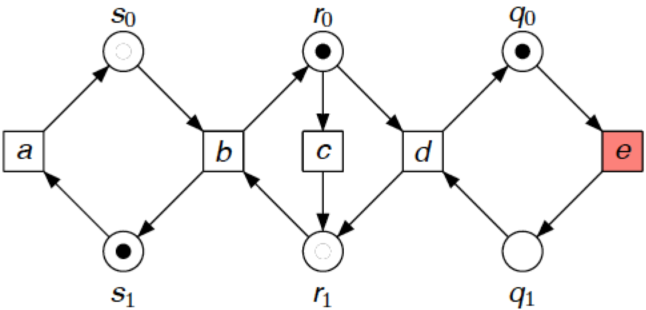
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b}$$

The interleaving semantics of Petri nets



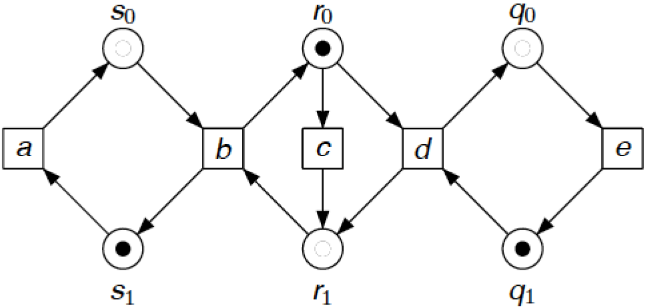
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}$$

The interleaving semantics of Petri nets



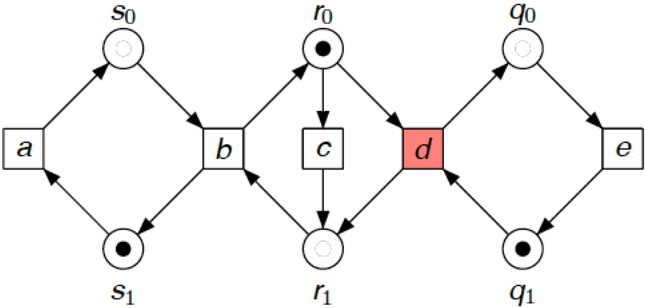
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{\color{red}{e}}$$

The interleaving semantics of Petri nets



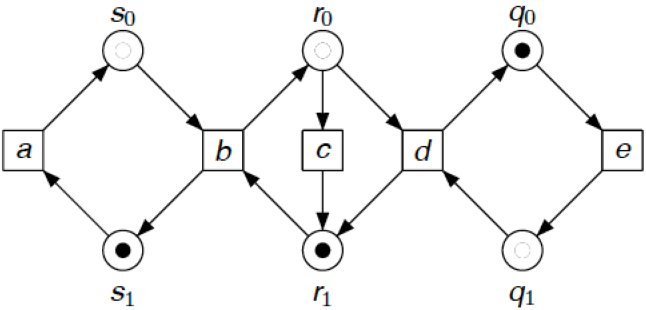
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}$$

The interleaving semantics of Petri nets



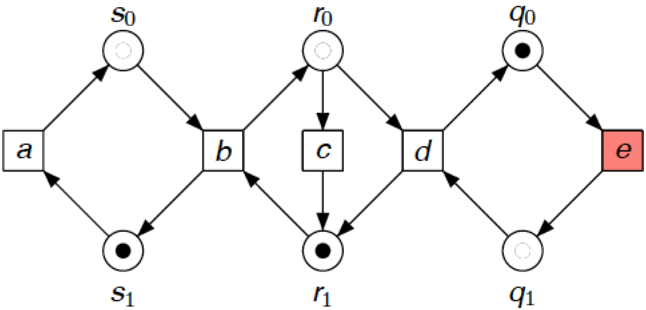
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \xrightarrow{d}$$

The interleaving semantics of Petri nets



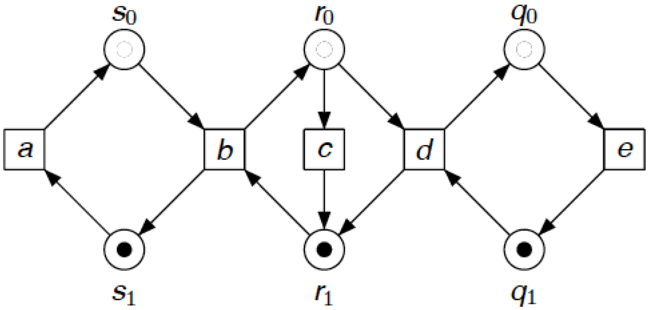
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \xrightarrow{d} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}$$

The interleaving semantics of Petri nets



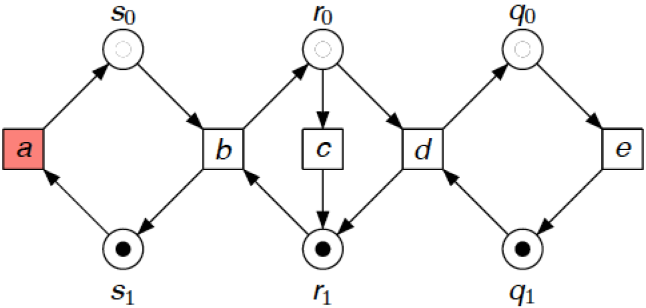
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \xrightarrow{d} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{e}$$

The interleaving semantics of Petri nets



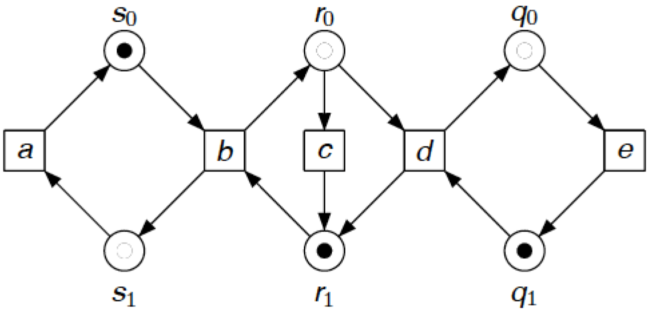
$$\begin{array}{c} s_1 \\ r_1 \\ q_1 \end{array} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \xrightarrow{d} \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{e} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$$

The interleaving semantics of Petri nets



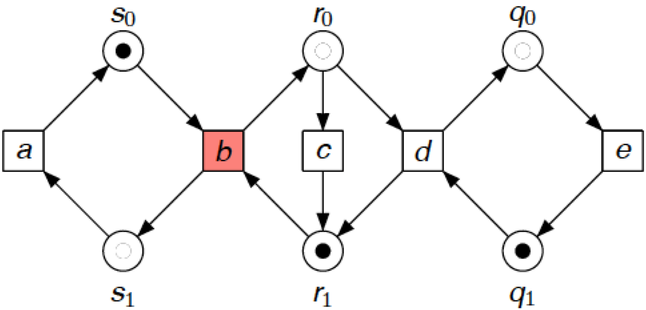
$$
\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix}
\begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c}
\begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b}
\begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{e}
\begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} \xrightarrow{d}
\begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{e}
\begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \xrightarrow{a}
$$

The interleaving semantics of Petri nets



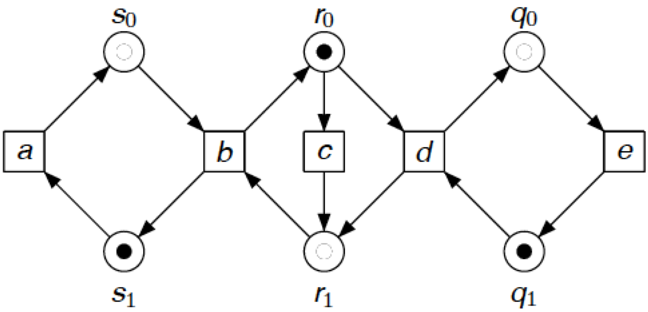
$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \cdots \xrightarrow{e} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \xrightarrow{a} \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}$$

The interleaving semantics of Petri nets



$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \cdots \xrightarrow{e} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \xrightarrow{a} \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix} \xrightarrow{b}$$

The interleaving semantics of Petri nets



$$\begin{matrix} s_1 \\ r_1 \\ q_1 \end{matrix} \quad \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \xrightarrow{c} \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \cdots \xrightarrow{a} \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix} \xrightarrow{b} \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix}$$

Reachable markings

Reachable markings

Step sequence

A sequence of transitions $\sigma = t_1\, t_2 \ldots t_n$ is an step sequence if there exist markings M_1 through M_n such that:

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \ldots \xrightarrow{t_{n-1}} M_{n-1} \xrightarrow{t_n} M_n$$

Reachable markings

Step sequence

A sequence of transitions $\sigma = t_1\, t_2 \ldots t_n$ is an step sequence if there exist markings M_1 through M_n such that:

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \cdots \xrightarrow{t_{n-1}} M_{n-1} \xrightarrow{t_n} M_n$$

Marking M_n is reached by the occurrence of σ, denoted $M_0 \xrightarrow{\sigma} M_n$.

Reachable markings

Step sequence

A sequence of transitions $\sigma = t_1\, t_2 \ldots t_n$ is an step sequence if there exist markings M_1 through M_n such that:

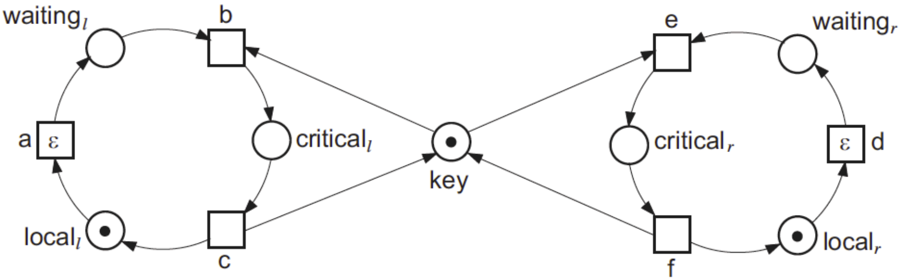
$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \cdots \xrightarrow{t_{n-1}} M_{n-1} \xrightarrow{t_n} M_n$$

Marking M_n is reached by the occurrence of σ, denoted $M_0 \xrightarrow{\sigma} M_n$.

M is a reachable marking if there exists a step sequence σ with $M_0 \xrightarrow{\sigma} M$.

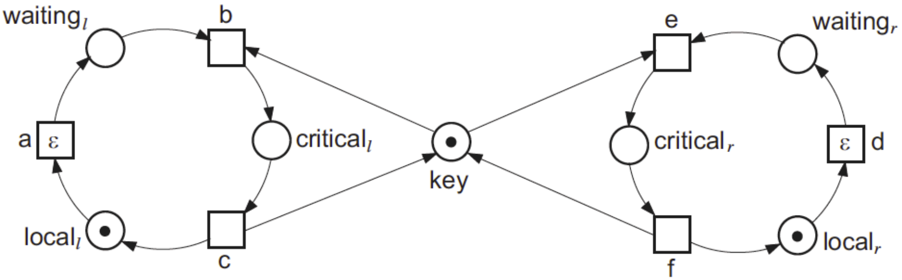
Mutual exclusion

Two processes cycling through the states local, waiting and critical.



Mutual exclusion

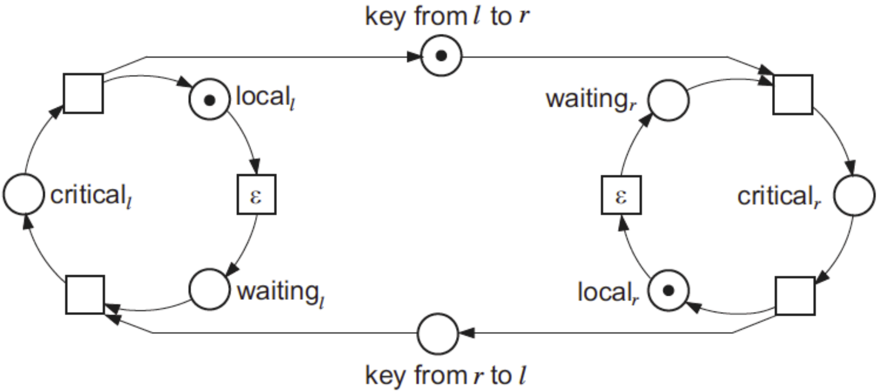
Two processes cycling through the states local, waiting and critical.



Between transitions b and e a conflict can arise infinitely often. No strategy has been modeled to solve this conflict.

Mutual exclusion

A strategy where processes are acquired access in an alternating fashion:



One-bounded elementary system nets

1-bounded elementary net system

An elementary net system N is called 1-bounded if for each reachable marking M and place p of N:

$$M(p) \leqslant 1.$$

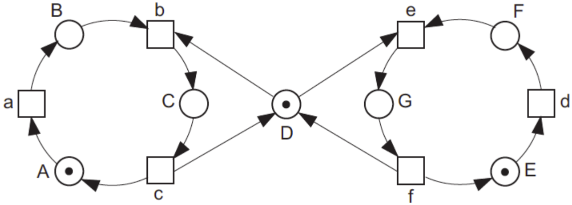
One-bounded elementary system nets

1-bounded elementary net system

An elementary net system N is called 1-bounded if for each reachable marking M and place p of N:

$$M(p) \leqslant 1.$$

Markings of 1-bounded elementary net systems can be described as a string of marked places, e.g., ADE. Two steps begin with this marking: $ADE \xrightarrow{a} BDE$ and $ADE \xrightarrow{d} ADF$.



Overview

Sequential runs

Sequential run

Let N be an elementary net system. A sequential run of N is a sequence

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \cdots$$

of steps of N starting with the initial marking M_0. A run can be finite or infinite. A finite run $M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_1} \cdot \xrightarrow{t_n} M_n$ is complete if M_n does not enable any transition.

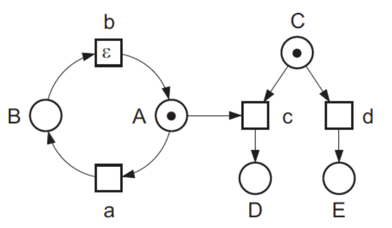
Sequential runs

Sequential run

Let N be an elementary net system. A sequential run of N is a sequence

$$M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_2} \cdots$$

of steps of N starting with the initial marking M_0. A run can be finite or infinite. A finite run $M_0 \xrightarrow{t_1} M_1 \xrightarrow{t_1} \cdot \xrightarrow{t_n} M_n$ is complete if M_n does not enable any transition.



A sample complete run is:

$$AC \xrightarrow{a} BC \xrightarrow{b} AC \xrightarrow{c} D$$

A sample incomplete run is:

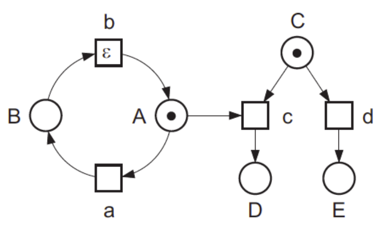
$$AC \xrightarrow{d} AE \xrightarrow{a} BE$$

Marking graph

The marking graph of N has as nodes the reachable markings of N and as edges the reachable steps of N.

Marking graph

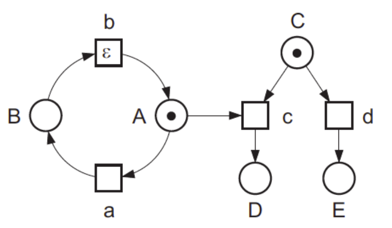
The marking graph of N has as nodes the reachable markings of N and as edges the reachable steps of N.



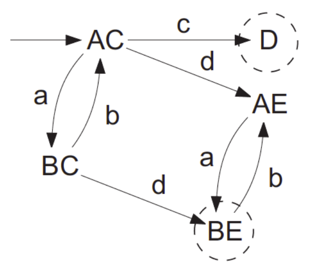
A sample elementary net system

Marking graph

The marking graph of N has as nodes the reachable markings of N and as edges the reachable steps of N.



A sample elementary net system



Its marking graph

Overview

Summary

- A Petri net consists of places, transitions and arcs

Summary

- A Petri net consists of places, transitions and arcs
- An elementary net is a Petri net plus a marking

Summary

- A Petri net consists of places, transitions and arcs
- An elementary net is a Petri net plus a marking
- Firing a single transition in a marking is a step

Summary

- A Petri net consists of places, transitions and arcs
- An elementary net is a Petri net plus a marking
- Firing a single transition in a marking is a step
- A sequential run is a sequence of steps starting in the initial marking

Summary

- A Petri net consists of places, transitions and arcs
- An elementary net is a Petri net plus a marking
- Firing a single transition in a marking is a step
- A sequential run is a sequence of steps starting in the initial marking
- A marking graph has as nodes the reachable markings of the net and as edges its reachable steps.

Summary

- A Petri net consists of places, transitions and arcs
- An elementary net is a Petri net plus a marking
- Firing a single transition in a marking is a step
- A sequential run is a sequence of steps starting in the initial marking
- A marking graph has as nodes the reachable markings of the net and as edges its reachable steps.
- The marking graph is the interleaving semantics of a net.