



General Remarks

- Please hand in your solutions in groups of 3. Either hand in your solutions at the beginning of the exercise class or put them into the box at the chair.
- If you have questions regarding the exercises and/or lecture, feel free to write us an email or visit us at our office.

Exercise 1 (Worklist Algorithm):

(2 Points)

Perform a live variable analysis on the following program using the worklist algorithm.

```
skip;  
while  $\neg(x > 3 * y)$  do  $x := x + y$  end  
 $y := x$ ;
```

Exercise 2 (Complexity of LVA Fixpoint Iteration):

(3 Points)

In the lecture we saw that fixpoint iteration requires at most $m \cdot n$ steps, where m is the height of the PO while n is the number of program points (i.e. labels). But how fast is the iterative algorithm for a concrete analysis (here: live variables)?

- Show that LVA has the following property:
Let $c \in \text{Cmd}$, $x \in \text{Var}_c$ and $l \in L_c$. If x is live on the exit of l , then there exists an **acyclic** path from B^l to a use of x that does not re-define x .
- Show that (standard) fixpoint iteration as seen in the lecture requires at most $|L_c|$ steps for convergence in case of LVA.

Exercise 3 (Definedness Analysis):

(3 Points)

- Formally describe how the constant propagation analysis can be turned into a “definedness” analysis that determines for every block which of the variables have a defined value (i.e. were initialised with an expression containing constants and initialised variables only). Assume that initially all variables have an undefined value.
- Prove or disprove: The MOP solution for the resulting analysis is decidable.

Exercise 4 (Distributivity of Transfer Functions):

(1 Points)

Consider a partial order of the form (D, \subseteq) , where $D = 2^M$ and M a finite set. A transfer function $\varphi_l : D \mapsto D$ is called distributive iff for any $d_1, d_2 \in D$:

$$\varphi_l(d_1 \sqcup d_2) = \varphi_l(d_1) \sqcup \varphi_l(d_2)$$

Show that any transfer function of the form $\varphi_l(d) = (d \setminus \text{kill}(B^l)) \cup \text{gen}(B^l)$ is distributive in this setting.