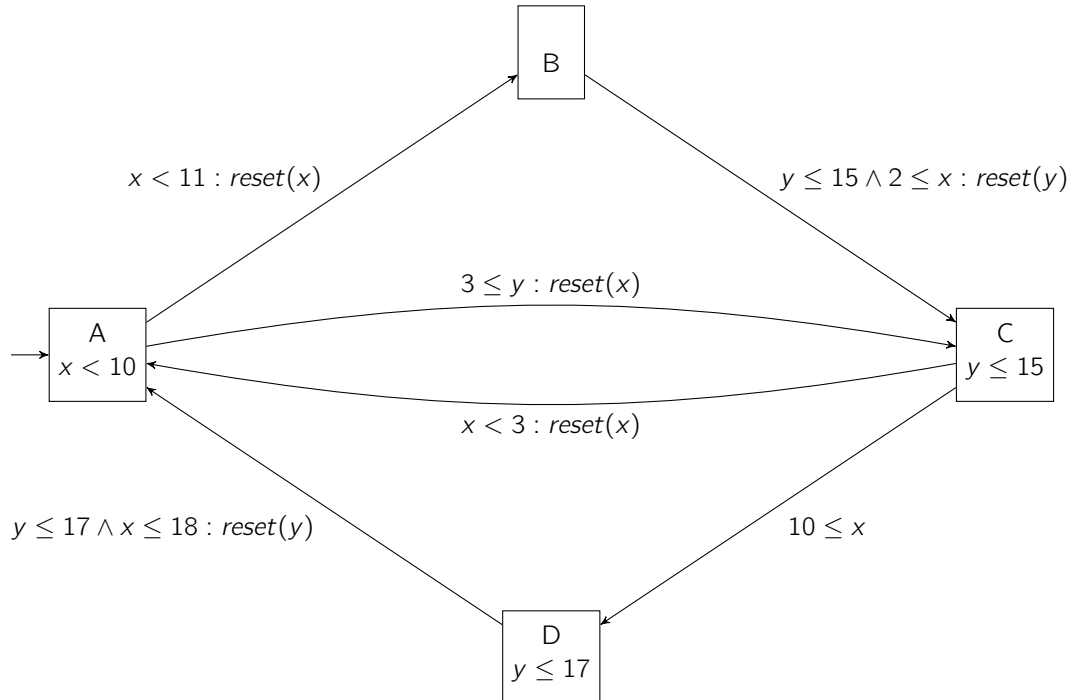


**Exercise 1 (Timelocks, zenoness):**

**(1+1 points)**

Consider the following TA and answer the following questions. Justify your answers.

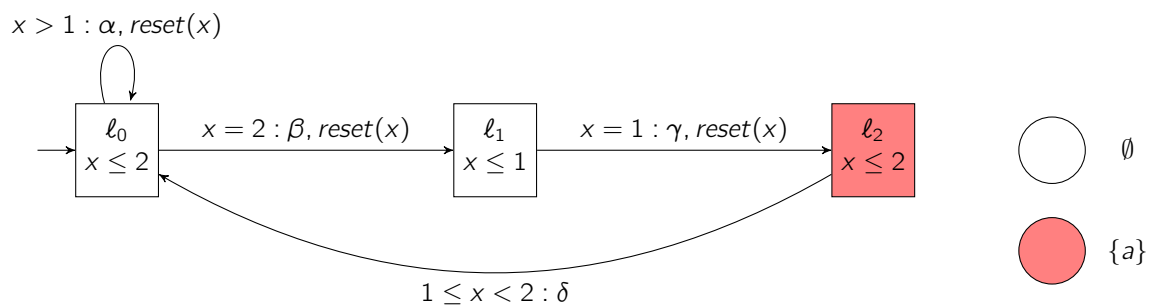


- (a) Is the TA non-zeno?
- (b) Is the TA timelock-free?

**Exercise 2 (Analysis of timed automata):**

**(1+2+1 points)**

Consider the following TA.



(a) TA

(b) Legend

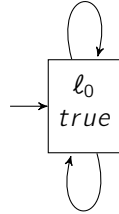
- (a) Determine the set of states  $Sat(\exists \diamond^{\leq 4} a)$ .
- (b) Determine the region transition system  $RTS(TA, true)$ .
- (c) Is the TA timelock-free? Justify your answer.

**Exercise 3 (Reachable states):**

**(2 points)**

Consider the following two timed automata  $TA_1$  and  $TA_2$ .

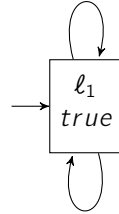
$y \leq 4 : \text{reset}(y)$



$x \leq 4 : \text{reset}(x)$

(c)  $TA_1$

$y \leq 4 : \text{reset}(x)$



$x \leq 4 : \text{reset}(y)$

(d)  $TA_2$

As these automata only have a single location, their states can be thought of as a point in the real plane. A point  $(d, e) \in \mathbb{R}_{\geq 0}^2$  then represents that clock  $x$  has value  $d$  and clock  $y$  has value  $e$ . Determine the reachable state space of both timed automata. Justify your answer.

**Exercise 4 (TCTL):**

**(1+1 points)**

(a) Find a counterexample showing that

$$\forall \diamond^{=d} \phi \wedge \forall \diamond^{=d} \psi \not\equiv \forall \diamond^{=d} (\phi \wedge \psi)$$

where  $\diamond^{=d} = \diamond^{[d,d]}$  for  $d \in \mathbb{R}_{\geq 0}$ .

(b) Does this also hold for  $J = [0, \infty)$ ? Justify your answer.