

Static Program Analysis

Lecture 9: Dataflow Analysis VIII

(Conditional Interval Analysis & Java Virtual Machine)

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(Software Modeling and Verification)



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<http://moves.rwth-aachen.de/teaching/ws-1415/spa/>

Winter Semester 2014/15

- 1 Recap: Taking Conditional Branches into Account
- 2 Interval Analysis with Assertions
- 3 The Java Virtual Machine
- 4 The Java Bytecode Verifier
- 5 The Type-Level Abstract Interpreter

- **Solution:** introduce **transfer functions for branches**
- **First approach:** attach (negated) conditions as **labels to control flow edges**
 - advantage: no language modification required
 - disadvantage: entails extension of DFA framework
 - will not further be considered here
- **Second approach:** encode conditions as **assertions** (statements)
 - advantage: DFA framework can be reused
 - disadvantage: entails extension of WHILE language
 - the way we will follow

Extending the Syntax of WHILE Programs

Definition (Labelled WHILE programs with assertions)

The **syntax of labelled WHILE programs with assertions** is defined by the following context-free grammar:

$$\begin{aligned} a &::= z \mid x \mid a_1 + a_2 \mid a_1 - a_2 \mid a_1 * a_2 \in AExp \\ b &::= t \mid a_1 = a_2 \mid a_1 > a_2 \mid \neg b \mid b_1 \wedge b_2 \mid b_1 \vee b_2 \in BExp \\ c &::= [\text{skip}]' \mid [x := a]' \mid c_1 ; c_2 \mid \\ &\quad \text{if } [b]' \text{ then } c_1 \text{ else } c_2 \mid \text{while } [b]' \text{ do } c \mid [\text{assert } b]' \in Cmd \end{aligned}$$

To be done:

- Definition of **transfer functions** for **assert** blocks (depending on analysis problem)
- Idea: assertions as **filters** that let only “valid” information pass

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Interval Analysis with Assertions I

So far:

- The domain (Int, \subseteq) of **intervals over \mathbb{Z}** is defined by
$$Int := \{[z_1, z_2] \mid z_1 \in \mathbb{Z} \cup \{-\infty\}, z_2 \in \mathbb{Z} \cup \{+\infty\}, z_1 \leq z_2\} \cup \{\emptyset\}$$

where

- $-\infty \leq z, z \leq +\infty$, and $-\infty \leq +\infty$ (for all $z \in \mathbb{Z}$)
- $\emptyset \subseteq J$ (for all $J \in Int$)
- $[y_1, y_2] \subseteq [z_1, z_2]$ iff $y_1 \geq z_1$ and $y_2 \leq z_2$
- Transfer functions** $\{\varphi_I \mid I \in Lab\}$ are defined by

$$\varphi_I(\delta) := \begin{cases} \delta & \text{if } B^I = \text{skip or } B^I \in BExp \\ \delta[x \mapsto val_\delta(a)] & \text{if } B^I = (x := a) \end{cases}$$

where

$$\begin{array}{ll} val_\delta(x) := \delta(x) & val_\delta(a_1 + a_2) := val_\delta(a_1) \oplus val_\delta(a_2) \\ val_\delta(z) := [z, z] & val_\delta(a_1 - a_2) := val_\delta(a_1) \ominus val_\delta(a_2) \\ & val_\delta(a_1 * a_2) := val_\delta(a_1) \odot val_\delta(a_2) \end{array}$$

with

$$\emptyset \oplus J := J \oplus \emptyset := \emptyset \ominus J := \dots := \emptyset$$

$$[y_1, y_2] \oplus [z_1, z_2] := [y_1 + z_1, y_2 + z_2]$$

$$[y_1, y_2] \ominus [z_1, z_2] := [y_1 - z_2, y_2 - z_1]$$

$$[y_1, y_2] \odot [z_1, z_2] := [\bigsqcap\{y_1 z_1, y_1 z_2, y_2 z_1, y_2 z_2\}, \bigsqcup\{y_1 z_1, y_1 z_2, y_2 z_1, y_2 z_2\}]$$

Interval Analysis with Assertions II

Additionally for $B^I = (\text{assert } b)$, $\delta : \text{Var}_c \rightarrow \text{Int}$ and $x \in \text{Var}_c$:

$$\varphi_I(\delta)(x) := \begin{cases} \emptyset & \text{if } Z = \emptyset \\ \left[\prod_{Z \cup \{-\infty\}} Z, \bigsqcup_{Z \cup \{+\infty\}} Z \right] & \text{otherwise} \end{cases}$$

where

- $Z := \{\sigma(x) \mid \sigma \in \Sigma_\delta, \text{val}_\sigma(b) = \text{true}\}$
- $\Sigma_\delta := \{\sigma : \text{Var}_c \rightarrow \mathbb{Z} \mid \forall y \in \text{Var}_c : \sigma(y) \in \delta(y)\}$
(and thus $\Sigma_\delta = \emptyset$ iff $\delta(y) = \emptyset$ for some $y \in \text{Var}_c$)
- $\text{val}_\sigma : \text{BExp} \rightarrow \mathbb{B}$ as before

Example 9.1

$$\text{Var}_c = \{x, y\}, \delta = (\underbrace{[-\infty, 2]}_x, \underbrace{[0, +\infty]}_y)$$

$$\Rightarrow \varphi_{\text{assert } x > 0}(\delta) = ([1, 2], [0, +\infty])$$

$$\varphi_{\text{assert } x=y}(\delta) = ([0, 2], [0, 2])$$

$$\varphi_{\text{assert } x > y}(\delta) = ([1, 2], [0, 1])$$

$$\varphi_{\text{assert } x < y}(\delta) = ([-\infty, 2], [0, +\infty])$$

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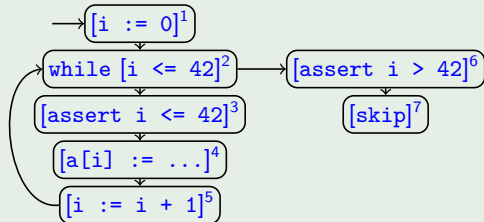
$$\varphi_{\text{assert } x<y}(\delta) = ([-\infty, 2], [0, +\infty])$$

Remarks:

- Again for $B^l = (\text{assert } b)$ and $\delta : Var_c \rightarrow Int$, $\varphi_l(\delta) \sqsubseteq \delta$ and hence $\Sigma_{\varphi_l(\delta)} \subseteq \Sigma_\delta$ (“filter”)
- Again if $Al_l(x) = \emptyset$ for some $l \in Lab_c$ and $x \in Var_c$, then l is **unreachable** (and $Al_l(y) = \emptyset$ for all $y \in Var_c$)

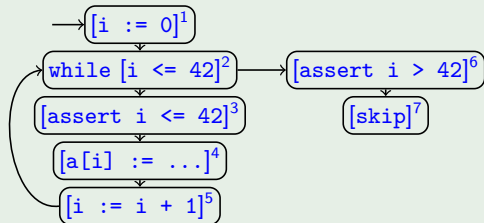
Interval Analysis with Assertions IV

Example 9.2 (Interval analysis for array index; cf. Example 7.6)



Interval Analysis with Assertions IV

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$$\varphi_1(J) = [0, 0]$$

$$\varphi_2(J) = J$$

$$\varphi_3(J) = J \cap [-\infty, 42]$$

$$\varphi_4(J) = J$$

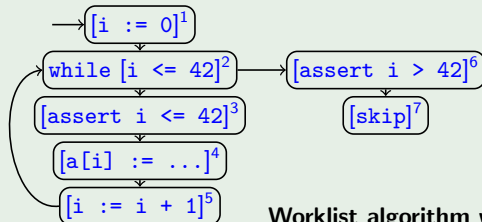
$$\varphi_5(\emptyset) = \emptyset$$

$$\varphi_5([i_1, i_2]) = [i_1 + 1, i_2 + 1]$$

$$\varphi_6(J) = J \cap [43, +\infty]$$

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Example 9.2 (Interval analysis for array index; cf. Example 7.6)



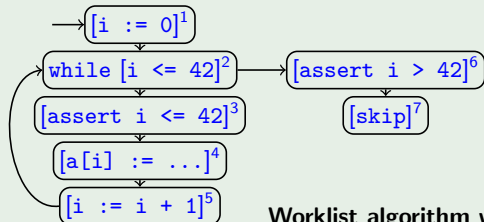
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Worklist algorithm with widening:

W	AI_1	AI_2	AI_3	AI_4	AI_5	AI_6	AI_7
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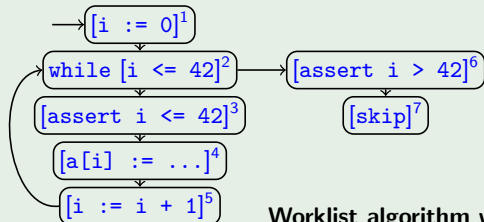
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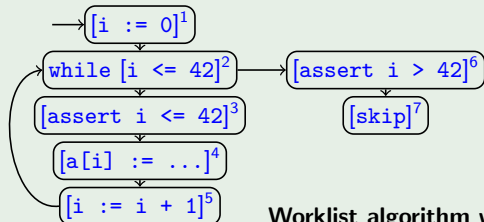
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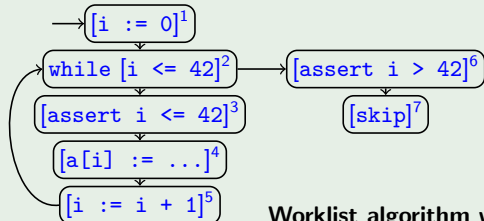
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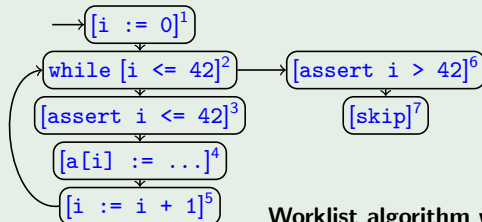
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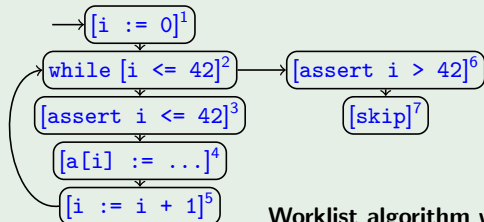
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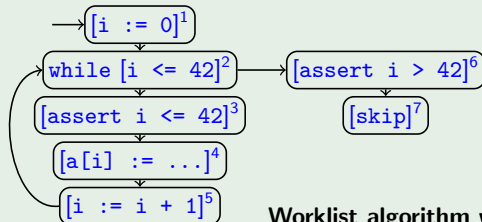
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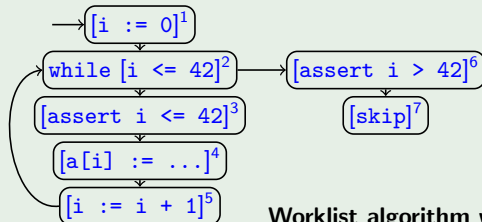
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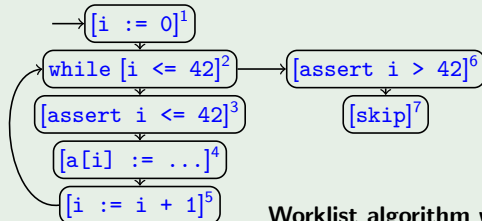
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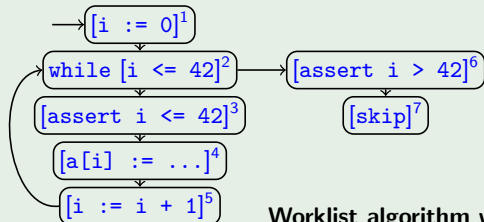
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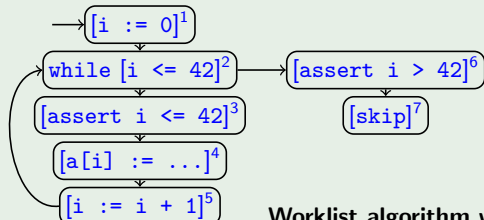
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52, 26, 67	$[-\infty, +\infty]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
23, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
34, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
45, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 0]$	\emptyset	\emptyset
52, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	\emptyset	\emptyset
26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	\emptyset	\emptyset
67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	\emptyset

Interval Analysis with Assertions IV

Example 9.2 (Interval analysis for array index; cf. Example 7.6)

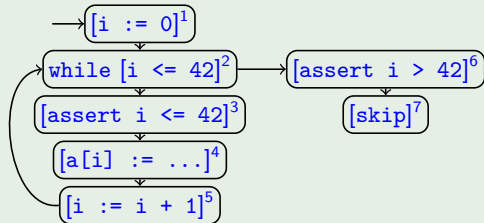


$$\begin{aligned} \varphi_1(J) &= [0, 0] \\ \varphi_2(J) &= J \\ \varphi_3(J) &= J \cap [-\infty, 42] \\ \varphi_4(J) &= J \\ \varphi_5(\emptyset) &= \emptyset \\ \varphi_5([i_1, i_2]) &= [i_1 + 1, i_2 + 1] \\ \varphi_6(J) &= J \cap [43, +\infty] \end{aligned}$$

Worklist algorithm with widening:

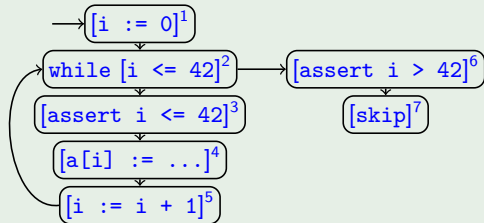
W	AI ₁	AI ₂	AI ₃	AI ₄	AI ₅	AI ₆	AI ₇
12, 23, 34, 45, 52, 26, 67	$[-\infty, +\infty]$	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset
23, 34, 45, 52, 26, 67	$[-\infty, +\infty]$	$[0, 0]$	\emptyset	\emptyset	\emptyset	\emptyset	\emptyset
34, 45, 52, 26, 67	$[-\infty, +\infty]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset	\emptyset	\emptyset
45, 52, 26, 67	$[-\infty, +\infty]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset	\emptyset
52, 26, 67	$[-\infty, +\infty]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
23, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, 0]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
34, 26, 67	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 0]$	$[0, 0]$	\emptyset	\emptyset
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ϵ	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$

Example 9.2 (Interval analysis for array index; continued)



Interval Analysis with Assertions V

Example 9.2 (Interval analysis for array index; continued)



$$\varphi_1(J) = [0, 0]$$

$$\varphi_2(J) = J$$

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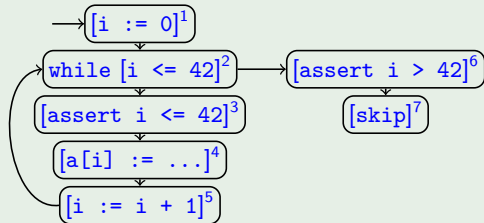
$$\varphi_5(\emptyset) = \emptyset$$

$$\varphi_5([i_1, i_2]) = [i_1 + 1, i_2 + 1]$$

$$\varphi_6(J) = J \cap [43, +\infty]$$

Interval Analysis with Assertions V

Example 9.2 (Interval analysis for array index; continued)



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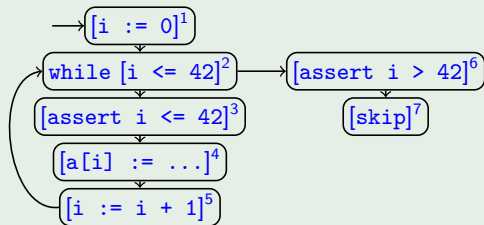
$$\varphi_6(J) = J \cap [43, +\infty]$$

Narrowing:

	Al_1	Al_2	Al_3	Al_4	Al_5	Al_6	Al_7
$\text{fix}^\nabla(\Phi_S)$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$

Interval Analysis with Assertions V

Example 9.2 (Interval analysis for array index; continued)



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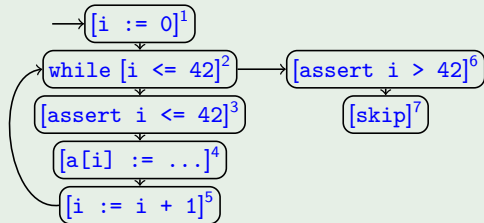
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$\Phi_S(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$

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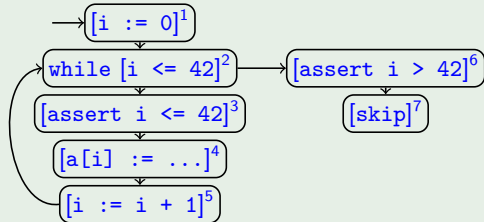
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$\Phi_S(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$
$\Phi_S^2(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$

Interval Analysis with Assertions V

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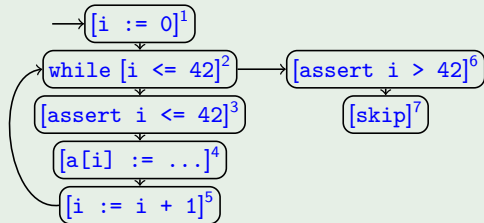
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$\Phi_S(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$
$\Phi_S^2(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$
$\Phi_S^3(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, 43]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$

Interval Analysis with Assertions V

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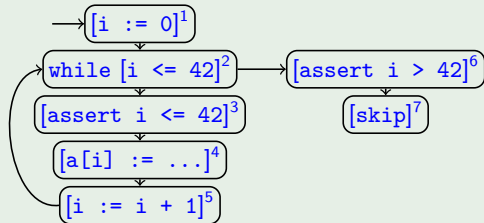
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$\Phi_S^2(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$
$\Phi_S^3(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, 43]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$
$\Phi_S^4(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, 43]$	$[0, 43]$	$[0, 42]$	$[0, 42]$	$[0, 43]$	$[43, +\infty]$

Interval Analysis with Assertions V

Example 9.2 (Interval analysis for array index; continued)



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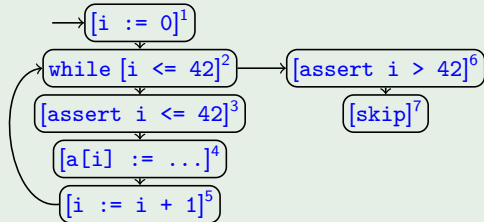
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Narrowing:

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$\Phi_S(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$
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$\Phi_S^4(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, 43]$	$[0, 43]$	$[0, 42]$	$[0, 42]$	$[0, 43]$	$[43, +\infty]$
$\Phi_S^5(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, 43]$	$[0, 43]$	$[0, 42]$	$[0, 42]$	$[0, 43]$	$[43, 43]$

Example 9.2 (Interval analysis for array index; continued)



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	Al ₁	Al ₂	Al ₃	Al ₄	Al ₅	Al ₆	Al ₇
$\text{fix}^\nabla(\Phi_S)$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[43, +\infty]$
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$\Phi_S^2(\text{fix}^\nabla(\Phi_S))$	$[-\infty, +\infty]$	$[0, +\infty]$	$[0, +\infty]$	$[0, 42]$	$[0, 42]$	$[0, +\infty]$	$[43, +\infty]$
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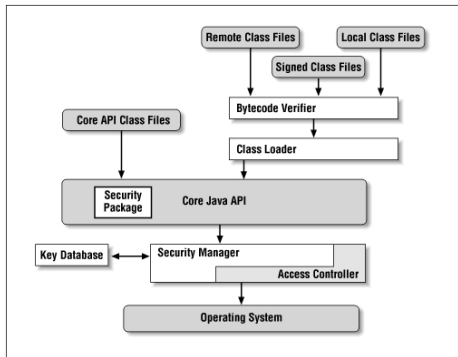
- 1 Recap: Taking Conditional Branches into Account
- 2 Interval Analysis with Assertions
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- Execution on **Java Virtual Machine** (JVM)
- **Advantages:**
 - architecture independency (especially for web applications)
 - faster than pure (i.e., source code) interpretation
- **Problem: security issues**
 - destruction of data
 - modification of data
 - disclosure of personal information
 - modification of other programs

Java Security: the Sandbox

- **Insulation layer** providing indirect access to system resources
- Hardware access via **API classes and methods**
- **Bytecode verification** upon uploading
 - well-typedness
 - proper object referencing
 - proper control flow



- Conventional **stack-based abstract machine**
- Supports **object-oriented features**: classes, methods, etc.
- **Stack** for intermediate results of expression evaluations
- **Registers** for source-level local variables and method parameters
- Both part of **method activation record**
(and thus preserved across method calls)
- Method entry point specifies **required number** of registers (m_r) and stack slots (m_s ; for memory allocation)
- (Most) instructions are **typed**

Example 9.3 (Factorial function)

Java source code:

```
static int factorial(int n)
{ int res;
  for (res = 1; n > 0; n--) res = res * n;
  return res; }
```

Example: Factorial Function

Example 9.3 (Factorial function)

Java source code:

```
static int factorial(int n)
{ int res;
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  return res; }
```

Corresponding JVM bytecode:

```
method static int factorial(int), 2 registers, 2 stack slots
 1: istore 0      // store n in register 0
 2: iconst_1     // push constant 1
 3: istore 1     // store res in register 1
 4: iload 0      // push n
 5: ifle 12      // if <= 0, go to end
 6: iload 1      // push res
 7: iload 0      // push n
 8: imul        // res * n on top of stack
 9: istore 1     // store in res
10: iinc 0, -1   // decrement n
11: goto 4       // go to loop header
12: iload 1      // push res
13: ireturn     // return res to caller
```

JVM Instruction Set (excerpt)

`iload n` : push integer from register n

`istore n` : pop integer into register n

`iconst_ z` : push integer z

`aconst_null`: push null reference

`iadd`: add two topmost integers on stack and push sum

`getfield C f τ` : pops reference to object (of class C) and pushes value of field f (of type τ)

`putfield C f τ` : pops value v (of type τ) and reference to object o (of class C) and assigns v to field f of o

`new C` : creates new object (of class C) and pushes reference

`invoke C M $\tau_0(\tau_1, \dots, \tau_n)$` : pops values v_1, \dots, v_n (of type τ_1, \dots, τ_n) and reference to object (of class C), calls method M with parameters v_1, \dots, v_n , and pushes return value (of type τ_0)

`if_icmpeq l` : pop two topmost integers from stack and jump to line l if equal

`ireturn`: return to caller with integer result on top of stack

(\approx 200 instructions in total)

Example 9.4 (Malicious bytecode)

```
1: iconst_5  
2: iconst_1  
3: putfield A f int
```

interprets second stack entry (5) as reference to object of class `A` and assigns first stack entry (1) to field `f` of this object

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Correctness of Bytecode

Conditions to ensure **proper operation**:

Type correctness: arguments of instructions always of expected type

No stack over-/underflow: never push to full stack or pop from empty stack

Code containment: PC must always point into the method code

Register initialization: load from non-parameter register only after store

Object initialization: constructor must be invoked before using class instance

Access control: operations must respect visibility modifiers
(**private/protected/public**)

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Options:

- **dynamic checking** at execution time (“defensive JVM approach”)
 - expensive, slows down execution
- **static checking** at loading time (here)
 - verified code executable at full speed without extra dynamic checks

Summary: **dataflow analysis** applied to **type-level abstract interpretation** of JVM

- ① Association of **type information** with register and stack contents
 - set of types forms a complete lattice
- ② Simulation of **execution of instructions** at type level
- ③ Use **dataflow analysis** to cover all concrete executions
- ④ Modularity: analysis proceeds method **per method**

(see X. Leroy: *Java Bytecode Verification: Algorithms and Formalizations*, Journal of Automated Reasoning 30(3-4), 2003, 235–269)

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The **set of types**, *Typ*, is composed of

- **Primitive** types:
 - *int* (covering *boolean*, *byte*, *char*, *short*)
 - *long*
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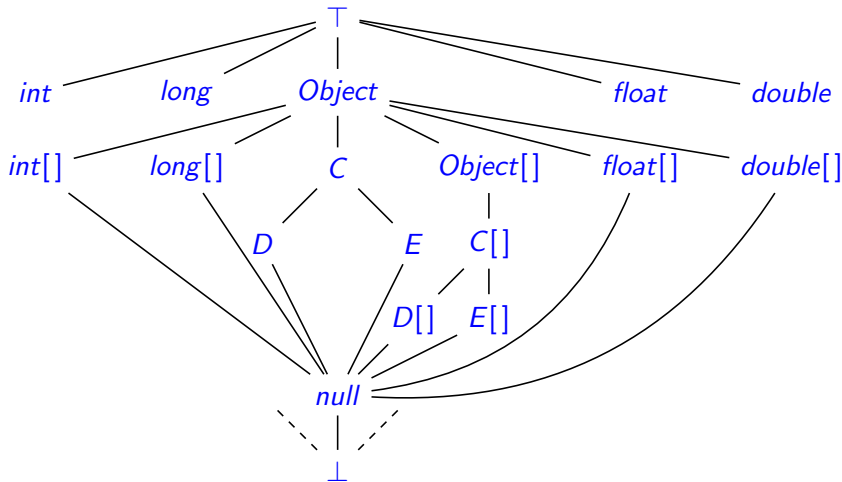
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- **Method** types: $\tau_0(\tau_1, \dots, \tau_n)$ for $n \in \mathbb{N}$, $\tau_i \in Typ$
- **Special** types:
 - $null$ (null reference)
 - $Object$ (any object)
 - \top (contents of uninitialized registers, i.e., any value)
 - \perp (absence of any value)

The Subtyping Relation (excerpt)

(C , D , E user-defined classes; D , E extending C)



Notation: $\tau_1 \sqsubseteq_t \tau_2$

The Type-Level Abstract Interpreter I

- **Idea:** execute JVM instructions on **types** (rather than concrete values)
 - **stack type** $S \in \text{Typ}^{\leq m_s}$ (top to the left)
 - **register type** $R : \{0, \dots, m_r - 1\} \rightarrow \text{Typ}$

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- Represented as **transition relation**

$$i : (S, R) \rightarrow (S', R')$$

where

- i : current instruction
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- i : current instruction
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- (S', R') : stack/register type after execution
- **Errors** (type mismatch, stack over-/underflow, ...) denoted by absence of transition

The Type-Level Abstract Interpreter II

Some transition rules:

<code>iconst_z :</code>	$(S, R) \rightarrow (int.S, R)$	if $ S < m_s$
<code>aconst_null :</code>	$(S, R) \rightarrow (null.S, R)$	if $ S < m_s$
<code>iadd :</code>	$(int.int.S, R) \rightarrow (int.S, R)$	
<code>if_icmpeq l :</code>	$(int.int.S, R) \rightarrow (S, R)$	
<code>iload n :</code>	$(S, R) \rightarrow (int.S, R)$	if $0 \leq n < m_r, R(n) = int, S < m_s$
<code>aload n :</code>	$(S, R) \rightarrow (R(n).S, R)$	if $0 \leq n < m_r, R(n) \sqsubseteq_t Object, S < m_s$
<code>istore n :</code>	$(int.S, R) \rightarrow (S, R[n \mapsto int])$	if $0 \leq n < m_r$
<code>astore n :</code>	$(\tau.S, R) \rightarrow (S, R[n \mapsto \tau])$	if $0 \leq n < m_r, \tau \sqsubseteq_t Object$
<code>getfield C f τ :</code>	$(D.S, R) \rightarrow (\tau.S, R)$	if $D \sqsubseteq_t C$
<code>putfield C f τ :</code>	$(\tau'.D.S, R) \rightarrow (S, R)$	if $\tau' \sqsubseteq_t \tau, D \sqsubseteq_t C$
<code>invoke C M σ :</code>	$(\tau'_n \dots \tau'_1.\tau'.S, R) \rightarrow (\tau_0.S, R)$	if $\sigma = \tau_0(\tau_1, \dots, \tau_n), \tau'_i \sqsubseteq_t \tau_i$ for $1 \leq i \leq n, \tau' \sqsubseteq_t C$

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- 3 (*Soundness*) If $i : (S, R) \rightarrow (S', R')$, then for all concrete states (s, r) matching (S, R) , the defensive JVM will not stop with a run-time type exception when applying i to (s, r) (but rather change to some (s', r') matching (S', R')).

Some Theoretical Properties

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Proof.

see X. Leroy: *Java Bytecode Verification: Algorithms and Formalizations* □