

Exercise Sheet 7

Due date: June 21th. Please hand in your solutions at the start of the exercise class.

Task 1: Interpretation of Hoare Triples (25 Points)

Match each of the following Hoare triples

- (a) $\{P\} c \{\text{true}\}$
- (b) $\{P\} c \{\Downarrow \text{true}\}$
- (c) $\{\text{true}\} c \{\text{false}\}$
- (d) $\{\text{true}\} c \{\Downarrow \text{false}\}$
- (e) $\{\text{false}\} c \{\Downarrow Q\}$

with their corresponding interpretation from 1–6. (There may be more than one Hoare triple with the same interpretation in 1–6.) Assume that P and Q contain no logical variable.

1. Program c diverges for all initial states.
2. The program never finishes in a final state satisfying Q .
3. The triple does not say anything about program c ; it is logically equivalent to **true**.
4. The triple is logically equivalent to **false**.
5. None of the above. Provide yourself the interpretation of the triple.
6. Program c always terminate whenever executed in an initial state that satisfies P .

Task 2: Derived Rules (25 Points)

Which of the following rules hold true? For valid rules, provide a justification (no formal proof required). For rules that do not hold, provide a counterexample.

- (a)
$$\frac{\{P\} c \{Q_1\} \quad \{P\} c \{Q_2\}}{\{P\} c \{Q_1 \wedge Q_2\}}$$
- (b)
$$\frac{\{P\} c \{Q_1\} \quad \{P\} c \{Q_2\}}{\{P\} c \{Q_1 \vee Q_2\}}$$
- (c)
$$\frac{\{P_1\} c \{Q\} \quad \{P_2\} c \{Q\}}{\{P_1 \wedge P_2\} c \{Q\}}$$
- (d)
$$\frac{\{P_1\} c \{Q\} \quad \{P_2\} c \{Q\}}{\{P_1 \vee P_2\} c \{Q\}}$$
- (e)
$$\frac{\{P\} c \{Q\}}{\{P \wedge R\} c \{Q \wedge R\}}$$
- (f)
$$\frac{\{P\} c \{Q\}}{\{P \wedge R\} c \{Q \vee R\}}$$

Task 3: Strongest Postconditions (50 Points)

Intuitively, the *strongest liberal postcondition* $\text{s1p}[\![c]\!](A)$ of a program c and a precondition A is the strongest postcondition B that holds when running c on a state satisfying A . In contrast to weakest liberal preconditions, we thus apply forward reasoning.

- (a) Formalize the intuitive definition of strongest liberal postconditions from above, i.e. give an exact definition of the set of states described by $\text{s1p}[\![c]\!](A)$.
- (b) Express the validity of a Hoare triple $\{A\}c\{B\}$ in terms of strongest liberal postconditions (cf. Lecture 11, Corollary 11.5.1).
- (c) Determine $\text{s1p}[\!\text{while true do skip end}]\!(\text{true})$.
- (d) Define $\text{s1p}[\![c]\!](A)$ by induction on the structure of c .
- (e) Apply your definition from (d) to compute the following strongest liberal postcondition:

$$\text{s1p}[\![x := 2 * x; y := x + 2; z := y + x]\!](x = 1)$$

- (f) Prove or disprove: For every program $c \in \text{Cmd}$, we have

$$\models \{\text{w1p}[\![c]\!](\text{s1p}[\![c]\!](\text{false}))\}c\{\text{s1p}[\![c]\!](\text{w1p}[\![c]\!](\text{false}))\} .$$