

# Introduction to Model Checking (Summer Term 2018)

## — Exercise Sheet 6 (due 11th June) —

### General Remarks

- The exercises are to be solved in groups of *three* students.
- You may hand in your solutions for the exercises just before the exercise class starts at 12:15 or by dropping them into the “Introduction to Model Checking” box at our chair *before 12:00*. Do *not* hand in your solutions via L2P or via e-mail.

### Exercise 1

(1+2+3+4 Points)

Consider the transition system  $TS_{Sem}$  for mutual exclusion with a semaphore.

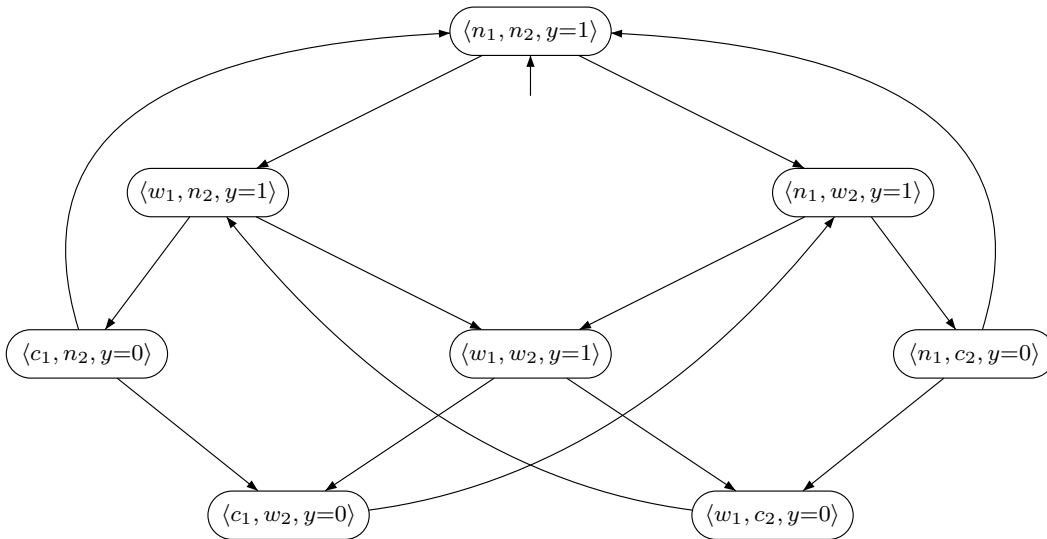


Figure 6.1: Mutual exclusion with semaphore (transition system representation).

Let  $P_{live}$  be the following  $\omega$ -regular property over  $AP = \{wait_1, crit_1\}$ :

“whenever process 1 is waiting for the critical section,  
it will eventually (potentially at the very same time) be in its critical section.”

Check whether  $TS_{Sem} \models P_{live}$  with the following steps:

- Introduce the necessary labels in  $TS_{Sem}$ .
- Depict an NBA  $\bar{A}$  for the complement property  $\overline{P_{live}} = (2^{AP})^\omega \setminus P_{live}$ .  
*Hint:* There is an NBA  $\bar{A}$  for  $\overline{P_{live}}$  with 3 states.

- (c) Depict the reachable fragment of the product  $TS_{Sem} \otimes \bar{\mathcal{A}}$ .  
*Hint:* There is an NBA for  $\bar{\mathcal{A}}$  with 3 states that is a solution to task (b) and will lead to a product transition system with 19 states.
- (d) Apply the nested depth-first search (lecture 11, slides 150 and 159) to  $TS_{Sem} \otimes \bar{\mathcal{A}}$  for the persistence property “eventually forever  $\neg F$ ”, where  $F$  is the acceptance set of  $\bar{\mathcal{A}}$ . To illustrate the steps:
- before each *Pop* operation give:
    - for the first DFS the contents of stack  $\pi$  and set  $U$ , and
    - for the second DFS the contents of stack  $\xi$  and set  $V$ .
  - indicate whenever *cycle\_check()* is called.
- Does  $TS_{Sem} \models P_{live}$  hold? In case the property is refuted, give the counterexample returned by the algorithm.

## Exercise 2★

(4 Points)

Recall the following LT properties from exercise sheet 3.

- (i) **“Winter is coming.”**  
 $P_1 = \emptyset^* \{winter\} (2^{AP})^\omega$
- (ii) **“Everything is awesome.”**  
 $P_2 = \{awesome\}^\omega$
- (iii) **“I’ll be back.”**  
 $P_3 = \{here\}^+ \emptyset^+ \{here\}^+ (2^{AP})^\omega$
- (iv) **“You either die a hero, or you live long enough to see yourself become the villain.”**  
 $P_4 = \{live, hero\}^+ \{hero\} (2^{AP})^\omega + \{live, hero\}^+ \{live\} (2^{AP})^\omega$
- (v) **“By night one way, by day another  
 Thus shall be the norm  
 Till you receive true love’s kiss  
 then, take love’s true form.”**  

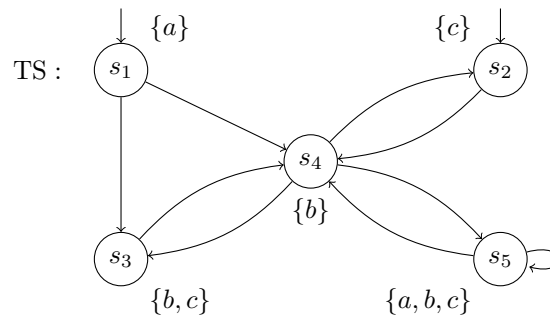
$$P_5 = ((\{form_1\} \{day, form_2\})^+ + \{form_1\} (\{day, form_2\} \{form_1\})^*) \{kiss, true\_form\} \\ (\{true\_form\} \{true\_form, day\})^\omega$$
- (vi) **“A Lannister always pays his debts.”**  
 $P_6 = \emptyset^* (\{in\_debt\}^+ \emptyset^+)^* \emptyset^\omega$
- (vii) **“Anything is possible [if you just believe]”**  
 $P_7 = (2^{AP})^\omega$
- (viii) **“It’s gonna be legen... wait for it... dary!”**  
 $P_8 = \{legen\} \{wait\_for\_it\}^+ \{dary\} (2^{AP})^\omega$

Express each property  $P_i$  as an LTL formula  $\varphi_i$ .

**Exercise 3★**

**(6 Points)**

Consider the following transition system TS where we omit the transition labels, because they are all  $\tau$ .



For each of the LTL formulae  $\varphi_i$  below, decide whether  $TS \models \varphi_i$ . Justify your answer and, in particular, provide a path  $\pi_i \in Paths(TS)$  such that  $\pi_i \not\models \varphi_i$  in case you find  $TS \not\models \varphi_i$ .

- $\varphi_1 = \diamond \square c$ ,
- $\varphi_2 = \square \diamond c$ ,
- $\varphi_3 = \bigcirc \neg c \rightarrow \bigcirc \bigcirc c$ ,
- $\varphi_4 = \square a$ ,
- $\varphi_5 = a \text{ U } \square (b \vee c)$ ,
- $\varphi_6 = (\bigcirc \bigcirc b) \text{ U } (b \vee c)$ ,
- $\varphi_7 = c \text{ R } b$ ,

where the *release operator*  $\varphi \text{ R } \psi$  for two LTL formulae  $\varphi, \psi$  is defined by  $\varphi \text{ R } \psi \equiv \neg(\neg\varphi \text{ U } \neg\psi)$ .