Theoretical Foundations of the UML

Lecture 10: Realisability

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Outline

1. Introduction
2. Properties of CFMs
   - Deterministic CFMs
   - Deadlock-free CFMs
   - Synchronisation messages add expressiveness
3. Realisability
4. Inference of MSCs
5. Characterisation and complexity of realisability by weak CFMs

Overview

1. Introduction
2. Properties of CFMs
   - Deterministic CFMs
   - Deadlock-free CFMs
   - Synchronisation messages add expressiveness
3. Realisability
4. Inference of MSCs
5. Characterisation and complexity of realisability by weak CFMs

Motivation

Practical use of MSCs and CFMs

- MSCs and MSGs are used by software engineers to capture requirements.
- These are the expected behaviours of the distributed system under design.
- Distributed systems can be viewed as a collection of communicating automata.

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Central problem

Can we synthesize, preferably in an automated manner, a CFM whose behaviours are precisely the behaviours of the MSCs (or MSG)?

This is known as the realisability problem.

From requirements to implementation

Realisability problem

INPUT: a set of MSCs

OUTPUT: a CFM $\mathcal{A}$ such that $L(\mathcal{A})$ equals the set of input MSCs.

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2. If so, how complex is it to obtain such CFM?

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Questions:

1. Is this possible? (That is, is this decidable?)
2. If so, how complex is it to obtain such CFM?
3. If so, how do such algorithms work?

Problem variants (1)

Realisability problem

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- Consider MSCs whose set of linearisations is a regular word language.

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- Consider MSGs, that may describe an infinite set of MSCs.
- Consider MSCs whose set of linearisations is a regular word language.
- Consider MSGs that are non-local choice.

Problem variants (2)

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Different system models

- Consider CFMs without synchronisation messages.

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Different system models

- Consider CFMs without synchronisation messages.
- Allow CFMs that may deadlock. Possibly, a realisation deadlocks.
- Forbid CFMs that deadlock. No realisation will ever deadlock.

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- Consider CFMs that are deterministic.

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- Consider CFMs that are bounded.

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Realisation of a finite set of MSCs by a CFM without synchronisation messages and that may possibly deadlock.

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Realisation of a finite set of well-formed words (= language) by a CFM without synchronisation messages and that may possibly deadlock.

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1. CFMs without synchronisation messages are weaker than CFMs.

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1. CFMs without synchronisation messages are weaker than CFMs.
2. Conditions for realisability of a finite set of MSCs by a weak CFM.

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Results:

1. CFMs without synchronisation messages are weaker than CFMs.
2. Conditions for realisability of a finite set of MSCs by a weak CFM.
3. Checking realisability for such sets is co-NP complete.

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Determinism

Definition (Deterministic CFM)

A CFM $\mathcal{A}$ is *deterministic* if for all $p \in \mathcal{P}$, the transition relation Δ_p satisfies the following two conditions:

1. $(s, !(p, q, (a, m_1)), s_1) \in \Delta_p$ and $(s, !(p, q, (a, m_2)), s_2) \in \Delta_p$ implies $m_1 = m_2$ and $s_1 = s_2$
2. $(s, ?(p, q, (a, m)), s_1) \in \Delta_p$ and $(s, ?(p, q, (a, m)), s_2) \in \Delta_p$ implies $s_1 = s_2$

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Note:

From a given state, process p may have the possibility of sending messages to more than one process.

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Note:

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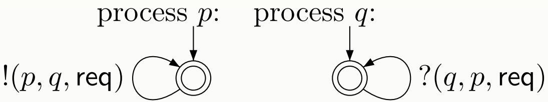
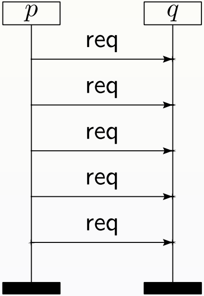
Example:

Example CFM (1) and (2) are deterministic, while (3) is not.

Example (1)

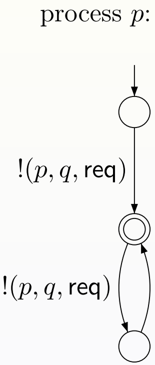


process p:
process q:
$!(p, q, \mathsf{req})$
$?(q, p, \mathsf{req})$



p
q
req
req
req
req
req

Example (2)


process p:
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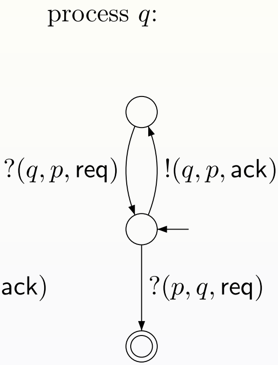
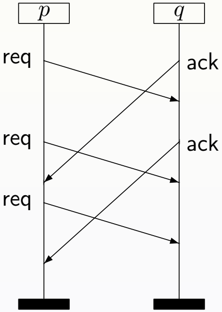

process q:
$?(q, p, \mathsf{req})$
$!(q, p, \mathsf{ack})$
$?(p, q, \mathsf{req})$


p
q
req
ack
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ack
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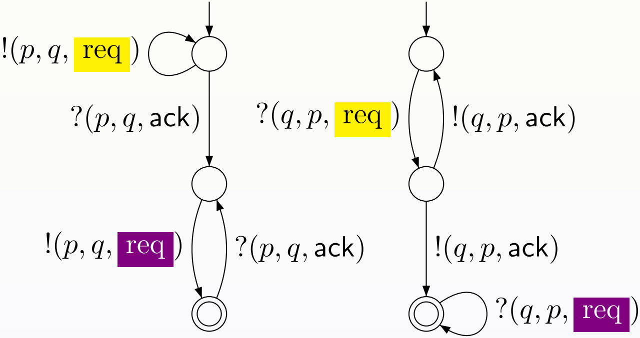
Example (3)



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Deadlock-freeness

Definition (Deadlock-free CFM)

A CFM $\mathcal{A}$ is *deadlock-free* if, for all $w \in Act^*$ and all runs γ of $\mathcal{A}$ on w, there exist $w' \in Act^*$ and run γ' in $\mathcal{A}$ such that $\gamma \cdot \gamma'$ is an accepting run of $\mathcal{A}$ on $w \cdot w'$.

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Theorem: [Genest et. al, 2006]

For any $\exists B$-bounded CFM $\mathcal{A}$, the decision problem "is $\mathcal{A}$ deadlock-free?" is decidable (and is PSPACE-complete).

Weak CFMs

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Q: Are CFMs more expressive than weak CFMs? That is, do there exist languages (over linearizations or, equivalently, MSCs) that can be generated by CFMs but **not** by weak CFMs?

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Q: Are CFMs more expressive than weak CFMs? That is, do there exist languages (over linearizations or, equivalently, MSCs) that can be generated by CFMs but **not** by weak CFMs? Yes.

CFM vs. weak CFM

Theorem:

Weak CFMs are strictly less expressive than CFMs.

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Proof.

For $m, n \geqslant 1$, let $M(m,n) \in \mathbb{M}$ over $\mathcal{P} = \{1,2\}$ and $\mathcal{C} = \{\text{req}, \text{ack}\}$ be:

- $M \upharpoonright 1 = (!(1,2,\text{req}))^m \; (?(1,2,\text{ack}) \; !(1,2,\text{req}))^n$
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Claim: there is no weak CFM over $\mathcal{P} = \{1,2\}$ and $\mathcal{C} = \{\text{req}, \text{ack}\}$ whose language is $L = \{M(n,n) \mid n > 0\}$.

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- $\mathcal{A}_1$ visits a cycle of length $i > 0$ to read the first n letters of $M(n,n) \upharpoonright 1$

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- $\mathcal{A}_2$ visits a cycle of length $j > 0$ to read the last n letters of $M(n,n) \upharpoonright 2$

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Then there is an accepting run of $\mathcal{A}$ on $M(n + (i \cdot j), n) \notin L$. $\square$

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Claim: there is no weak CFM over $\mathcal{P} = \{1,2\}$ and $\mathcal{C} = \{\text{req}, \text{ack}\}$ whose language is $L = \{M(n,n) \mid n > 0\}$. By contraposition. Suppose there is a weak CFM $\mathcal{A} = ((\mathcal{A}_1, \mathcal{A}_2), s_{init}, F)$ with $L(\mathcal{A}) = L$. For any $n > 0$, there is an accepting run of $\mathcal{A}$ on $M(n,n)$. If n is sufficiently large, then

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Then there is an accepting run of $\mathcal{A}$ on $M(n + (i \cdot j), n) \notin L$. Contradiction.

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Intuition proof

If $\mathcal{A}_1$ traverses a cycle of size i at least once to “generate” $(!(1,2,\mathrm{req}))^n$, then it can autonomously traverse this cycle more often and thus “pump” to an expression of the form $(!(1,2,\mathrm{req}))^{n \cdot i}$.

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Similar reasoning applies to automaton $\mathcal{A}_2$ for the last n letters of the input word $M \upharpoonright 2$. Suppose its cycle is of size j.

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Similar reasoning applies to automaton $\mathcal{A}_2$ for the last n letters of the input word $M \upharpoonright 2$. Suppose its cycle is of size j.

Now if $\mathcal{A}_1$ traverses its cycle of size i, j times, and $\mathcal{A}_2$ traverses its cycle of size j, i times, then the number of requests sent by process 1 matches the number of receipts by process 2.

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Theorem:

Weak CFMs are strictly less expressive than CFMs.

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Similar reasoning applies to automaton $\mathcal{A}_2$ for the last n letters of the input word $M \restriction 2$. Suppose its cycle is of size j.

Now if $\mathcal{A}_1$ traverses its cycle of size i, j times, and $\mathcal{A}_2$ traverses its cycle of size j, i times, then the number of requests sent by process 1 matches the number of receipts by process 2.

But this yields a word in $M(n + (i \cdot j), n)$ that is not in L.

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2. A finite set $\{M_1, \ldots, M_n\}$ of MSCs is realisable whenever $\{M_1, \ldots, M_n\} = \mathcal{L}(\mathcal{A})$ for some CFM $\mathcal{A}$.

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We will consider realisability using its characterisation by linearisations.

Overview

1 Introduction

2 Properties of CFMs
- Deterministic CFMs
- Deadlock-free CFMs
- Synchronisation messages add expressiveness

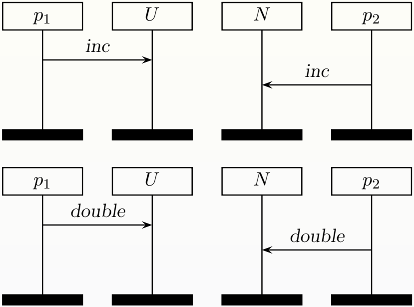
3 Realisability

4 Inference of MSCs

5 Characterisation and complexity of realisability by weak CFMs

Two example MSCs

Consider the MSCs M_{inc} (top) and M_{db} (bottom):



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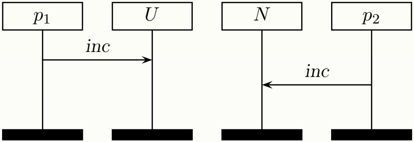
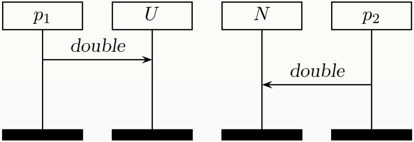
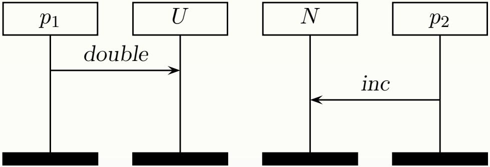




Intuition

In M_{inc}, the volume of U (uranium) and N (nitric acid) is increased by one unit; in M_{db} both volumes are doubled. For safety reasons, it is essential that both ingredients are increased by the same amount!

A third, inferred fatal scenario



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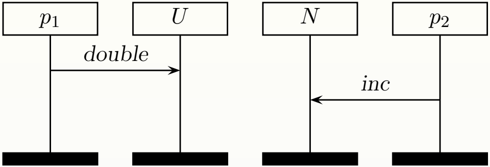

So:

The set $\{ M_{inc}, M_{db} \}$ is not realisable, as any CFM that realises this set also realises the inferred MSC M_{bad} above.

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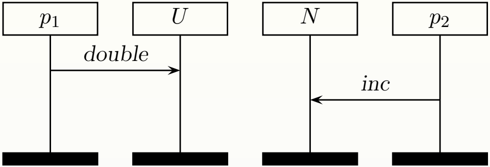

So:

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Note that:

MSCs M_{inc} or M_{db} alone do not imply M_{bad}. Together they do.

Inference

Definition (Inference)

The set L of MSCs is said to infer MSC $M \notin L$ if and only if:

$$\text{for any CFM } \mathcal{A}. \ (L \subseteq \mathcal{L}(\mathcal{A}) \text{ implies } M \in \mathcal{L}(\mathcal{A}))\,.$$

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What we will show later on:

The set L of MSCs is realisable iff L contains all MSCs that it infers.

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A realisable set of MSCs contains all its implied scenarios.

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For computational purposes, an alternative characterisation is required.

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Definition (MSC projection)

For MSC M and process p let $M \upharpoonright p$, the projection of M on process p, be the ordered sequence of actions occurring at process p in M.

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For MSC M and process p let $M \restriction p$, the projection of M on process p, be the ordered sequence of actions occurring at process p in M.

Lemma

An MSC M over the processes $\mathcal{P} = \{ p_1, \ldots, p_n \}$ is uniquely determined by the projections $M \restriction p_i$ for $0 < i \leqslant n$.

Projection (2)

Definition (Word projection)

For word $w \in Act^*$ and process p, the projection of w on process p, denoted $w \upharpoonright p$, is defined by:

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\begin{array}{rcl}
\epsilon \upharpoonright p & = & \epsilon \\
(!(r,q,a) \cdot w) \upharpoonright p & = & \begin{cases} !(r,q,a) \cdot (w \upharpoonright p) & \text{if } r = p \\ w \upharpoonright p & \text{otherwise} \end{cases}
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$w = !(1,2,\mathrm{req})!(1,2,\mathrm{req})?(2,1,\mathrm{req})!(2,1,\mathrm{ack})?(2,1,\mathrm{req})!(2,1,\mathrm{ack})?(1,2,\mathrm{ack})!(1,2,\mathrm{req})$

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Lemma

A well-formed word w over Act^* given as projections $w \upharpoonright p_1, \ldots, w \upharpoonright p_n$ uniquely characterises an MSC $M(w)$ over $\mathcal{P} = \{ p_1, \ldots, p_n \}$.

Closure

Definition (Inference relation)

For well-formed[a] $L \subseteq Act^*$, and well-formed word $w \in Act^*$, let:

$$L \models w \quad \text{iff} \quad (\forall p \in \mathcal{P}.\, \exists v \in L.\, w \upharpoonright p = v \upharpoonright p)$$

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Intuition

The closure condition says that the set of MSCs (or, equivalently, well-formed words) can be obtained from the projections of the MSCs in L onto individual processes.

Closure: example

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Realisability by a weak CFM

A finite set $\{M_1, \ldots, M_n\}$ of MSCs is realisable (by a weak CFM) whenever $\{M_1, \ldots, M_n\} = L(\mathcal{A})$ for some weak CFM $\mathcal{A}$

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For any weak CFM $\mathcal{A}$, $Lin(\mathcal{A})$ is closed under $\models$.

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In absence of synchronisation messages, the "local" accepting runs $\pi \restriction p$ for all processes p together can be combined to obtain an accepting run of $\mathcal{A}$ on w.

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To show that $Lin(\mathcal{A})$ is closed under $\models$, we prove that $w \in Lin(\mathcal{A})$.
By definition of $\models$, for any process p there is $v^p \in Lin(\mathcal{A})$ with $v^p \upharpoonright p = w \upharpoonright p$.
Let π be an accepting run of $\mathcal{A}$ on v^p and let run $\pi \upharpoonright p$ visit only states of $\mathcal{A}_p$ while taking only transitions in Δ_p. Then, $\pi \upharpoonright p$ is an accepting run of "local" automaton $\mathcal{A}_p$ on the word $v^p \upharpoonright p = w \upharpoonright p$.
In absence of synchronisation messages, the "local" accepting runs $\pi \upharpoonright p$ for all processes p together can be combined to obtain an accepting run of $\mathcal{A}$ on w.
Thus, $w \in Lin(\mathcal{A})$. $\square$

Overview

1. Introduction
2. Properties of CFMs
   - Deterministic CFMs
   - Deadlock-free CFMs
   - Synchronisation messages add expressiveness
3. Realisability
4. Inference of MSCs
5. Characterisation and complexity of realisability by weak CFMs

Characterisation of realisability

Theorem: [Alur et al., 2001]

Finite $L \subseteq Act^*$ is realisable (by a weak CFM) iff L is closed under $\models$.

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Proof.

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Corollary

The finite set of MSCs $\{M_1, \ldots, M_n\}$ is realisable (by a weak CFM) iff $\bigcup_{i=1}^{n} Lin(M_i)$ is closed under $\models$.

Theorem

For any well-formed $L \subseteq Act^*$:

$$L \text{ is regular and closed under } \models$$
$$\text{if and only if}$$
$$L = Lin(\mathcal{A}) \text{ for some } \forall\text{-bounded weak CFM } \mathcal{A}.$$

Complexity of realisability

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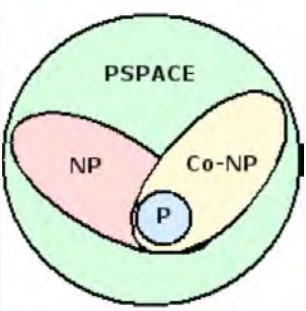

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Theorem: [Alur et al., 2001]

The decision problem "is a given finite set of MSCs realisable by a weak CFM?" is decidable and is co-NP complete.

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Proof.

1. Membership in co-NP is proven by showing that its complement is in NP. This is rather standard.
2. The co-NP hardness proof is based on a polynomial reduction of the join dependency problem to the above realisability problem. (Details on the black board.)

□