



**Advanced Model Checking
Summer term 2014**

– Series 5 –

Hand in on May 21'th before the exercise class.

Exercise 1

(7 points)

Let $CTL \setminus U$ be the sublogic of CTL that does not permit the until operator. Similarly, $CTL^* \setminus U$ means CTL^* without U. Which of the following statements are correct for finite transition systems.

- a) $CTL \setminus U$ equivalence is finer than $CTL \setminus O$ equivalence.
- b) $CTL \setminus U$ equivalence is finer than divergence-sensitive stutter trace equivalence.
- c) $CTL \setminus O$ equivalence is finer than $LTL \setminus O$ equivalence.
- d) Divergence-sensitive stutter bisimulation equivalence is finer than $CTL \setminus U$ equivalence.
- e) Stutter trace equivalence is finer than $CTL \setminus U$ equivalence.
- f) For AP-deterministic transition systems, stutter trace equivalence is finer than trace-equivalence.
- g) For AP-deterministic transition systems, trace equivalence is finer than $CTL^* \setminus U$ equivalence.

Exercise 2

(3 points)

Consider the transition systems T_1 and T_2 in Figure .

- (a) Show that $T_2 \preceq T_1$ by providing a simulation relation for (T_2, T_1) .
- (b) Show that $T_1 \not\preceq T_2$ by providing a $\forall CTL$ formula Φ_{\forall} and a $\exists CTL$ formula Φ_{\exists} such that $T_1 \models \Phi_{\forall}$, but $T_2 \not\models \Phi_{\forall}$ and $T_1 \models \Phi_{\exists}$, but $T_2 \not\models \Phi_{\exists}$.

