Concurrency Theory

Winter Semester 2015/16

Lecture 3: Hennessy-Milner Logic

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http://moves.rwth-aachen.de/teaching/ws-1516/ct/

Outline of Lecture 3

Recap: Calculus of Communicating Systems

Infinite State Spaces

Process Traces

Hennessy-Milner Logic

Closure under Negation

Syntax of CCS I

Definition (Syntax of CCS)

- Let A be a set of (action) names.
- $\overline{A} := \{\overline{a} \mid a \in A\}$ denotes the set of co-names.
- $Act := A \cup \overline{A} \cup \{\tau\}$ is the set of actions with the silent (or: unobservable) action τ.
- Let Pid be a set of process identifiers.
- The set Prc of process expressions is defined by the following syntax:

$$
\begin{array}{rlll}
P ::= & \mathrm{nil} & & \text{(inaction)} \\
\mid & \alpha.P & & \text{(prefixing)} \\
\mid & P_1 + P_2 & & \text{(choice)} \\
\mid & P_1 \parallel P_2 & & \text{(parallel composition)} \\
\mid & P \setminus L & & \text{(restriction)} \\
\mid & P[f] & & \text{(relabelling)} \\
\mid & C & & \text{(process call)}
\end{array}
$$

where $\alpha \in Act$, $L \subseteq A$, $C \in Pid$, and $f : Act \to Act$ such that $f(\tau) = \tau$ and $f(\overline{a}) = \overline{f(a)}$ for each $a \in A$.

Syntax of CCS II

Definition (continued)

- A (recursive) process definition is an equation system of the form
$$(C_i = P_i \mid 1 \leq i \leq k)$$
where $k \geq 1$, $C_i \in Pid$ (pairwise distinct), and $P_i \in Prc$ (with identifiers from $\{C_1, \ldots, C_k\}$).

Notational Conventions:

- $\overline{\overline{a}}$ means a
- $\sum_{i=1}^{n} P_i$ ($n \in \mathbb{N}$) means $P_1 + \ldots + P_n$ (where $\sum_{i=1}^{0} P_i := \mathrm{nil}$)
- $P \setminus a$ abbreviates $P \setminus \{a\}$
- $[a_1 \mapsto b_1, \ldots, a_n \mapsto b_n]$ stands for $f : Act \to Act$ with $f(a_i) = b_i$ ($i \in [n]$) and $f(\alpha) = \alpha$ otherwise
- restriction and relabelling bind stronger than prefixing, prefixing stronger than composition, composition stronger than choice:
$$P \setminus a + b.Q \parallel R \quad \text{means} \quad (P \setminus a) + ((b.Q) \parallel R)$$

Labelled Transition Systems

Goal: represent behaviour of system by (infinite) graph

- nodes = system states
- edges = transitions between states

Definition (Labelled transition system)

A (*Act*-)labelled transition system (LTS) is a triple $(S, Act, \longrightarrow)$ consisting of

- a set S of states
- a set Act of (action) labels
- a transition relation $\longrightarrow \subseteq S \times Act \times S$

For $(s, \alpha, s') \in \longrightarrow$ we write $s \xrightarrow{\alpha} s'$. An LTS is called finite if S is so.

Remarks:

- sometimes an initial state $s_0 \in S$ is distinguished ("$LTS(s_0)$")
- (finite) LTSs correspond to (finite) automata without final states

Semantics of CCS I

Definition (Semantics of CCS)

A process definition $(C_i = P_i \mid 1 \leq i \leq k)$ determines the LTS $(Prc, Act, \longrightarrow)$ whose transitions can be inferred from the following rules ($P, P', Q, Q' \in Prc$, $\alpha \in Act$, $\lambda \in A \cup \overline{A}$, $a \in A$):

$$\text{(Act)}\frac{}{\alpha.P \xrightarrow{\alpha} P} \qquad \text{(Sum}_1\text{)}\frac{P \xrightarrow{\alpha} P'}{P + Q \xrightarrow{\alpha} P'} \qquad \text{(Sum}_2\text{)}\frac{Q \xrightarrow{\alpha} Q'}{P + Q \xrightarrow{\alpha} Q'}$$

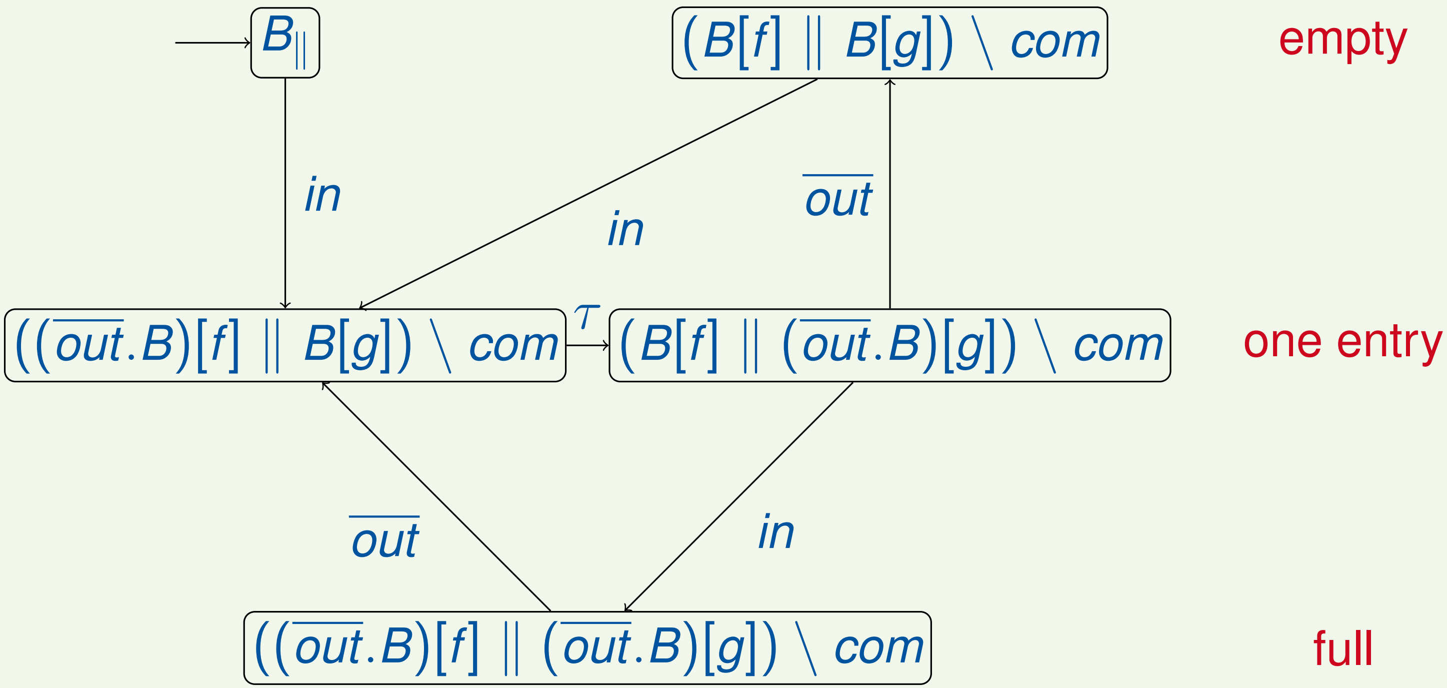
$$\text{(Par}_1\text{)}\frac{P \xrightarrow{\alpha} P'}{P \parallel Q \xrightarrow{\alpha} P' \parallel Q} \qquad \text{(Par}_2\text{)}\frac{Q \xrightarrow{\alpha} Q'}{P \parallel Q \xrightarrow{\alpha} P \parallel Q'} \qquad \text{(Com)}\frac{P \xrightarrow{\lambda} P' \quad Q \xrightarrow{\overline{\lambda}} Q'}{P \parallel Q \xrightarrow{\tau} P' \parallel Q'}$$

$$\text{(Res)}\frac{P \xrightarrow{\alpha} P' \quad (\alpha, \overline{\alpha} \notin L)}{P \setminus L \xrightarrow{\alpha} P' \setminus L} \qquad \text{(Rel)}\frac{P \xrightarrow{\alpha} P'}{P[f] \xrightarrow{f(\alpha)} P'[f]} \qquad \text{(Call)}\frac{P \xrightarrow{\alpha} P' \quad (C = P)}{C \xrightarrow{\alpha} P'}$$

Semantics of CCS II

Example

Parallel two-place buffer: $B_{\|} = (B[f] \parallel B[g]) \setminus com$, $B = in.\overline{out}.B$
where $f := [out \mapsto com]$ and $g := [in \mapsto com]$



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gives rise to infinite LTS (abbreviating $down := down.\mathrm{nil}$):

$$\begin{tikzcd}
{} \arrow[r] & \boxed{C} \arrow[r, "up"] & \boxed{C \parallel down} \arrow[r, "up"] \arrow[d, "down"] & \boxed{C \parallel down \parallel down} \arrow[r, "up"] \arrow[d, "down"] & \boxed{\cdots} \\
 & & \boxed{C \parallel \mathrm{nil}} \arrow[r, "up"] & \boxed{C \parallel down \parallel \mathrm{nil}} \arrow[r, "up"] \arrow[d, "down"] & \boxed{\cdots} \\
 & & & \boxed{C \parallel \mathrm{nil} \parallel \mathrm{nil}} \arrow[r, "up"] & \boxed{\cdots}
\end{tikzcd}$$

The Power of Recursive Definitions

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& & & C \parallel \mathrm{nil} \parallel \mathrm{nil} \arrow[r, "up"] & \cdots
\end{tikzcd}$$

Sequential "specification": $C_0 = up.C_1$

$$C_n = up.C_{n+1} + down.C_{n-1} \qquad (n > 0)$$

Outline of Lecture 3

Process Traces I

Goal: reduce processes to the action sequences they can perform

Definition 3.2 (Trace language)

For every $P \in Prc$, let

$$Tr(P) := \{w \in Act^* \mid \text{ex. } P' \in Prc \text{ such that } P \xrightarrow{w} P'\}$$

be the trace language of P

(where $\xrightarrow{w} := \xrightarrow{a_1} \circ \ldots \circ \xrightarrow{a_n}$ for $w = a_1 \ldots a_n$).

$P, Q \in Prc$ are called trace equivalent if $Tr(P) = Tr(Q)$.

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Example 3.3 (One-place buffer)

$B = in.\overline{out}.B$

$\implies \quad Tr(B) = (in \cdot \overline{out})^* \cdot (in + \varepsilon)$

Process Traces II

Remarks:

- The trace language of $P \in Prc$ is accepted by the LTS of P, interpreted as a (finite or infinite) automaton with initial state P and where every state is final.

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- Trace equivalence is obviously an equivalence relation (i.e., reflexive, symmetric, and transitive).
- Trace equivalence identifies processes with identical LTSs: the trace language of a process consists of the (finite) paths in the LTS. Thus:

$$LTS(P) = LTS(Q) \implies Tr(P) = Tr(Q)$$

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- Later we will see: trace equivalence is too coarse, i.e., identifies too many processes $\implies$ bisimulation

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Motivation

Goal: check processes for simple properties

- action a is initially enabled
- action b is initially disabled
- a deadlock never occurs
- always sends a reply after receiving a request

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- action *a* is initially enabled
- action *b* is initially disabled
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- formalisation in Hennessy-Milner Logic (HML)
- M. Hennessy, R. Milner: *On Observing Nondeterminism and Concurrency*, ICALP 1980, Springer LNCS 85, 299–309
- checking by exploration of state space

Syntax of HML

Definition 3.4 (Syntax of HML)

The set *HMF* of Hennessy-Milner formulae over a set of actions *Act* is defined by the following syntax:

$$
\begin{array}{rcll}
F & ::= & \mathrm{tt} & \text{(true)} \\
 & \mid & \mathrm{ff} & \text{(false)} \\
 & \mid & F_1 \wedge F_2 & \text{(conjunction)} \\
 & \mid & F_1 \vee F_2 & \text{(disjunction)} \\
 & \mid & \langle\alpha\rangle F & \text{(diamond)} \\
 & \mid & [\alpha]F & \text{(box)}
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where $\alpha \in Act$.

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$$

where $\alpha \in Act$.

Abbreviations for $L = \{\alpha_1, \ldots, \alpha_n\}$ ($n \in \mathbb{N}$):

- $\langle L \rangle F := \langle \alpha_1 \rangle F \vee \ldots \vee \langle \alpha_n \rangle F$
- $[L]F := [\alpha_1]F \wedge \ldots \wedge [\alpha_n]F$
- In particular, $\langle \emptyset \rangle F :=$ ff and $[\emptyset]F :=$ tt

Meaning of HML Constructs

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- A process satisfies $[\alpha]F$ for some $\alpha \in Act$ iff all its α-labelled transitions lead to a state satisfying F (necessity).

Semantics of HML

Definition 3.5 (Semantics of HML)

Let $(S, Act, \longrightarrow)$ be an LTS and $F \in HMF$. The set of processes in S that satisfy F, $[\![F]\!] \subseteq S$, is defined by:

$$[\![\mathrm{tt}]\!] := S \qquad [\![\mathrm{ff}]\!] := \emptyset$$
$$[\![F_1 \wedge F_2]\!] := [\![F_1]\!] \cap [\![F_2]\!] \qquad [\![F_1 \vee F_2]\!] := [\![F_1]\!] \cup [\![F_2]\!]$$
$$[\![\langle\alpha\rangle F]\!] := \langle\cdot\alpha\cdot\rangle([\![F]\!]) \qquad [\![[\alpha]F]\!] := [\cdot\alpha\cdot]([\![F]\!])$$

where $\langle\cdot\alpha\cdot\rangle, [\cdot\alpha\cdot] : 2^S \to 2^S$ are given by

$$\langle\cdot\alpha\cdot\rangle(T) := \{s \in S \mid \exists s' \in T : s \xrightarrow{\alpha} s'\}$$
$$[\cdot\alpha\cdot](T) := \{s \in S \mid \forall s' \in S : s \xrightarrow{\alpha} s' \implies s' \in T\}$$

We write $s \models F$ iff $s \in [\![F]\!]$. Two HML formulae are equivalent (written $F \equiv G$) iff they are satisfied by the same processes in every LTS.

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Example 3.6 ($\langle\cdot\alpha\cdot\rangle$, $[\cdot\alpha\cdot]$ operators)

on the board

Simple Properties Revisited

Example 3.7

1. action a is initially enabled: $\langle a \rangle \mathrm{tt}$

$$\begin{aligned} [\![\langle a \rangle \mathrm{tt}]\!] &= \langle \cdot a \cdot \rangle [\![\mathrm{tt}]\!] = \langle \cdot a \cdot \rangle (S) \\ &= \{s \in S \mid \exists s' \in S : s \xrightarrow{a} s'\} =: \{s \in S \mid s \xrightarrow{a}\} \end{aligned}$$

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   - initially: $\langle Act \rangle \mathrm{tt}$
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4. responsiveness:
   - initially: $[request]\langle\overline{reply}\rangle\mathrm{tt}$
   - always: later (requires recursion)

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Closure under Negation

Observation: negation is *not* one of the HML constructs

Reason: HML is closed under negation

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Lemma 3.8

For every $F \in HMF$ *there exists* $F^c \in HMF$ *such that* $[\![F^c]\!] = S \setminus [\![F]\!]$ *for every LTS* $(S, Act, \longrightarrow)$.

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Proof.

Definition of F^c:

$$\begin{aligned} \mathrm{tt}^c &:= \mathrm{ff} & \mathrm{ff}^c &:= \mathrm{tt} \\ (F_1 \wedge F_2)^c &:= F_1^c \vee F_2^c & (F_1 \vee F_2)^c &:= F_1^c \wedge F_2^c \\ (\langle\alpha\rangle F)^c &:= [\alpha]F^c & ([\alpha]F)^c &:= \langle\alpha\rangle F^c \end{aligned}$$

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$[\![F^c]\!] = S \setminus [\![F]\!]$: on the board □