Concurrency Theory

Winter Semester 2015/16

Lecture 1: Introduction

Joost-Pieter Katoen and Thomas Noll
Software Modeling and Verification Group
RWTH Aachen University

http://moves.rwth-aachen.de/teaching/ws-1516/ct/

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Outline of Lecture 1

People

- Lectures:
  - Joost-Pieter Katoen (`katoen@cs.rwth-aachen.de`)
  - Thomas Noll (`noll@cs.rwth-aachen.de`)
- Exercise classes:
  - Benjamin Kaminski (`benjamin.kaminski@cs.rwth-aachen.de`)
  - Christoph Matheja (`matheja@cs.rwth-aachen.de`)
- Student assistant: Wanted!
  - Evaluation of exercises
  - Organizational support
  - 12 hrs/week contract
  - Previous CT lecture not a prerequisite (but of course helpful)

Preliminaries

Target Audience

- Master program Informatik
  - Theoretische Informatik
- Master program Software Systems Engineering
  - Theoretical Foundations of SSE

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- Master program Informatik
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- Master program Software Systems Engineering
  - Theoretical Foundations of SSE
- In general:
  - interest in formal models for concurrent (software) systems
  - application of mathematical modelling and reasoning methods
- Expected: basic knowledge in
  - essential concepts of operating systems and system software
  - formal languages and automata theory
  - mathematical logic

Course Objectives

Objectives

- Understand the foundations of concurrent systems
- Model (and compare) concurrent systems in a rigorous manner
- Understand the main semantical underpinnings of concurrency

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- Supporting the design phase
  - “Programming Concurrent Systems”
  - synchronization, scheduling, semaphores, ...

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- Verifying functional correctness properties
  - "Model Checking"
  - validation of mutual exclusion, fairness, absence of deadlocks, ...

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  - synchronization, scheduling, semaphores, ...
- Verifying functional correctness properties
  - “Model Checking”
  - validation of mutual exclusion, fairness, absence of deadlocks, ...
- Comparing expressivity of models of concurrency
  - “interleaving” vs. “true concurrency”
  - equivalence, refinement, abstraction, ...

Organization

- Schedule:
  - Lecture Mon 14:15–15:45 AH 1 (starting 19 Oct)
  - Lecture Thu 14:15–15:45 AH 2 (starting 12 Nov)
  - Exercise class Mon 10:15–11:45 AH 6 (starting 26 Oct with “0th exercise”)
- Irregular lecture dates – checkout web page!

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- 1st assignment sheet: next Monday (26 Oct) on web page
  - submission by 2 Nov before exercise class
  - presentation on 2 Nov
- Work on assignments in groups of three

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- Work on assignments in groups of three
- Examination (6 ECTS credits):
  - oral or written (depending on number of participants)
  - date to be fixed
- Admission requires at least 50% of the points in the exercises
- Solutions to exercises and exam in English or German

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Concurrency and Interaction by Example

Observation: concurrency introduces new phenomena

Example 1.1

$$x := 0;$$
$$(x := x + 1 \parallel x := x + 2)$$

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$$x := 0;$$
$$(x := x + 1 \parallel x := \underset{2}{x + 2}) \qquad \text{value of } x\text{: } 2$$

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- But: both parallel components could read x before it is written
- Thus: x is assigned 2,

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- Thus: x is assigned 2, 1,

Concurrency and Interaction by Example

Observation: concurrency introduces new phenomena

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$$x := 0;$$
$$(x := x + 1 \parallel x := \underset{2}{x + 2})$$

value of x: 0

- At first glance: x is assigned 3
- But: both parallel components could read x before it is written
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$$2$$

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Concurrency and Interaction by Example

Observation: concurrency introduces new phenomena

Example 1.1

$$\begin{array}{c} x := 0; \\ (\underbrace{x := x + 1}_{3} \parallel x := x + 2) \end{array} \qquad \text{value of } x\text{: } 2$$

- At first glance: x is assigned 3
- But: both parallel components could read x before it is written
- Thus: x is assigned 2, 1,

Concurrency and Interaction by Example

Observation: concurrency introduces new phenomena

Example 1.1

$$x := 0;$$
$$(\underset{3}{x := x + 1} \parallel x := x + 2) \qquad \text{value of } x\text{: } 3$$

- At first glance: x is assigned 3
- But: both parallel components could read x before it is written
- Thus: x is assigned 2, 1, or 3

Concurrency and Interaction by Example

Observation: concurrency introduces new phenomena

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$$x := 0;$$
$$(x := x + 1 \parallel x := x + 2)$$

- At first glance: x is assigned 3
- But: both parallel components could read x before it is written
- Thus: x is assigned 2, 1, or 3
- If exclusive access to shared memory and atomic execution of assignments guaranteed
  $\Longrightarrow$ only possible outcome: 3

Concurrency and Interaction

The problem arises due to the combination of

- concurrency and
- interaction (here: via shared memory)

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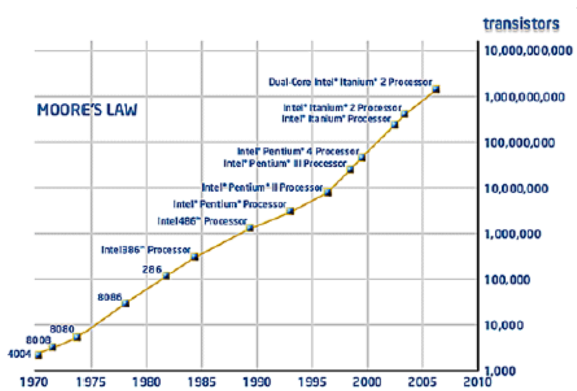
Conclusion

When modelling concurrent systems, the precise description of the mechanisms of both concurrency and interaction is crucially important.

Concurrency Everywhere

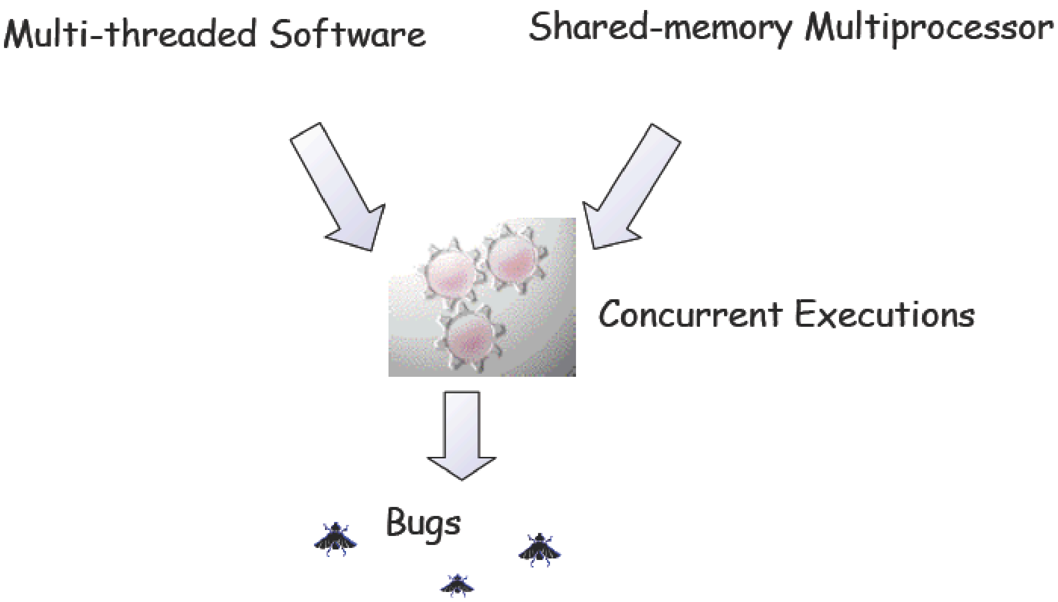
- Operating systems
- Embedded/reactive systems:
  - parallelism (at least) between hardware, software, and environment
- High-end parallel hardware infrastructure
  - high-performance computing
- Low-end parallel hardware infrastructure:
  - increasing performance only achievable by parallelism
  - multi-core computers, GPGPUs, FPGAs



Moore's Law: Transistor density doubles every 2 years

Problems Everywhere

- Operating systems:
  - mutual exclusion
  - fairness
  - no deadlocks, ...
- Shared-memory systems:
  - memory models
  - inconsistencies ("sequential consistency" vs. relaxed notions)
- Embedded systems:
  - safety
  - liveness, ...



Outline of Lecture 1

Preliminaries

Concurrency and Interaction

A Closer Look at Memory Models

A Closer Look at Reactive Systems

Overview of the Course

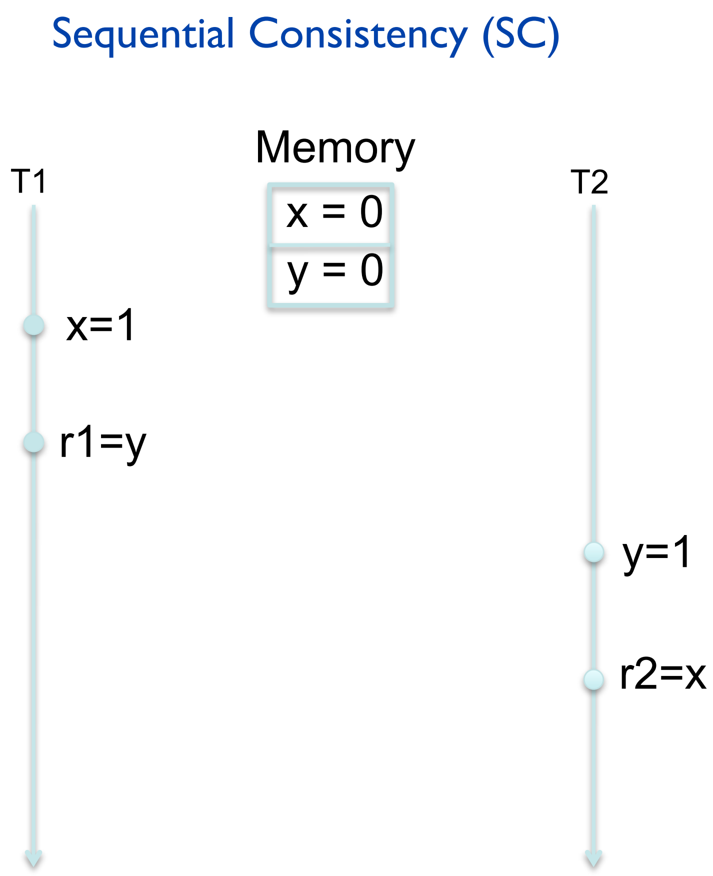
Memory Models

An illustrative example

Initially: x = y = 0

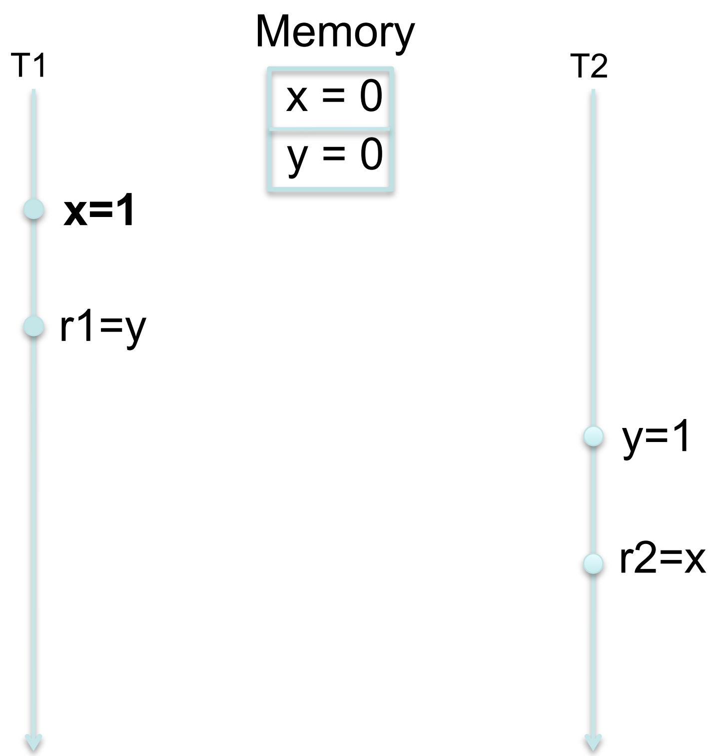
thread1:	thread2:
1: x = 1	3: y = 1
2: r1 = y	4: r2 = x

Memory Models



Memory Models

Sequential Consistency (SC)



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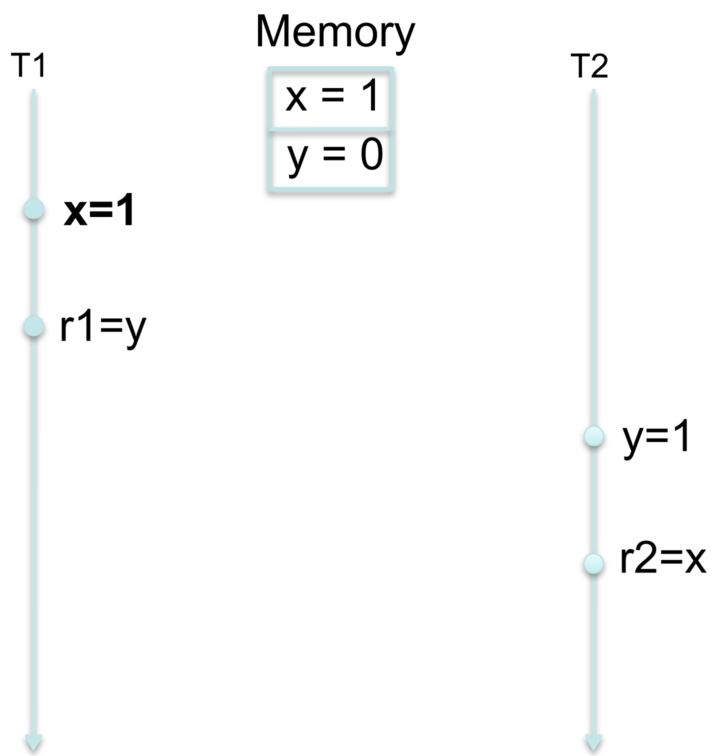

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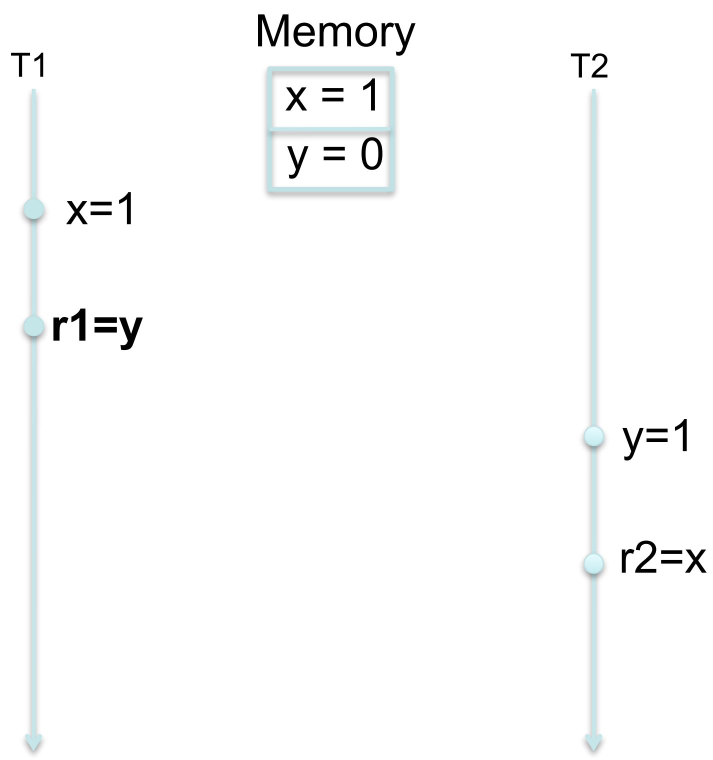

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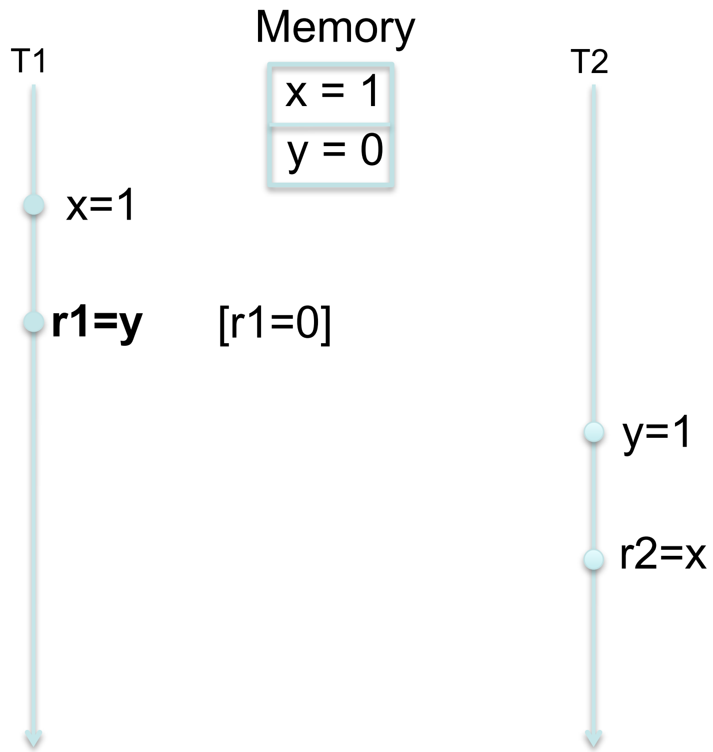

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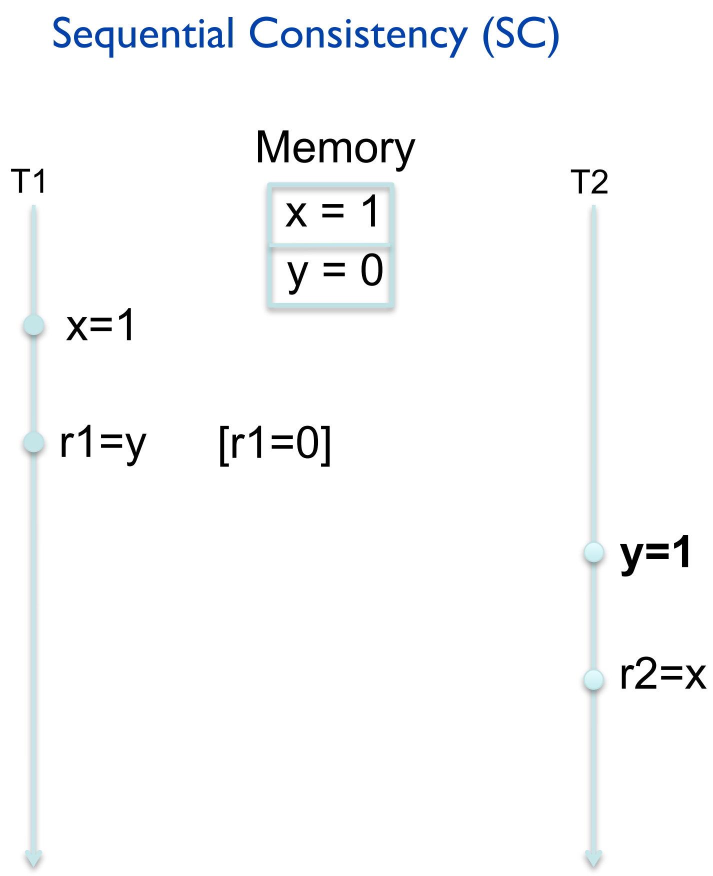

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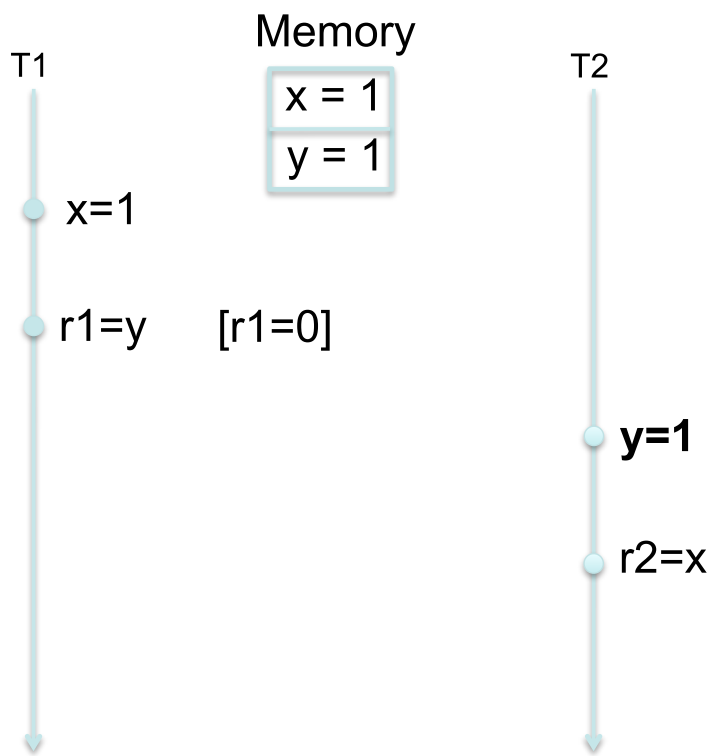

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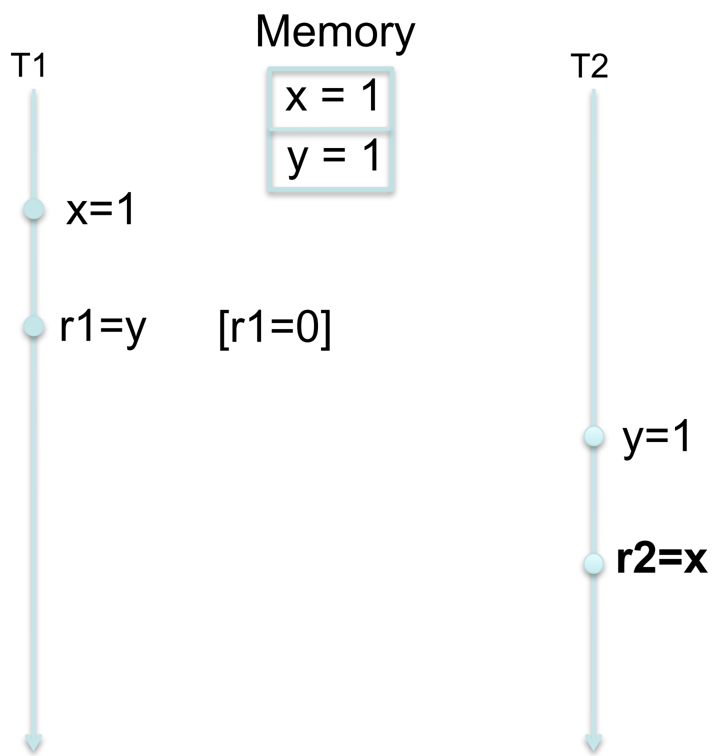

A Closer Look at Memory Models

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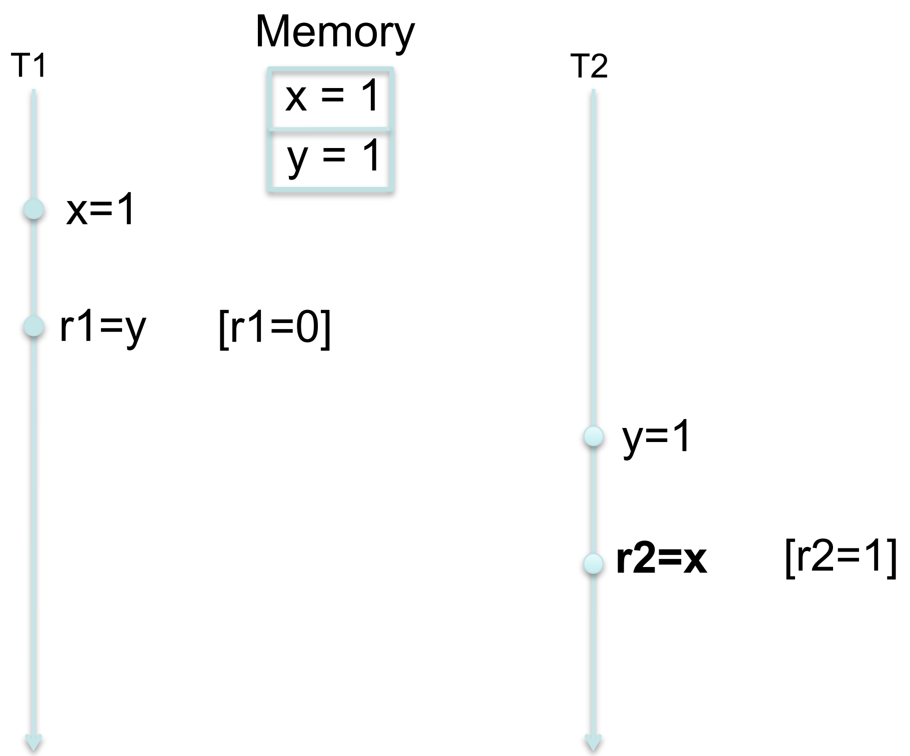

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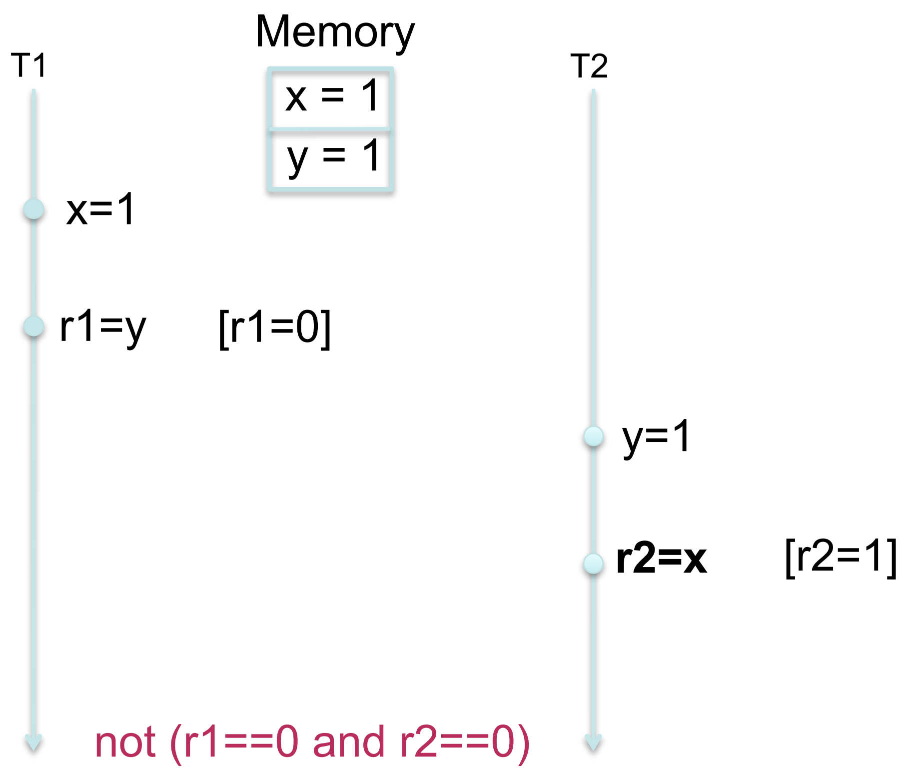
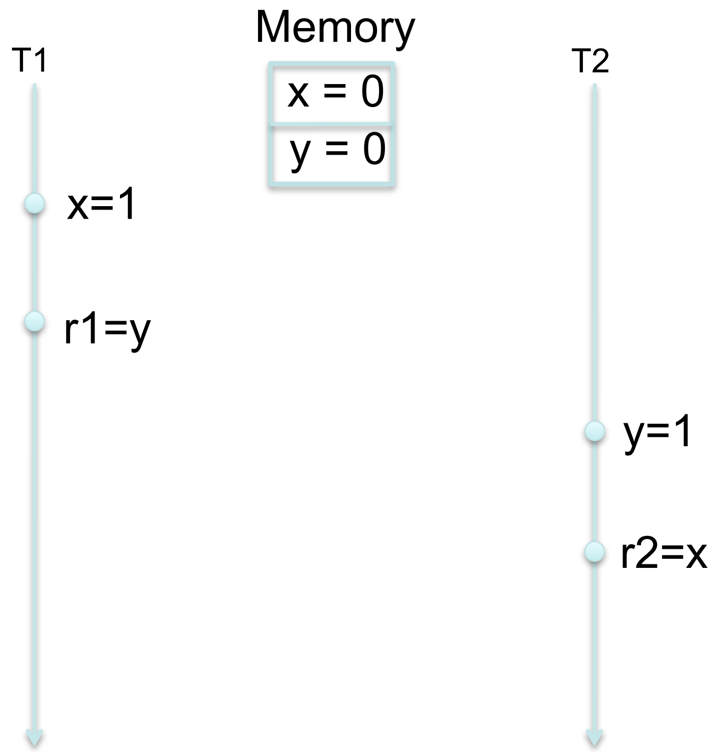


Memory Models

Total Store Ordering (TSO)



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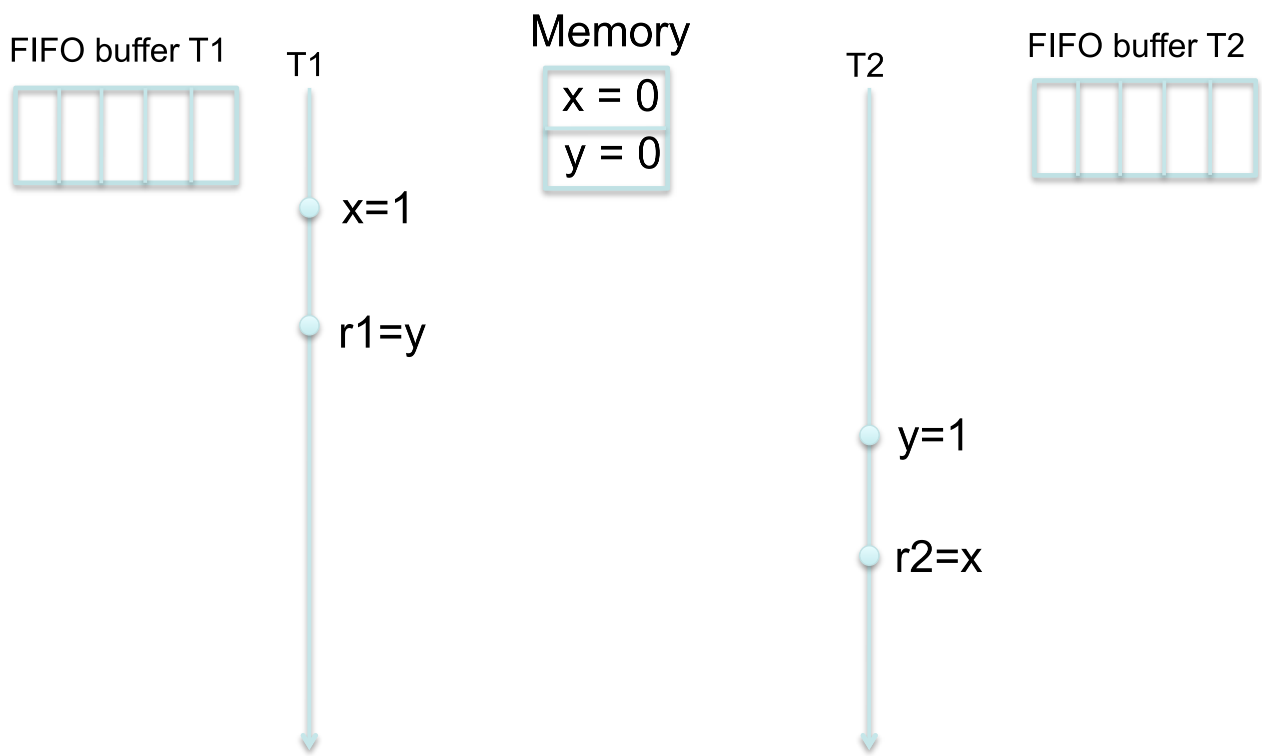

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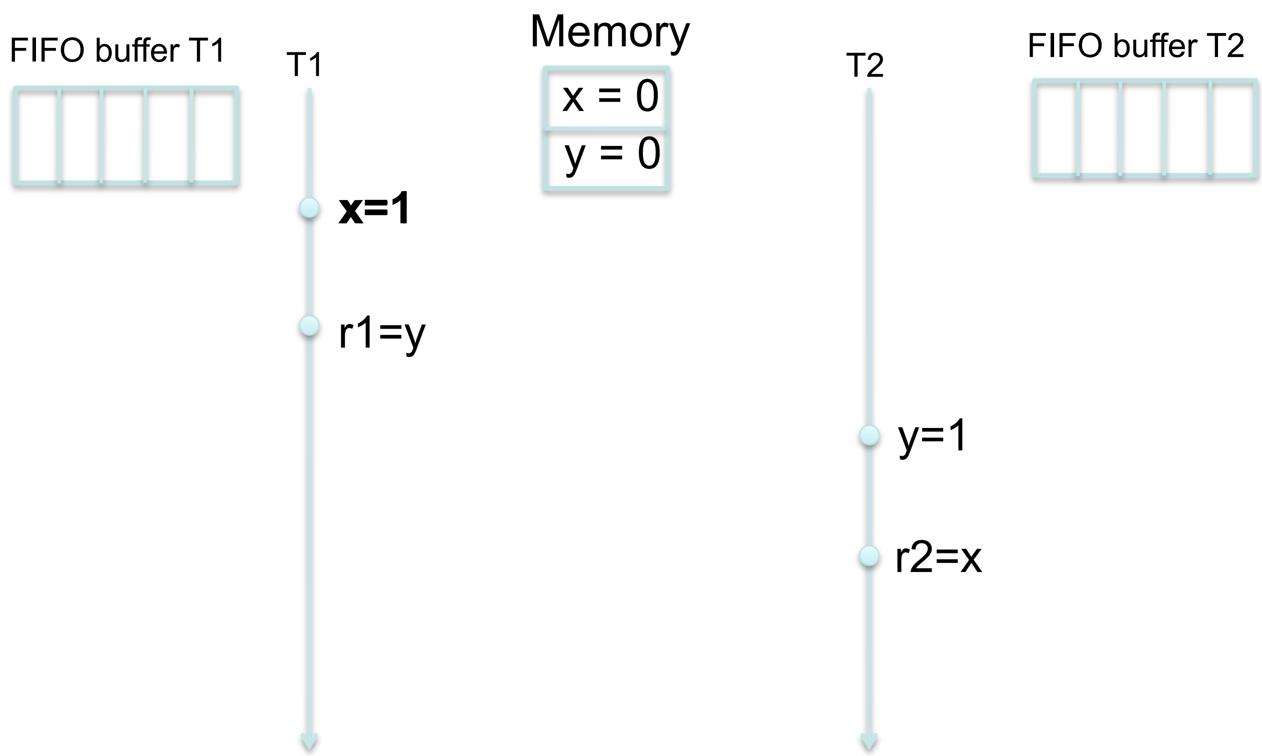

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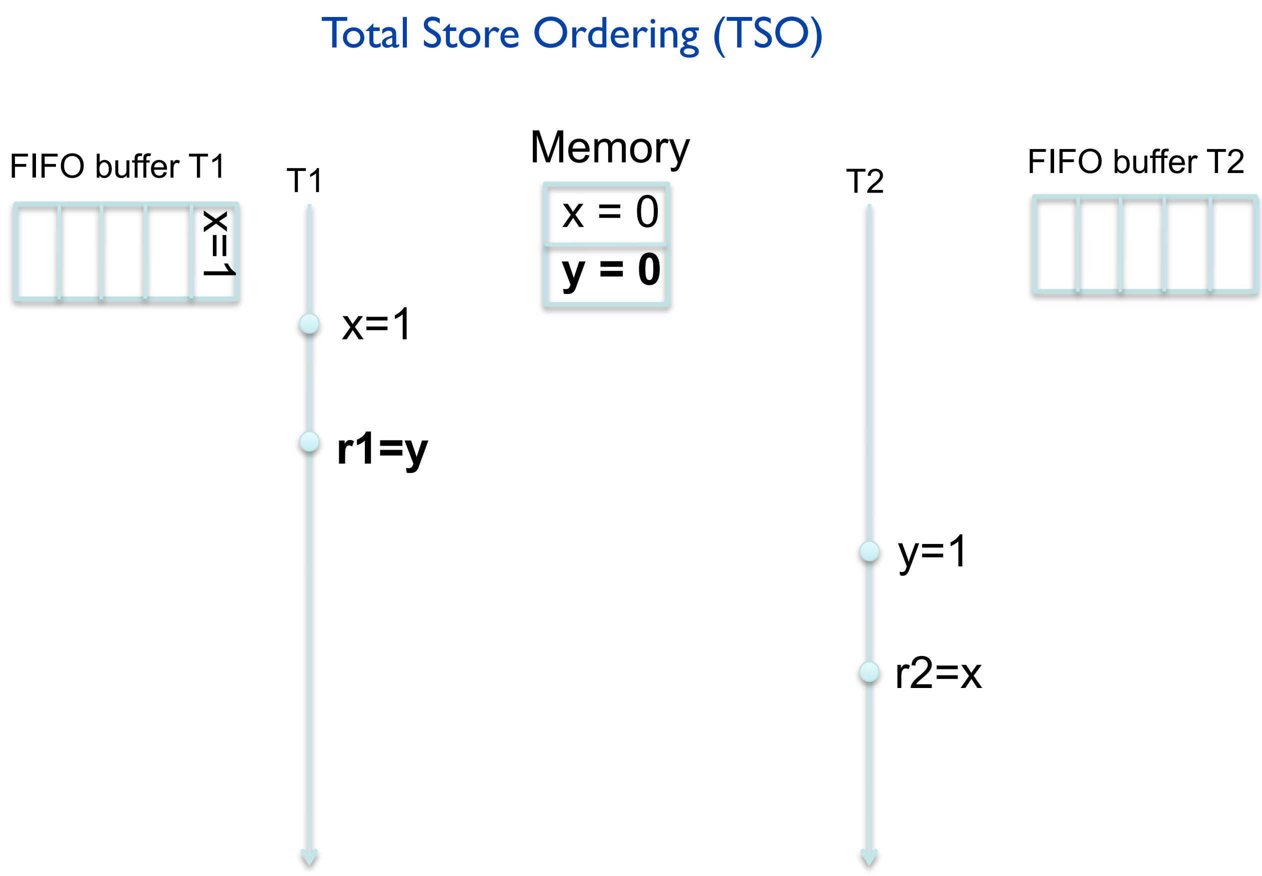

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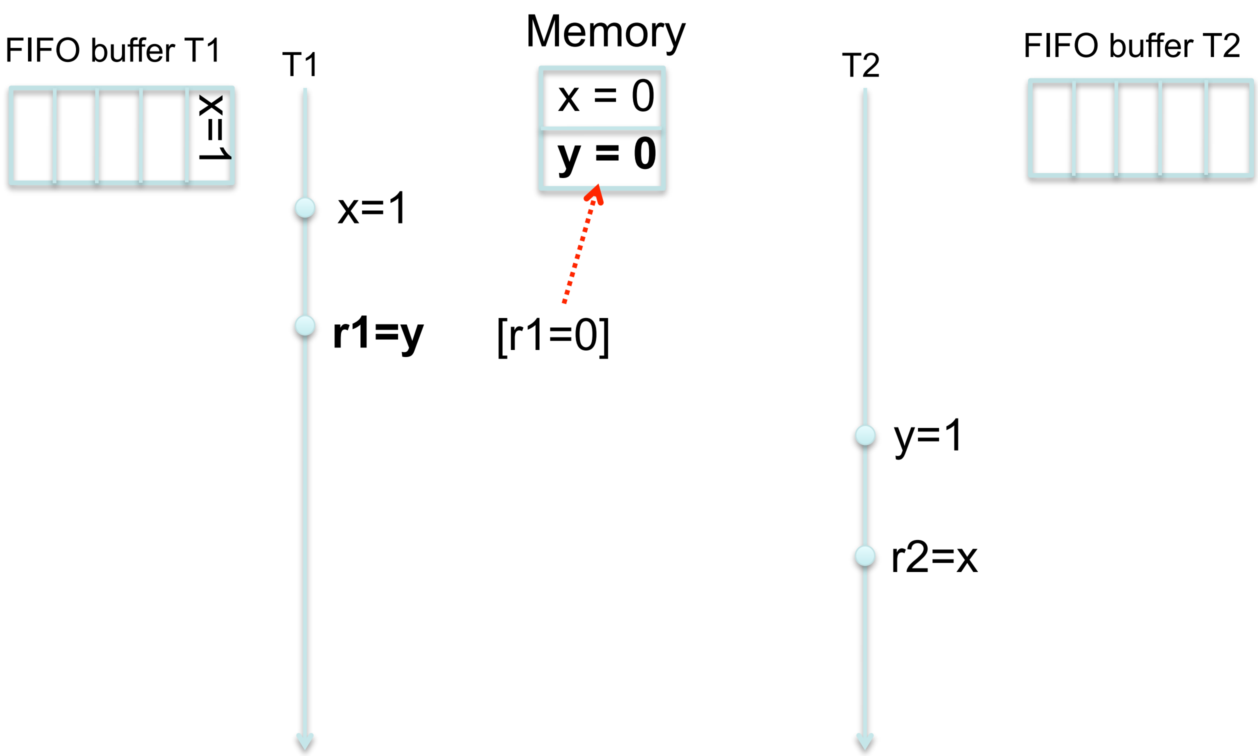

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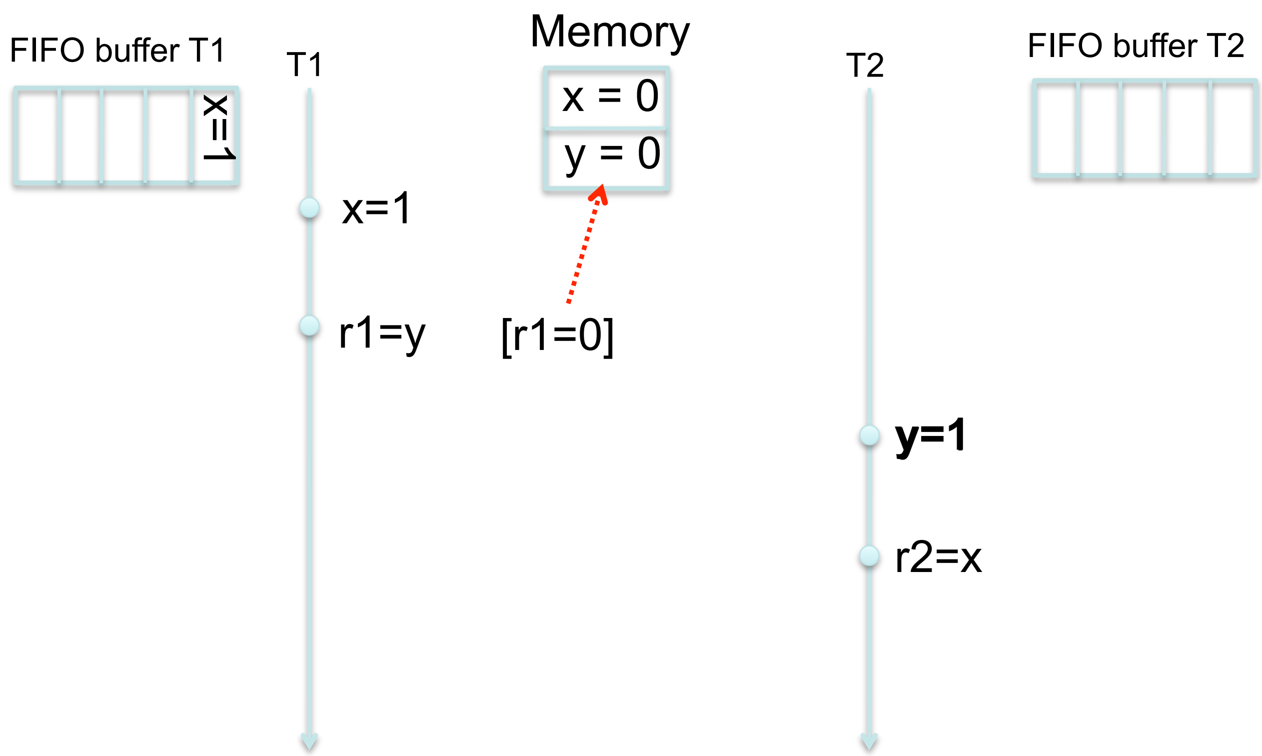

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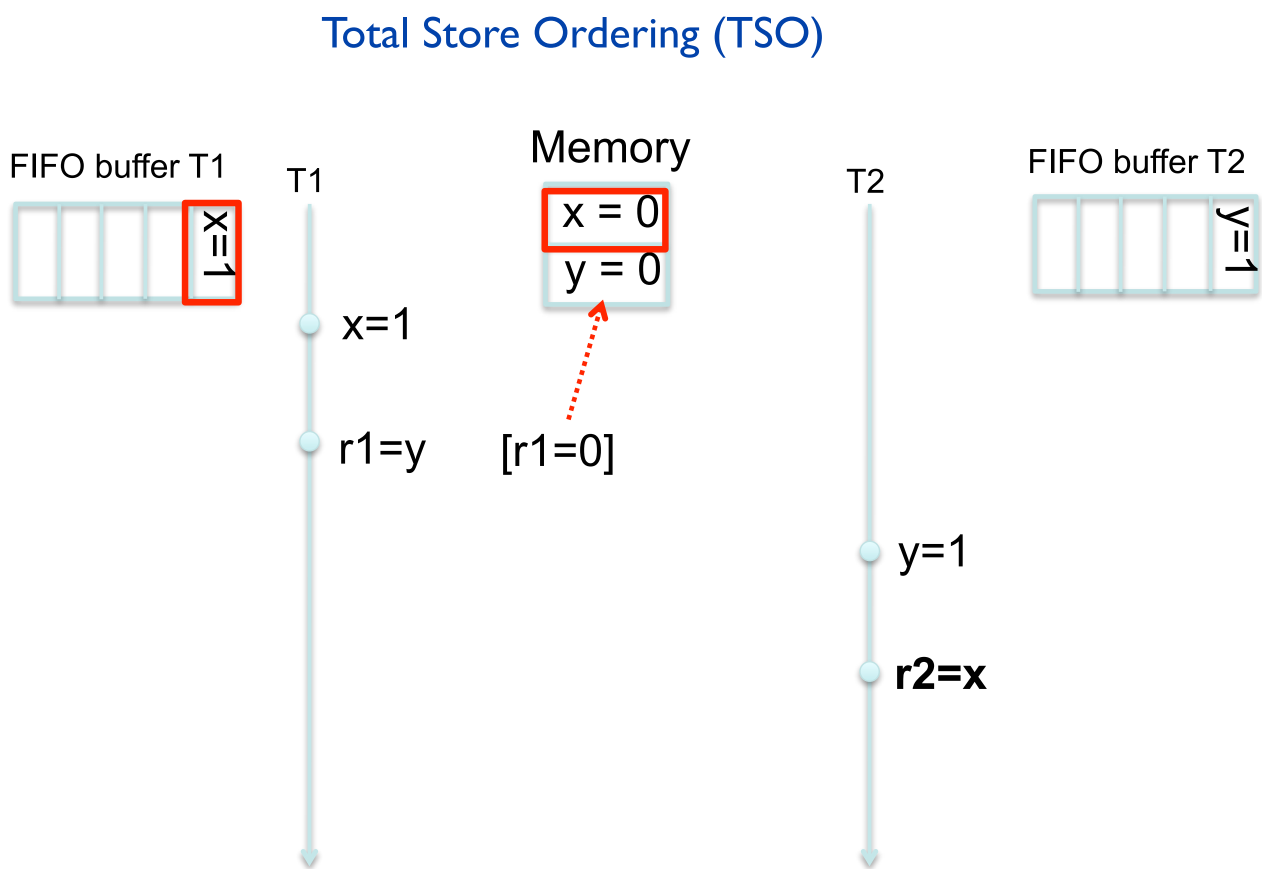

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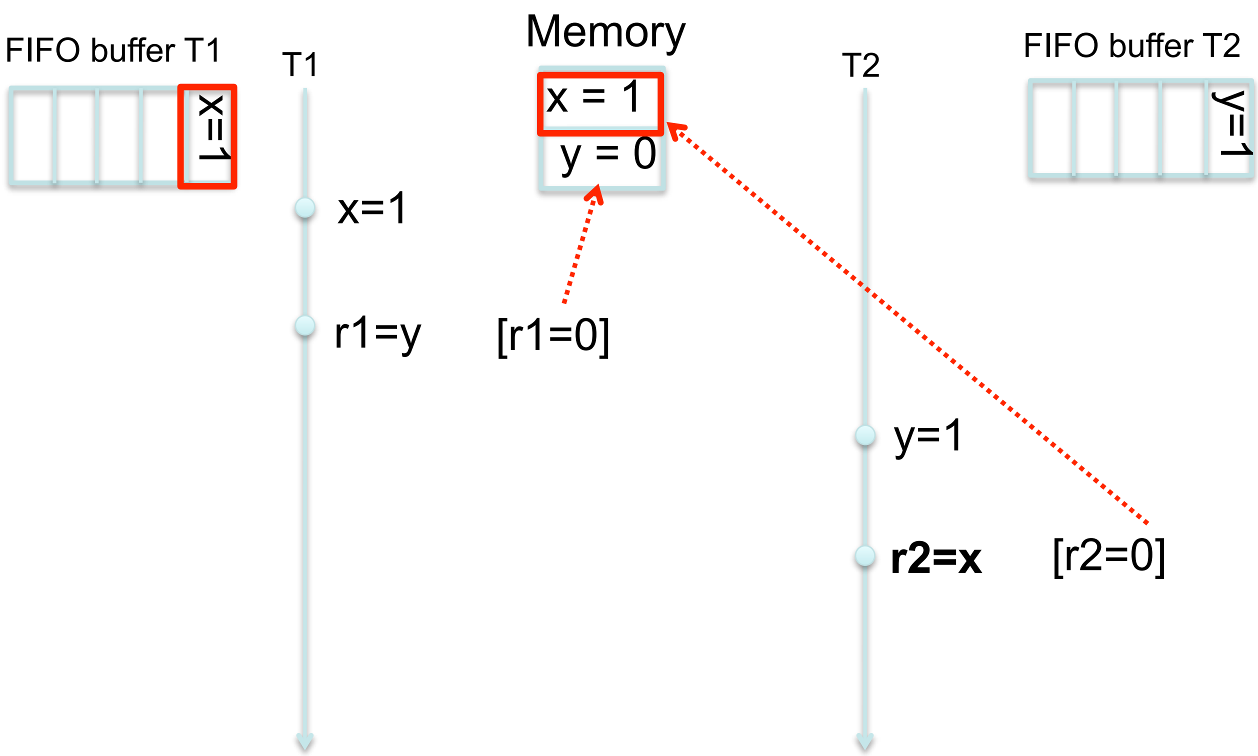

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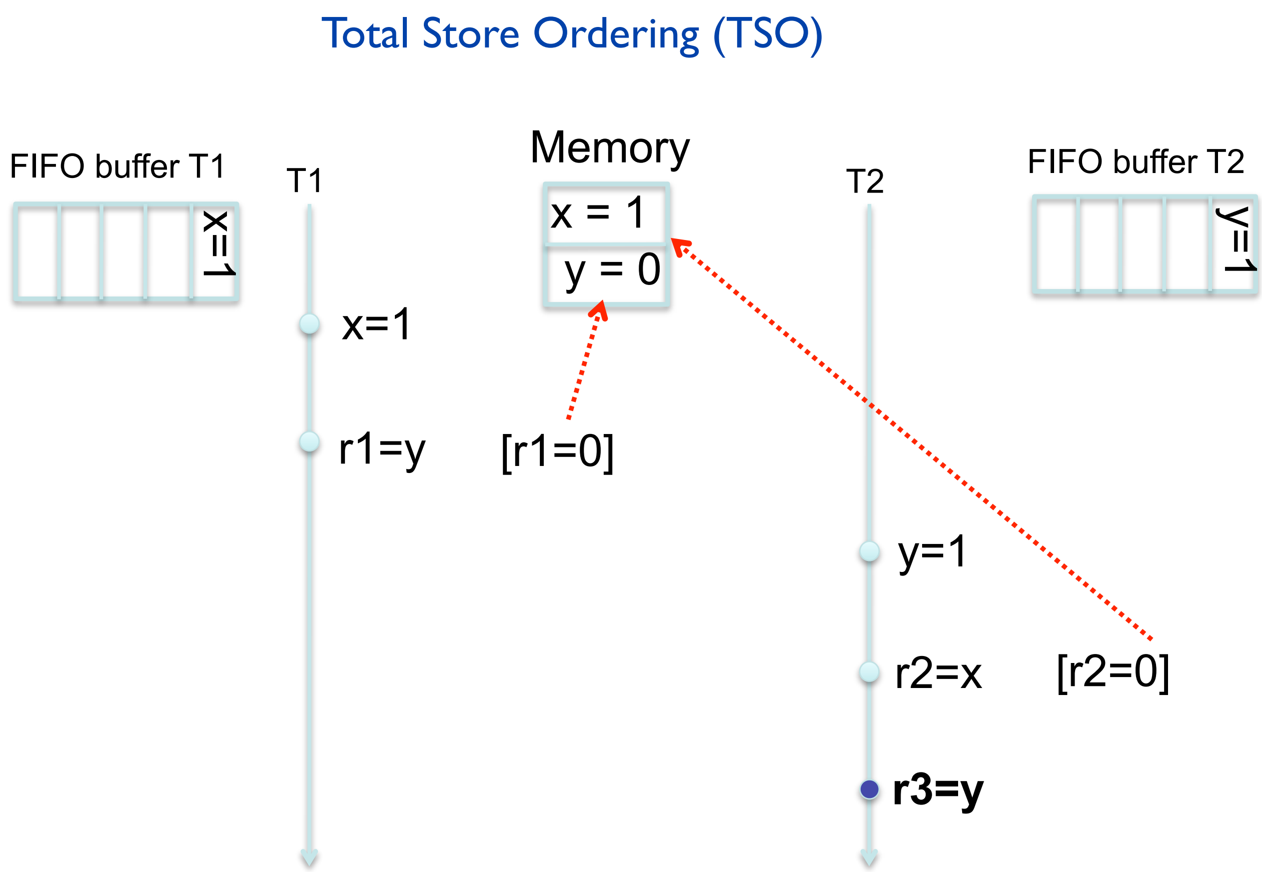

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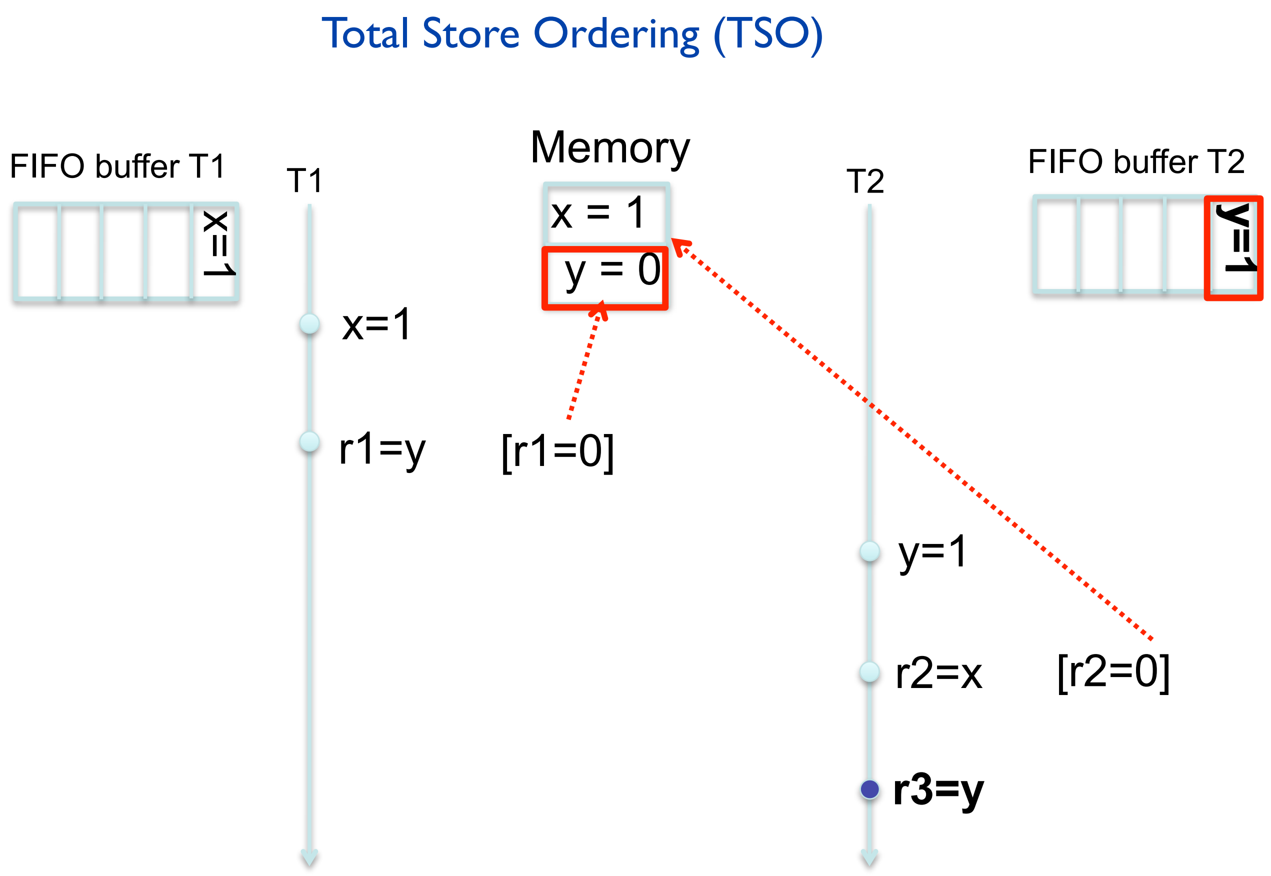

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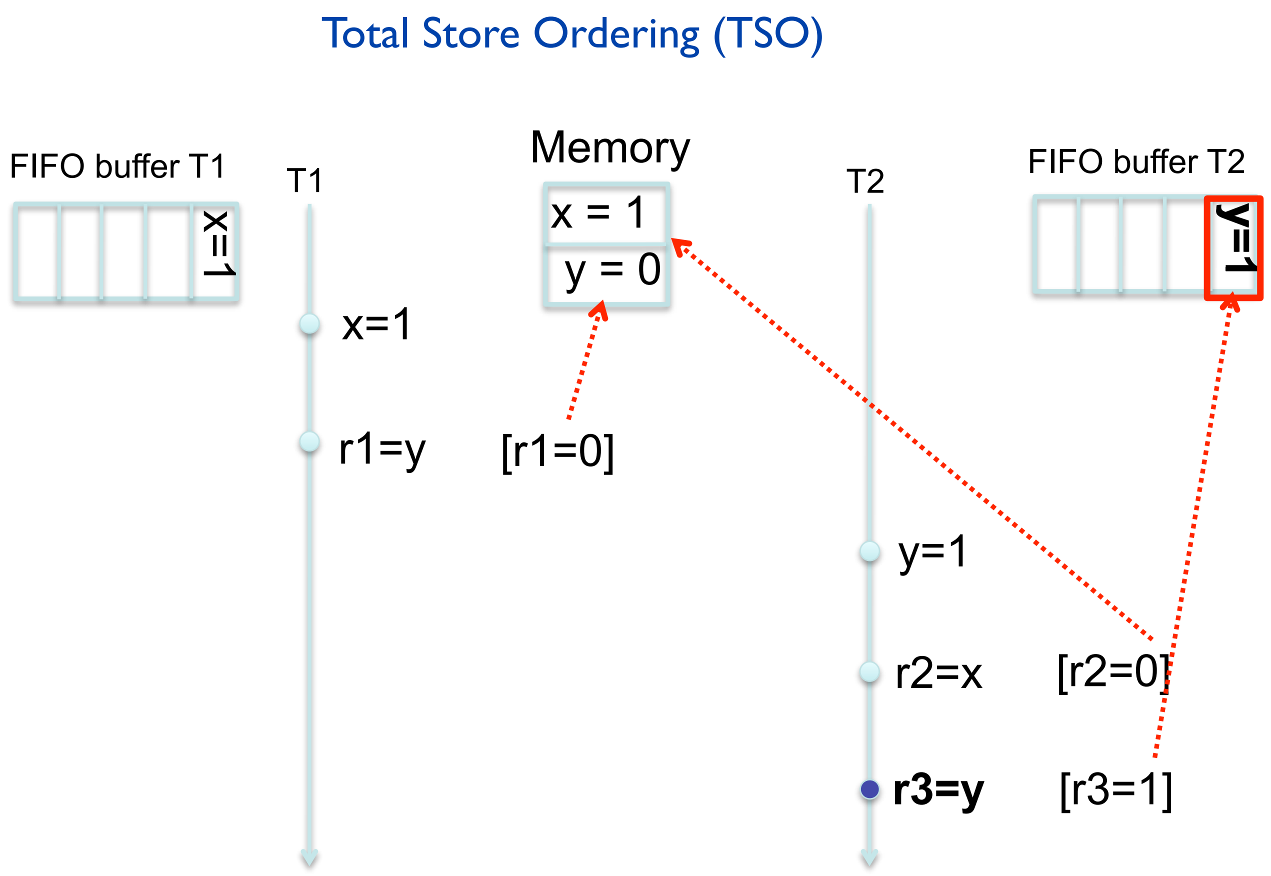

Outline of Lecture 1

A Closer Look at Reactive Systems

Reactive Systems I

- Thus: "classical" model for sequential systems

$$\mathit{System} : \mathit{Input} \rightarrow \mathit{Output}$$

  (transformational systems) is not adequate
- Missing: aspect of interaction

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- Rather: reactive systems which interact with environment and among themselves

Reactive Systems I

- Thus: "classical" model for sequential systems

$$\textit{System} : \textit{Input} \to \textit{Output}$$

  (transformational systems) is not adequate
- Missing: aspect of interaction
- Rather: reactive systems which interact with environment and among themselves
- Main interest: not terminating computations but infinite behavior
  (system maintains ongoing interaction with environment)
- Examples:
  - embedded systems controlling mechanical or electrical devices
    (planes, cars, home appliances, ...)
  - power plants, production lines, ...

Reactive Systems II

Observation: reactive systems often safety critical
$\Longrightarrow$ correct behavior has to be ensured

- Safety properties: “Nothing bad is going to happen.”
  E.g., “at most one process in the critical section”
- Liveness properties: “Eventually something good will happen.”
  E.g., “every request will finally be answered by the server”
- Fairness properties: “No component will starve to death.”
  E.g., “any process requiring entry to the critical section will eventually be admitted”

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Preliminaries

Concurrency and Interaction

A Closer Look at Memory Models

A Closer Look at Reactive Systems

Overview of the Course

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1. Introduction and Motivation
2. The “Interleaving” Approach
   – Syntax and semantics of CCS
   – Hennessy-Milner Logic
   – Case study: mutual exclusion
   – Extensions and alternative approaches (value passing, mobility, CSP, ACP, ...)
3. Equivalence, Refinement and Compositionality
   – Behavioural equivalences ((bi-)simulation)
   – Case study: mutual exclusion
   – (Pre-)congruences and compositional abstraction
   – HML and bisimilarity
4. The “True Concurrency” Approach
   – Petri nets: basic concepts
   – Case study: mutual exclusion
   – Branching processes and net unfoldings
   – Analyzing Petri nets
   – Alternative models (trace languages, event structures, ...)
5. Extensions (timed models, ...)

Literature

(also see the collection "Handapparat Softwaremodellierung und Verifikation" at the CS Library)

- Fundamental:
  - Luca Aceto, Anna Ingólfsdóttir, Kim Guldstrand Larsen and Jiří Srba: *Reactive Systems: Modelling, Specification and Verification*. Cambridge University Press, 2007.
  - Wolfgang Reisig: *Understanding Petri Nets: Modeling Techniques, Analysis Methods, Case Studies*. Springer Verlag, 2012.
- Supplementary:
  - Maurice Herlihy and Nir Shavit: *The Art of Multiprocessor Programming*. Elsevier, 2008.
  - Jan Bergstra, Alban Ponse and Scott Smolka (Eds.): *Handbook of Process Algebra*. Elsevier, 2001.