

Concurrency Theory WS 2015/2016 — Series 6 —

Hand in until December 14th before the exercise class.

Exercise 1 (Structural Congruence)

Prove that $P \rightarrow Q$ implies that there exists a derivation of this reduction in which the (Struct) rule (see Definition 9.6) is applied, if at all, only as the last rule.

Exercise 2 (Reaction Relation)

Let

$$\begin{split} S &= \operatorname{new} x(\\ & (x(u) \,.\, u(y) \,.\, u(z) \,.\, \bar{y} \langle z \rangle \,.\, \operatorname{nil} \\ & || \,\, x(t) \,.\, t(w) \,.\, t(v) \,.\, \bar{v} \langle w \rangle \,.\, \operatorname{nil}) \\ & || \,\, \operatorname{lnew} s(\bar{x} \langle s \rangle \,.\, \bar{s} \langle a \rangle \,.\, \bar{s} \langle b \rangle \,.\, \operatorname{nil}) \\ &). \end{split}$$

Show that

$$S \longrightarrow^{\leq 12} (\bar{a} \langle b \rangle . \operatorname{nil} \mid\mid \bar{b} \langle a \rangle . \operatorname{nil}) \mid\mid \operatorname{new} x(\operatorname{!new} s(\bar{x} \langle s \rangle . \bar{s} \langle a \rangle . \bar{s} \langle b \rangle . \operatorname{nil}))$$

where $\longrightarrow^{\leq 12}$ denotes at most 12 applications of the reaction relation.

Exercise 3 (Polyadic π -Calculus)

Consider the following process definition in polyadic π -calculus:

 $x(y_1, y_2) \cdot P \parallel \overline{x}\langle z_1, z_2 \rangle \cdot Q \parallel \overline{x}\langle z_1', z_2' \rangle \cdot Q'.$

Provide the corresponding encoding in monadic π -calculus. Furthermore, do at least two reduction sequences to the resulting process definition in order to convince yourself of the correctness of your translation.

(3 Points)

(4 Points)

(3 Points)