Static Program Analysis

Lecture 1: Introduction to Program Analysis

Thomas Noll

Lehrstuhl für Informatik 2
(Software Modeling and Verification)

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http://moves.rwth-aachen.de/teaching/ws-1415/spa/

Winter Semester 2014/15

Outline

People

- Lectures:
  - Thomas Noll (noll@cs.rwth-aachen.de)
  - Christina Jansen (christina.jansen@cs.rwth-aachen.de)
- Exercise classes:
  - Christian Dehnert (dehnert@cs.rwth-aachen.de)
  - Benjamin Kaminski (benjamin.kaminski@cs.rwth-aachen.de)
- Student assistant:
  - Frederick Prinz

Target Audience

- MSc Informatik:
  - Theoretische Informatik
- MSc Software Systems Engineering:
  - Theoretical Foundations

Expectations

- What you can expect:
  - Foundations of static analysis of computer software
  - Implementation and tool support
  - Applications in, e.g., program optimization and software validation

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  - Foundations of static analysis of computer software
  - Implementation and tool support
  - Applications in, e.g., program optimization and software validation
- What we expect: basic knowledge in
  - Programming (essential concepts of imperative and object-oriented programming languages and elementary programming techniques)
  - helpful: Theory of Programming (such as Semantics of Programming Languages or Software Verification)

Organization

- Schedule:
  - Lecture Mon 14:15–15:45 AH 1 (starting October 13)
  - Lecture Thu 14:15–15:45 AH 2 (starting October 23)
  - Exercise class Mon 10:15–11:45 AH 6 (starting October 27)
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- Written material in English, lecture and exercise classes "on demand", rest up to you

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What Is It All About?

Static (Program) Analysis

Static analysis is a general method for automated reasoning on artefacts such as requirements, design models, and programs.

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Distinguishing features:

Static: based on source code, not on (dynamic) execution
(in contrast to testing, profiling, or run-time verification)

Automated: "push-button" technology, i.e., little user intervention
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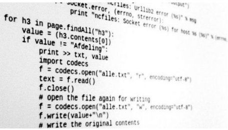
(Main) Applications:

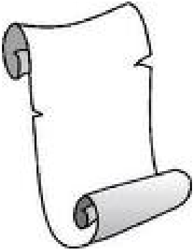
Optimizing compilers: exploit program properties to improve runtime or memory efficiency of generated code (dead code elimination, constant propagation, ...)

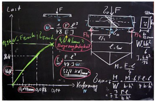
Software validation: verify program correctness (bytecode verification, shape analysis, ...)

Dream of Static Program Analysis

Program **Analyzer** **Result**







Property specification

Fundamental Limits

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Example 1.2 (Detection of constants)

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read(x);                    read(x);
if x > 0 then               if x > 0 then
  P;                          P;
  y := x;          ?~         y := x;
else                        else
  y := 1;                     y := 1;
end;                        end;
write(y);                   write(1);
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Thus: constant detection is undecidable

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   - employ abstract models of systems
     - finite automata, labeled transition systems, ...
   - perform exact analyses
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     - model checking, theorem proving, ...
2. **Weaker analyses** (here):
   - employ concrete models of systems
     - source code
   - perform approximate analyses
     - dataflow analysis, abstract interpretation, type checking, ...

Soundness vs. Completeness

- **Soundness:**
  - Predicted results must apply to every system execution
  - Examples:
    - constant detection: replacing expression by appropriate constant does not change program results
    - pointer analysis: analysis finds pointer variable $x \neq 0$ $\Longrightarrow$ no run-time exception when dereferencing x
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    - value of variable in $[0, 255]$ $\Longrightarrow$ interval analysis finds out
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- **Correctness** := Soundness $\wedge$ Completeness
  (often for logical axiomatizations and such, usually not guaranteed for program analyses)

Outline

Syntactic Categories

WHILE: simple imperative programming language without procedures or advanced data structures

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Syntactic categories:

Category	Domain	Meta variable
Numbers	$\mathbb{Z} = \{0, 1, -1, \ldots\}$	z
Truth values	$\mathbb{B} = \{\mathsf{true}, \mathsf{false}\}$	t
Variables	$Var = \{\mathtt{x}, \mathtt{y}, \ldots\}$	x
Arithmetic expressions	$AExp$ (next slide)	a
Boolean expressions	$BExp$ (next slide)	b
Commands (statements)	Cmd (next slide)	c

Syntax of WHILE Programs

Definition 1.3 (Syntax of WHILE)

The syntax of WHILE Programs is defined by the following context-free grammar:

$$
\begin{aligned}
a &::= z \mid x \mid a_1\texttt{+}a_2 \mid a_1\texttt{-}a_2 \mid a_1\texttt{*}a_2 \in AExp \\
b &::= t \mid a_1\texttt{=}a_2 \mid a_1\texttt{>}a_2 \mid \neg b \mid b_1 \wedge b_2 \mid b_1 \vee b_2 \in BExp \\
c &::= \texttt{skip} \mid x \texttt{ := } a \mid c_1\texttt{;}c_2 \mid \texttt{if } b \texttt{ then } c_1 \texttt{ else } c_2 \mid \texttt{while } b \texttt{ do } c \in Cmd
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Remarks: we assume that

- the syntax of numbers, truth values and variables is predefined (i.e., no "lexical analysis")
- the syntax of ambiguous constructs is uniquely determined (by brackets, priorities, or indentation)

A WHILE Program

Example 1.4

```
x := 6;
y := 7;
z := 0;
while x > 0 do
  x := x - 1;
  v := y;
  while v > 0 do
    v := v - 1;
    z := z + 1
```

A WHILE Program and its Flow Diagram

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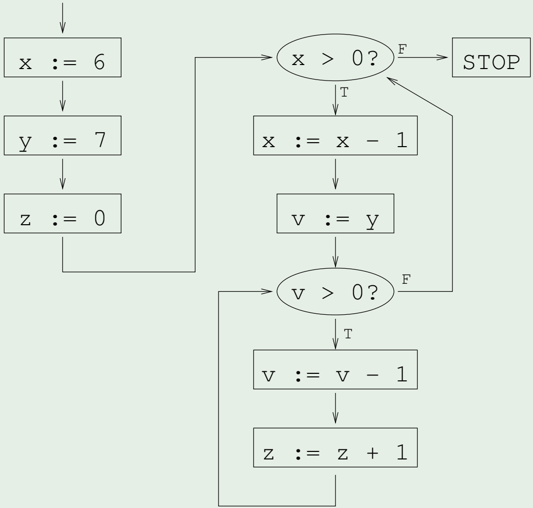

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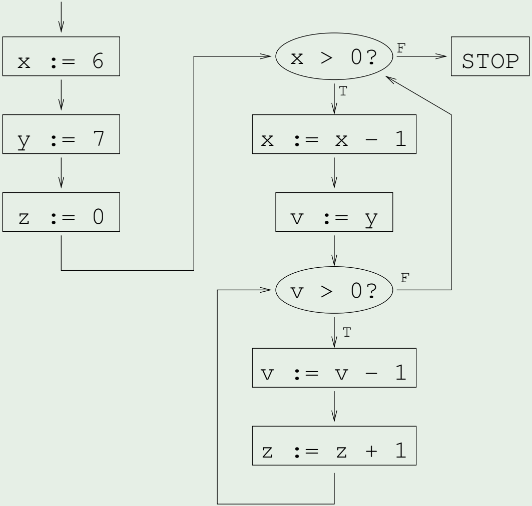

Effect: `z := x * y = 42`

Outline

(Preliminary) Overview of Contents

1. Introduction to Program Analysis
2. Dataflow analysis (DFA)
   1. Available expressions problem
   2. Live variables problem
   3. The DFA framework
   4. Solving DFA equations
   5. The meet-over-all-paths (MOP) solution
   6. Case study: Java bytecode verifier
3. Abstract interpretation (AI)
   1. Working principle
   2. Program semantics & correctness
   3. Galois connections
   4. Instantiations (sign analysis, interval analysis, ...)
   5. Case study: 16-bit multiplication
4. Interprocedural analysis
5. Pointer analysis

Outline

Additional Literature

- Flemming Nielson, Hanne R. Nielson, Chris Hankin: Principles of Program Analysis, 2nd edition, Springer, 2005
  [available in CS Library]
- Michael I. Schwartzbach: Lecture Notes on Static Analysis
  [http://www.itu.dk/people/brabrand/UFPE/Data-Flow-Analysis/static.pdf]
- Helmut Seidl, Reinhard Wilhelm, Sebastian Hack: Übersetzerbau 3: Analyse und Transformation, Springer, 2010
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