

Exercise 1 (Reaching Definitions Analysis using Abstract Interpretation): (8 Points)

We will now consider the Reaching Definitions Analysis (RDA). Given a labeled WHILE-program, this analysis computes for every program location and every variable all other program locations in which the variable might have been most recently written (i.e. written without being re-written in between). As an example, consider the following program.

```
[x := 2]1;  
[x := 3]2;  
while [y < 10]3;  
    [y := y + 1]4;  
[x := y * 2]5;
```

$(x, 2)$ is a reaching definition at label 4, because there is a path reaching label 4 such that x is most recently written at label 2. On the other hand $(x, 1)$ is not a reaching definition at label 4. If the most recent definition of a variable is "before the program", this is indicated by a question mark as the label information. For example, for label 5 we have the reaching definitions $\{(y, 4), (y, ?), (x, 2)\}$.

Your task is to implement this analysis in the abstract interpretation framework. Assume that you are given a **labeled** WHILE-program.

- Describe informally why you need to extend the concrete semantics of the WHILE-language presented in the lecture.
- Formally (re-)define the parts of the concrete semantics of WHILE that need to be changed.
- Formally adapt the abstract semantics of WHILE.
- Build the abstract transition system for your abstraction and the following WHILE-program.

```
[x := 2]1;  
[y := x + 2]2;  
while [x > 3]3  
    [y := y + x]4;  
[x := y]5;
```

- What do you ultimately have to do with the abstract transition system to derive the desired output of the analysis?

Exercise 2 (Lattice): (2 Points)

Consider the following predicates q_1 and q_2 . Calculate $q_1 \sqcup q_2$ and $q_1 \sqcap q_2$.

- $q_1 := p_1 \wedge p_2 \wedge p_3, q_2 := \neg p_3$
- $q_1 := p_1 \wedge \neg p_3, q_2 := p_1 \wedge p_2 \wedge \neg p_3$
- $q_1 := \neg p_1 \wedge \neg p_2 \wedge \neg p_3, q_2 := p_1 \wedge \neg p_2$