

Lehrstuhl für Informatik 2

Softwaremodellierung und Verifikation

Due date: May 12<sup>th</sup>. You can hand in your solutions at the start of the exercise class.

## Exercise 1 (Chain Complete Partial Orders)

Determine whether each of the following statements is true or false. For true statements present a formal proof, and for false statements provide a counterexample.

- (a) [7.5%] Every continuous function  $f: (D_1, \sqsubseteq_1) \to (D_1, \sqsubseteq_2)$  between two CCPOs  $(D_1, \sqsubseteq_1)$  and  $(D_2, \sqsubseteq_2)$  is monotonic.
- (b) [7.5%] Consider the partial order  $(\mathbb{Q}, \leq)$  of the rational numbers ordered by the natural order in the reals.  $(\mathbb{Q}, \leq)$  is chain complete.
- (c) [7.5%] If  $f: (D_1, \sqsubseteq_1) \to (D_1, \sqsubseteq_2)$  is a monotonic function between two CCPOs and  $D \subseteq D_1$  is a chain, then  $f(\bigsqcup D) \sqsubseteq_2 \bigsqcup f(D)$ .
- (d) [7.5%] Let  $(D, \Box)$  be a partial order and let  $f: (D, \Box) \to (D, \Box)$  be monotonic. If p is the least element in D satisfying  $f(p) \sqsubseteq p$ , then p is a fixed point of f.

Exercise 2 (repeat-until Loops)

(a) [10%] Define a transformer  $F: (\Sigma \dashrightarrow \Sigma) \to (\Sigma \dashrightarrow \Sigma)$  such that

$$\mathfrak{C}[\operatorname{repeat} c \text{ until } b]] = \operatorname{fix}(F)$$
.

The transformer F is allowed to depend on the semantics only of c and b (i.e.  $\mathfrak{B}[b]$ and  $\mathfrak{C}[c]$ . You cannot rely on the existence of while-loops within the language to define F.

(b) [5%] Use the definition provided in (a) to compute the transformer  $\hat{F}: (\Sigma \dashrightarrow \Sigma) \rightarrow \Sigma$  $(\Sigma \rightarrow \Sigma)$  whose least fixed point gives the semantics of program repeat skip until false. In other words, compute  $\hat{F}$  such that

 $\mathfrak{C}$ [repeat skip until false] = fix( $\hat{F}$ ).

(c) [10%] Show that  $\operatorname{fix}(\hat{F}) = f_{\emptyset}$ .

## Exercise 3 (Closed Sets)

A set  $C \subseteq D$  is *closed* if and only if for each chain  $G \subseteq C$ ,  $||G \in C$ . In the following, let  $(D, \sqsubseteq)$  be a chain complete partial order and  $f: D \to D$  be a continuous function. Prove the following two statements.

- (a) [7.5%] For each closed set  $C \subseteq D$  with  $f(x) \in C$  for each  $x \in C$ , we have fix $(f) \in C$ .
- (b) [7.5%]  $f(x) \sqsubseteq x$  implies fix $(f) \sqsubseteq x, x \in D$ .

## Exercise 4 (Pointwise Ordering)

Let  $(D, \sqsubseteq)$  be a CCPO and define  $(D \to D, \sqsubseteq')$  by setting

$$f_1 \sqsubseteq' f_2$$
 if and only if  $f_1(d) \sqsubseteq f_2(d)$  for all  $d \in D$ 

(a) [10%] Show that  $(D \to D, \Box')$  is a CCPO.



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(b) [20%] Show that fixpoints of chains are "continous", i.e.

$$\mathsf{fix}(\bigsqcup' \mathcal{F}) = \bigsqcup\{\mathsf{fix}(f) \mid f \in \mathcal{F}\}\$$

holds for all non-empty chains  $\mathcal{F}\subseteq D\to D$  of continous functions.