Introduction to Model Checking 2015: Exercise 3.

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Hand in before: 6th May

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We will use the following nomenclature in the rest of the exercise. For a set A, $\mathcal{P}(A)$ is the power set of A. The universe of infinite words (or strings) is Σ^{ω} where Σ is the alphabet. The set of finite words is Σ^+ . The i^{th} element of a word t is denoted by t[i]. A finite word x is a *prefix* of a word y (finite or infinite) is denoted as $x \prec y$. The linear time closure operator $\mathsf{cl} : \mathcal{P}(\Sigma^{\omega}) \to \mathcal{P}(\Sigma^{\omega})$ is defined as:

 $\mathsf{cl}(T) = \{ t \in \Sigma^{\omega} \mid \forall x \prec t \exists t' \in T : x \prec t' \}.$

We say, $A \subseteq \Sigma^{\omega}$ is safe iff cl(A) = A and A is live iff $cl(A) = \Sigma^{\omega}$.

Exercise 1

Prove the following simple equivalences:

- 1. $cl(\emptyset) = \emptyset$.
- 2. $A \subseteq B$ implies $cl(A) \subseteq cl(B)$.
- 3. $T \subseteq \mathsf{cl}(T)$.
- 4. cl(T) = cl(cl(T)).
- 5. $\operatorname{cl}(A \cup B) = \operatorname{cl}(A) \cup \operatorname{cl}(B)$.

Exercise 2

Show the following:

- 1. A singleton set is safe.
- 2. Consider the set $R_a = \{ t \in \Sigma^{\omega} \mid \exists i \in \mathbb{N} : t[i] = a \}$, where $a \in \Sigma$. R_a is not safe but it is live.
- 3. $A \cup \overline{cl(A)}$ is live for any set A. (where \overline{B} is the set complement of B.)

Exercise 3

Give examples of:

- 1. If A and B are live then $A \cap B$ is not live.
- 2. A set T such that $T \subsetneq cl(T)$.
- 3. The distributiveness of safe over \cup (union) can be easily extended to finite union that is: Finite union of safe sets are safe.

$$\mathsf{cl}(A_1\cup\ldots\cup A_n)=A_1\cup\ldots\cup A_n.$$

Show by an example that the statement does not hold for countable union, i.e.:

$$\mathsf{cl}\left(\bigcup_{i\in\omega}A_i\right) \neq \bigcup_{i\in\omega}A_i.$$

Hint: Exercise 2.2.

(4 points)

(2 points)

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