Compiler Construction

Lecture 1: Introduction

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http://moves.rwth-aachen.de/teaching/ss-14/cc14/

Summer Semester 2014







- Lectures:
 - Thomas Noll (noll@cs.rwth-aachen.de)
- Exercise classes:
 - Friedrich Gretz (fgretz@cs.rwth-aachen.de)
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- Student assistant:
 - Philipp Berger
 - Samiro Discher



• BSc Informatik:

- Wahlpflicht Theoretische Informatik
- MSc Informatik:
 - Theoretische Informatik
- MSc Software Systems Engineering:
 - Theoretical Foundations of SSE (was: Theoretical CS)



• What you can expect:

- how to implement (imperative) programming languages
- application of theoretical concepts
- compiler = example of a complex software architecture
- gaining experience with tool support
- What we expect: basic knowledge in
 - imperative programming languages
 - algorithms and data structures
 - formal languages and automata theory

- Schedule:
 - Lecture Mon 14:15–15:45 AH 6 (starting 14 April)
 - Lecture Wed 10:15–11:45 AH 6 (starting 9 April)
 - Exercise class Fri 08:15–09:45 AH 2 (starting 16 April)
 - Special: 16 April (exercise), 2/4 June (itestra)
 - see overview at http://moves.rwth-aachen.de/teaching/ss-14/cc14/
- 1st assignment sheet next week, presented 25 April
- Work on assignments in groups of 2-3 people
- Written exams (2 h, 6 Credits) on 25 July/3 September
- Admission requires at least 50% of the points in the exercises
- Written material in English, lecture and exercise classes in German, rest up to you









Compiler = Program: Source code \rightarrow Target code

Source code: in high-level programming language, tailored to problem

- imperative vs. declarative (functional, logic) vs. object-oriented
- sequential vs. concurrent

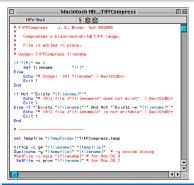
Target code: low-level code, tailored to machine

- platform-independent byte code (for virtual machine such as JVM)
- platform-dependent assembly/machine code (RISC/CISC/parallel/...)

Usage of Compiler Technology I

Programming language interpreters

- Ad-hoc implementation of small programs in scripting languages (perl, bash, ...)
- Programs usually interpreted, i.e., executed stepwise
- Moreover: many non-scripting languages also involve interpreters (e.g., JVM as byte code interpreter)



Usage of Compiler Technology II

Web browsers

- Receive HTML (XML) pages from web server
- Analyse (parse) data and translate it to graphical representation

```
<!DOCTYPE html PUBLIC "-//W3C//DTD HTML
2
   <html>
3
        <head>
4
              <title>Example</title>
5
              k href="screen.css" rel="sty
         </head>
6
7
8
9
         <body>
              \langle h1 \rangle
                   <a href="/">Header</a>
10
              </hi>
11
              \langle 1i \rangle
12
13
                        <a href="one/">One</a>
14
                   \langle /1i \rangle
15
                   \langle 1 i \rangle
16
                        <a href="two/">Two</a>
17
                   \langle /1i \rangle
```

Usage of Compiler Technology III

Text processors

• ETEX = "programming language" for texts of various kinds

• Translated to DVI, PDF, ...

```
\documentclass[12pt] {article}
%options include 12pt or 11pt or 10pt
%classes include article, report, book, letter, thesis
\title{This is the title}
\author Author One \\ Author Two}
\date{\today}
\begin {document}
\maketitle
This is the content of this document.
This is the 2nd paragraph.
Here is an inline formula:
$V=\frac{4 \pi r^3}{3}$
And appearing immediately below
is a displayed formula:
$$V=\frac{4 \pi r^3}{3}$$
\end{document}
```



Efficiency of generated code

Goal: target code as fast and/or memory efficient as possible

- program analysis and optimization
- cf. course on Static Program Analysis (WS 2012/13, 2014/15)

Efficiency of compiler

Goal: translation process as fast and/or memory efficient as possible

(for inputs of arbitrary size)

- fast (linear-time) algorithms
- sophisticated data structures

Correctness

Goals: conformance to source and target language specifications;

"equivalence" of source and target code

- compiler validation and verification
- proof-carrying code, ...
- cf. course on Semantics and Verification of Software (SS 2013, 2015)

Remark: mutual tradeoffs!

Aspects of a Programming Language

yntax: "How does a program look like?"

• hierarchical composition of programs from structural components

Semantics: "What does this program mean?"

"Static semantics": properties which are not (easily) definable in syntax (declaredness of identifiers, type correctness, ...)

"Dynamic semantics": execution evokes state transformations of an (abstract) machine

Pragmatics

- length and understandability of programs
- learnability of programming language
- appropriateness for specific applications

• ...

Motivation for Rigorous Formal Treatment

Example

```
From NASA's Mercury Project: FORTRAN DO loop
     • D0 5 K = 1,3: DO loop with index variable K
     • DO 5 K = 1.3: assignment to (real) variable DO5K
I How often is the following loop traversed?
                      for i := 2 to 1 do ...
   FORTRAN IV: once
      PASCAL: never
What if p = nil in the following program?
             while p <> nil and p^.key < val do ...
        Pascal: strict Boolean operations 4
       Modula: non-strict Boolean operations \checkmark
```



Code generation: since 1940s

- ad-hoc techniques
- concentration on back-end
- first FORTRAN compiler in 1960

Formal syntax: since 1960s

- LL/LR parsing
- shift towards front-end
- semantics defined by compiler/interpreter

Formal semantics: since 1970s

- operational
- denotational
- axiomatic
- cf. course on Semantics and Verification of Software

Automatic compiler generation: since 1980s

- [f]lex, yacc, ANTLR, action semantics, ...
- cf. http://catalog.compilertools.net/



Lexical analysis (Scanner):

- recognition of symbols, delimiters, and comments
- by regular expressions and finite automata

Syntax analysis (Parser):

- determination of hierarchical program structure
- by context-free grammars and pushdown automata

Semantic analysis:

- checking context dependencies, data types, ...
- by attribute grammars

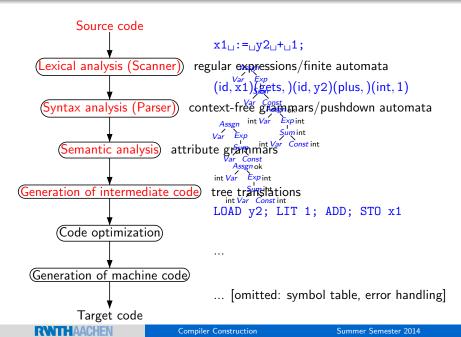
Generation of intermediate code:

- translation into (target-independent) intermediate code
- by tree translations

Code optimization: to improve runtime and/or memory behavior Generation of target code: tailored to target system Additionally: optimization of target code, symbol table, error handling

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Conceptual Structure of a Compiler



Analysis vs. synthesis

Analysis: lexical/syntax/semantic analysis (determination of syntactic structure, error handling)

Synthesis: generation of (intermediate/machine) code + optimization

Front-end vs. back-end

Front-end: machine-independent parts (analysis + intermediate code + machine-independent optimizations)

Back-end: machine-dependent parts (generation + optimization of machine code)

Historical: *n*-pass compiler

- *n* = number of runs through source program
- nowadays mainly one-pass

Literature

(CS Library: "Handapparat Softwaremodellierung und Verifikation")

General

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Special

- O. Mayer: Syntaxanalyse, BI-Wissenschafts-Verlag, 1978
- D. Brown, R. Levine T. Mason: *lex & yacc*, O'Reilly, 1995
- T. Parr: The Definite ANTLR Reference, Pragmatic Bookshelf, 2007

Historical

- W. Waite, G. Goos: Compiler Construction, 2nd edition, Springer, 1985
- N. Wirth: Grundlagen und Techniken des Compilerbaus, Addison-Wesley, 1996

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